Trade-off between incrementality and accuracy Seminar on Incremental Processing

Benjamin Weitz

June 30, 2011

Kato et al.

Summary

Outline









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Summary

Introduction

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Summary

Introduction

Incremental Dialogue Systems:

• fast(er)

- fast(er)
- sometimes wrong decisions

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- sometimes wrong decisions \Rightarrow revisions

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- components depend on each other

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Idea: decrease incrementality a bit to reduce revisions

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• Baumann et. al: Automatic Speech Recognition

Incremental Dialogue Systems:

- fast(er)
- $\bullet \ \text{sometimes wrong decisions} \Rightarrow \text{revisions}$
- components depend on each other

Idea: decrease incrementality a bit to reduce revisions

- Baumann et. al: Automatic Speech Recognition
- Kato et. al: Parsing

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Summary

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• Define measures to evaluate incremental ASR-Systems

- Define measures to evaluate incremental ASR-Systems
- Evaluate an existing system with these measures

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- Use the measures to improve ASR

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Definitions

Hypothesis at time t: $hyp_t = w_{hyp_t}$

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How to evaluate the quality of hyp_t ?

use **actually spoken input** as gold standard

use **final hypothesis of ASR** as gold standard

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How to evaluate the quality of hyp_t ?

use **actually spoken input** as gold standard use **final hypothesis of ASR** as gold standard

Why?

more meaningful:

- relates partial hypothesis to what can be expected from ASR
- correct interpretation might never be recognized

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Summary

Relative Correctness

w_{gold}: final, non incremental hypothesis

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Summary

Relative Correctness

```
w_{gold}: final, non incremental hypothesis
```

Relative Correctness

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Summary

Relative Correctness

w_{gold} : final, non incremental hypothesis

Relative Correctness

w_{gold}	sil	(e	Z	zwei			drei		•••		
time:	0 1 3	2 3	4	5	67	8	9	10	11	12	
w_{hyp_1}	sil									1	
w_{hyp_2}	sil										
w_{hyp_3}	sil										
w_{hyp_4}	sil	an									
w_{hyp_5}	sil	ei	n	j							
w_{hyp_6}	sil	Ē	ins	-	1						
w_{hyp_7}	sil	eiı	ns	z١	vei						
w_{hyp_8}	sil	E	ins	-	Zw	ar					
w_{hyp_9}	sil	eins			Z	zwei					
$w_{hyp_{10}}$	sil	eins			Ľ	zwei					
$w_{hyp_{11}}$	sil	eins			Z	zwei		sil			
$w_{hyp_{12}}$	sil	eins			Z	zwei		dr	ei	j	
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Summary

Relative Correctness

w_{gold} : final, non incremental hypothesis

Relative Correctness



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Summary

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Summary

Relative Correctness

```
w_{gold}: final, non incremental hypothesis
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Relative Correctness

 w_{hyp_t} is relatively correct, iff $w_{hyp_t} = w_{gold_t}$

Prefix Correctness

 w_{hyp_t} is prefix-correct, iff w_{hyp_t} is a prefix of w_{gold_t}

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Summary

Relative Correctness

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Edit Overhead

Three ways for hyp_{t+1} to differ from hyp_t :

- extension
- revokation
- revision

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- extension
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- revision

- \bullet add message: \oplus
- *revoke* message: ⊖

Edit Overhead

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- extension: \oplus
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Edit Overhead

Three ways for hyp_{t+1} to differ from hyp_t :

- \bullet extension: \oplus
- \bullet revokation: \ominus
- revision

- \bullet add message: \oplus
- revoke message: \ominus

Edit Overhead

Three ways for hyp_{t+1} to differ from hyp_t :

- extension: \oplus
- \bullet revokation: \ominus
- revision: \ominus , \oplus

- \bullet add message: \oplus
- *revoke* message: ⊖

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Edit Overhead



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Summary

Edit Overhead



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Summary

Edit Overhead



Perfect ASR-System: 1 extension for each word
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Edit Overhead



Perfect ASR-System: 1 extension for each word

Edit Overhead (EO)

rate of spurious edits

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Summary



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Summary



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Summary

Timing Measures

Word First Correct Response (WFC)

The first time a word appears in the correct position

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Summary

Timing Measures

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The first time a word appears in the correct position

w_{gold}	sil	eins			z	zwei		dre	ei)	•
time:	0 1 3	2 3	4	5	6 7	8	9	10 1	11 1	2	
w_{hyp_1}	sil										
w_{hyp_2}	sil										
w_{hyp_3}	sil										
w_{hyp_4}	sil	an	ן								
w_{hyp_5}	sil	ei	n	j							
w_{hyp_6}	sil		ins	_							
w_{hyp_7}	sil	ei	ns	Żzv	vei						
w_{hyp_8}	sil	C.	ins	_	zw	ar					
$w_{hyp_{\mathbf{q}}}$	sil	Ċ	eins	-	z	wei					
$w_{hyp_{10}}$	sil	C.	ins	_		zw	ei				
$w_{hyp_{11}}$	sil	Ľ	ins	_	Z	wei	Ľ.	sil	1		
$w_{hyp_{12}}$	sil	Ċ	ins	_	Z	wei	Ì	dre	ei	j	
÷	1 1		i T		\mathcal{Q}_{a}	1		T.	i	i	
			v	$V \Gamma$	\cup_{zw}	4 I	w F	I zwei			

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Summary

Timing Measures

Word First Correct Response (WFC)

The first time a word appears in the correct position

w_{gold}	sil	eins	zwe	i drei	
time:	0 1 3	2 3 4 5	5678	9 10 11	12
w_{hyp_1}	sil				
w_{hyp_2}	sil				
w_{hyp_3}	sil				
w_{hyp_4}	sil	an			
w_{hyp_5}	sil	ein			
w_{hyp_6}	sil	eins			
w_{hyp_7}	sil	eins	zwei		
w_{hyp_8}	sil	eins	zwar		
w_{hyp_9}	sil	eins	zwe		
$w_{hyp_{10}}$	sil	eins	ZW	/ei	
$w_{hyp_{11}}$	sil	eins	zwe	i <i>sil</i>	
$w_{hyp_{12}}$	sil	eins	zwe	i drei	
÷		יייי	VFC	WFF	'
			a C zuci	1 1 1 20VCS	

 $WFC_{zwei} = 7$

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Summary

Timing Measures

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The first time a word appears in the correct position

Word First Final Response (WFF)

The time a hypothesis remains stable / doesn't change anymore

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w_{gold}	sil	eins			Z	we	ei 🛛	dr	ei		
time:	0 1 2	2 3	4	5	6 7	7 8	3 9	10	11	12	
w_{hyp_1}	sil									1	
w_{hyp_2}	sil										
w_{hyp_3}	sil										
w_{hyp_4}	sil	an	ך								
w_{hyp_5}	sil	ei	n	j							
w_{hyp_6}	sil	e	ins	-							
w_{hyp_7}	sil	eir	ıs	Ζv	vei						
w_{hyp_8}	sil	e	ins	-	zw	ar					
w_{hyp_9}	sil	e	ins		Z	we	i				
$w_{hyp_{10}}$	sil	e	ins			z٧	/ei				
$w_{hyp_{11}}$	sil	e	ins		Z	we	i	sil			
$w_{hyp_{12}}$	sil	e	ins		Z	we	i	dı	ei	ב	
÷			' Ţ	' VF	C	ei I	WI	FF			

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w_{gold}	sil	ei	ns		zwe	ei 🛛	drei]
time:	0 1 2	2 3	45	6	78	8 9	10 11	12
w_{hyp_1}	sil							
w_{hyp_2}	sil							
w_{hyp_3}	sil							
w_{hyp_4}	sil	an						
w_{hyp_5}	sil	eir						
w_{hyp_6}	sil	ei	ns					
w_{hyp_7}	sil	ein	s i	zwei	j			
w_{hyp_8}	sil	ei	ns	z١	var			
w_{hyp_9}	sil	ei	ns		zwe	ei		
$w_{hyp_{10}}$	sil	ei	ns	T	zv	vei	1	
$w_{hyp_{11}}$	sil	ei	ns		zwe	i.	sil	
$w_{hyp_{12}}$	sil	ei	ns	Ţ	zwe	i	drei	
÷	1 1		1 1 147	EC.	J		i i T	1
			VV.	$r \phi_z$	wei	VV I	I zwei	

 $WFF_{zwei} = 9$

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Summary

Timing Measures

Word First Correct Response (WFC)

The first time a word appears in the correct position

Word First Final Response (WFF)

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Correction Time (CT)

CT = WFF - WFC

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CT = WFF - WFC

$w_{\scriptscriptstyle gold}$	sil	eins		zwei		drei	
time:	0 1 3	2 3 4	56	78	8 9	10 11	12
w_{hyp_1}	sil						
w_{hyp_2}	sil						
w_{hyp_3}	sil	()					
w_{hyp_4}	sil	an					
w_{hyp_5}	sil	ein	j				
w_{hyp_6}	sil	eins					
w_{hyp_7}	sil	eins	zwe	ei.			
w_{hyp_8}	sil	eins	Z	war			
w_{hyp_9}	sil	eins	Ţ	zwe	i		
$w_{hyp_{10}}$	sil	eins	#	zv	vei		
$w_{hyp_{11}}$	sil	eins		zwe	ei Ì	sil 🛛	
$w_{hyp_{12}}$	sil	eins	Ì	zwe	i (drei	
÷		''''''''''	VFC	۲.	WF	F.	ı
		,		zuci		- zwes	

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w_{gold}	sil	eins		zwe	ei 🛛	drei]
time:	0 1 2	2 3 4	5 6	78	8 9	10 11	12
w_{hyp_1}	sil						
w_{hyp_2}	sil						
w_{hyp_3}	sil						
w_{hyp_4}	sil	an					
w_{hyp_5}	sil	ein]				
w_{hyp_6}	sil	eins	\square				
w_{hyp_7}	sil	eins	zw	ei			
w_{hyp_8}	sil	eins		zwar			
w_{hyp_9}	sil	eins	Ì	zwe	ei		
$w_{hyp_{10}}$	sil	eins	\neg	zv	vei	ן ב	
$w_{hyp_{11}}$	sil	eins		zwe	i).	sil	
$w_{hyp_{12}}$	sil	eins	Ţ	zwe	i	drei	
÷			WFC		WF	F	·
				wwest.		19900	

 $CT_{zwei} = 9 - 7 = 2$

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Summary

Setup and Data

- continious speech framework Sphinx-4
- acoustic model
 - German
 - instructions in a puzzle domain
- trigram language model
- test data
 - 85 recodings
 - two speakers
 - sentence similar to training sentences

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Summary

SER	68.2%
WER	18.8%
r-correct	30.9%
p-correct	53.1%
edit overhead	90.5%
mean word duration	0.378 s
WFC	mean: 0.276 s
WFF	mean: 0.004 s
immediately correct	58.6%

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Summary

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• rather low correctness

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Summary

Measurements



58.6% immediately correct

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Summary

Measurements



58.6% immediately correct

90% correct after a correction time of 320 ms

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Summary

Measurements



58.6% immediately correct

90% correct after a correction time of 320 ms

95% correct after a correction time of 550 ms

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Summary

Independency of the measures?

Are the measures independent of specific settings?

- vary LM/AM-weight
- vary audio quality by adding white noise

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Summary

Independency of the measures?

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- Define measures to evaluate incremental ASR-Systems
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Summary

Which improvements?

main goal: improve edit overhead

- reduce amount of wrong hypotheses
- still as quick as possible

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Summary

Right Context

allow larger right context of size Δ :

- at time t: take into account output of ASR until $t \Delta$ only
- $hyp_{t-\Delta}$ has a lookahead up to t

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Summary

Right Context

allow larger right context of size Δ :

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- \Rightarrow reduction of edit overhead
- \Rightarrow hypothesis lags behind the gold standard
 - WFC increases by Δ
 - effects on correctness, because w_{gold_t} may contain more words

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Fair R-Correctness

 w_{hyp_t} is fairly r-correct, iff $w_{hyp_{t-\Delta}} = w_{gold_{t-\Delta}}$

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Summary

Right Context



correctness and EO improve with more right context

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Summary

Right Context



correctness and EO improve with more right context

timing measures increase with larger right context

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Summary

Right Context



correctness and EO improve with more right context

timing measures increase with larger right context

percentage of immediately correct hypotheses increases: 90% @ 580 ms 98% @ 1060 ms
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Message Smoothing

edit message must be result of N consecutive hypotheses before commitment

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Summary

Message Smoothing

edit message must be result of N consecutive hypotheses before commitment

N = 2:

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Summary

Message Smoothing

edit message must be result of N consecutive hypotheses before commitment

N = 2:

• *an, ein* and *zwar* would never be committed

w_{gold}	sil	6	eins		Z	Γ	drei				
time:	0 1 2	2 3	4	5	6 7	8	9	10	11	12	
w_{hyp_1}	sil										
w_{hyp_2}	sil										
w_{hyp_3}	sil										
w_{hyp_4}	sil	an									
w_{hyp_5}	sil	ei	n	j							
w_{hyp_6}	sil	e	eins	-							
w_{hyp_7}	sil	eii	ns	Ζv	vei						
w_{hyp_8}	sil	€	eins	-	zw	ar					
w_{hyp_9}	sil	Ē	eins	-	Ż	wei					
$w_{hyp_{10}}$	sil	e	eins	-	Ľ	zwe	ei	ב			
$w_{hyp_{11}}$	sil	e	ins		Z	wei	1	sil)		
$w_{hyp_{12}}$	sil	Ē	ins		Z	zwei		dre	ei	ב	
÷				' NF	r	. v	↓ VF	F		1	
					~ zwe	s /		- zwc	8		

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Message Smoothing

edit message must be result of *N* consecutive hypotheses before commitment

N = 2:

- *an, ein* and *zwar* would never be committed
- only 3 edit messages



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Summary

Message Smoothing



edit overhead falls rapidly: 50% after 110 ms 10% after 320 ms

smoothing in s (scale shows larger smoothings towards the left)

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Message Smoothing



edit overhead falls rapidly: 50% after 110 ms 10% after 320 ms decreasing (strict and fair)

r-correctness

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Summary

Message Smoothing



edit overhead falls rapidly: 50% after 110 ms 10% after 320 ms decreasing (strict and fair)

r-correctness

increasing timing measures

- Define measures to evaluate incremental ASR-Systems
- Evaluate an existing system with these measures
- Use the measures to improve ASR
- Conclusions

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Summary

Conclusions

goal: improve edit overhead

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Summary

Conclusions

goal: improve edit overhead

Right Context

improvements with larger delays, increasing correctness

Message Smoothing

improvements with shorter delays, decreasing correctness

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Summary

Conclusions

goal: improve edit overhead

Right Context

improvements with larger delays, increasing correctness

Message Smoothing

improvements with shorter delays, decreasing correctness

could be combined to yield a good effect

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• ITAG: incremental-parsing-oriented tree adjoining grammar

- ITAG: incremental-parsing-oriented tree adjoining grammar
- PITAG: probabilistic ITAG

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- PITAG: probabilistic ITAG
- Validity of Partial Parse Trees

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- Experimental Results and Conclusions

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Summary

Elementary Trees

Initial Trees

must be *leftmost-expanded*:

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1. $[t]_X$

t: terminal symbol

X: nonterminal symbol

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Х

t

Summary

Elementary Trees

Initial Trees

must be *leftmost-expanded*:

- 1. $[t]_X$
 - t: terminal symbol
 - X: nonterminal symbol









Auxiliary Trees

 $[X^* \sigma X_1 \cdots X_k]_X$ σ : leftmost expanded tree X, X_1, \dots, X_k : nonterminal symbols

Summary



$$[X^* \sigma X_1 \cdots X_k]_X$$

 σ : leftmost expanded tree
 X, X_1, \dots, X_k : nonterminal symbols



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Summary

Operations

Substitution

replaces a leftmost nonterminal leaf of a partial parse tree σ with an initial tree α having the same nonterminal symbol at its root

 s_{α} : substituting α $s_{\alpha}(\sigma)$: result of applying s_{α} to σ

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Summary

Operations

Substitution

replaces a leftmost nonterminal leaf of a partial parse tree σ with an initial tree α having the same nonterminal symbol at its root

```
s_{lpha}: substituting lpha
s_{lpha}(\sigma): result of applying s_{lpha} to \sigma
```



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Summary

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replaces a leftmost nonterminal leaf of a partial parse tree σ with an initial tree α having the same nonterminal symbol at its root

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replaces a leftmost nonterminal leaf of a partial parse tree σ with an initial tree α having the same nonterminal symbol at its root

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s_{lpha}: substituting lpha
s_{lpha}(\sigma): result of applying s_{lpha} to \sigma
```



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Summary

Operations

Adjunction

```
splits a partial parse tree \sigma at a nonterminal node having no
nonterminal leaf and inserts an auxiliary tree \beta having the
same nonterminal symbol at its root
a_{\beta}: adjoining \beta
```

```
a_{\beta}(\sigma): result of applying a_{\beta} to \sigma
```

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Summary

Operations

Adjunction

splits a partial parse tree σ at a nonterminal node having no nonterminal leaf and inserts an auxiliary tree β having the same nonterminal symbol at its root a_{β} : adjoining β

 $a_{\beta}(\sigma)$: result of applying a_{β} to σ



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Adjunction

splits a partial parse tree σ at a nonterminal node having no nonterminal leaf and inserts an auxiliary tree β having the same nonterminal symbol at its root a_{β} : adjoining β

 $a_{\beta}(\sigma)$: result of applying a_{β} to σ



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splits a partial parse tree σ at a nonterminal node having no nonterminal leaf and inserts an auxiliary tree β having the same nonterminal symbol at its root a_{β} : adjoining β

 $a_{\beta}(\sigma)$: result of applying a_{β} to σ



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Summary

Parsing with ITAG

at *i*-th word *w_i*:

Parsing with ITAG

at *i*-th word *w_i*:

combine (substitution, adjunction) elementary trees for w_i with partial parse trees for w₁ ··· w_{i-1}
Parsing with ITAG

at *i*-th word w_i:

combine (substitution, adjunction) elementary trees for w_i with partial parse trees for w₁ ··· w_{i-1}

 \Rightarrow partial parse trees for $w_1 \cdots w_i$

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Summary

Example



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Summary





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Summary



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Summary



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Summary



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word	#	partial parse tree
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Ι	2	$[[[I]_{prp}]_{np}vp]_s$
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constructing parse trees of initial fragments for every word input possible!

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Summary

Tree Construction

How to get the elementary trees?

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Summary

Tree Construction

How to get the elementary trees?

• Extraction from a treebank!

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- ITAG: incremental-parsing-oriented tree adjoining grammar
- PITAG: probabilistic ITAG
- Validity of Partial Parse Trees
- Experimental Results and Conclusions

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Summary

Probabilities

probabilities for events combining an elementary tree and another tree by **substitution and adjunction**

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Summary

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lpha: initial tree with root symbol X

Probability of substituting lpha

 $P(s_{\alpha}) = \frac{|s_{\alpha}| \text{ in the treebank}}{|\text{substitutions using other initial trees with root X| in the treebank}}$

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Probabilities

probabilities for events combining an elementary tree and another tree by **substitution and adjunction**

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 $P(s_{\alpha}) = \frac{|s_{\alpha}| \text{ in the treebank}}{|\text{substitutions using other initial trees with root X}| \text{ in the treebank}}$

β : auxiliary tree with root symbol X

Probability of adjoining β				
$P(a_eta)=rac{ a_eta $ in the treebank}{ X in the treebank				

probability of a parse-tree:

• product of probabilities of operations used for construction

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S

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s_{α_1}	1.0
s_{α_2}	0.7
s_{α_5}	0.3
s_{α_7}	0.5

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1.0

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Summary

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 $1.0\,\times\,0.7$

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Summary

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Summary

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probability of a parse-tree:

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 $1.0 \times 0.7 \times 0.3 \times 0.5 = 0.105$

Parsing with PITAG

Improve efficiency by:

- discard tree with lower probability when there are parse trees with same possible operations
- only keep n-best partial parse-trees

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- ITAG: incremental-parsing-oriented tree adjoining grammar
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Definitions

 σ, τ : partial parse trees

Subsumption

 σ subsumes τ , iff τ can be constructed by applying substitutionand adjunction-operations to σ and every parse tree subsumes itself

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Subsumption

 σ subsumes τ , iff τ can be constructed by applying substitutionand adjunction-operations to σ and every parse tree subsumes itself

Validity

a partial parse tree is valid for an input sentence, **iff** it subsumes the correct parse tree for the input sentence

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Evaluating Validity

validity for a partial parse tree depends on the rest of the sentence

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Conditional Validity of partial parse tree σ $V(\sigma|w_1 \cdots w_j) = \frac{\sum (\text{Probabilities of partial parse trees at } w_j \text{ subsumed by } \sigma)}{\sum (\text{Probabilities of partial parse trees constructed at } w_j)}$

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example: later

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Output Partial Parse Trees

delay output until validity is high enough:

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Output Partial Parse Trees

delay output until validity is high enough:

Parse Tree to be returned

Parse Tree with the longest initial fragment whose validity is greater than threshold $\boldsymbol{\theta}$

heta: threshold between 0 and 1

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Example



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Summary

Example



Parse Tree	Probability
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 $\theta = 0.8$

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When will #3 (valid) be output?

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- ITAG: incremental-parsing-oriented tree adjoining grammar
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Summary

Experimental Setup

- parser implemented in Lisp
- input: POS-Sequences
- elementary trees extracted from Penn Treebank
- only used correctly parsed sentences

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Summary

Measures

Degree of delay at j-th word

 $D(j,s) = j - |\text{output parse tree at } w_j|$

 $s = w_1 \cdots w_n$: Input sentence

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Summary

Measures

Degree	of	delay	at	j-th	word
--------	----	-------	----	------	------

 $D(j,s) = j - |\text{output parse tree at } w_j|$

 $s = w_1 \cdots w_n$: Input sentence

Precision

percentage of valid partial parse trees in the output

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Summary

 D_{ave}

6.4

2.9

2.2

1.8

1.3

0.9

0.0

Results and Conclusions

delay(number of words)



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Results and Conclusions

delay(number of words)



higher threshold

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Results and Conclusions

delay(number of words)



higher threshold \Rightarrow higher precision, greater delays

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Results and Conclusions

delay(number of words)



 $\begin{array}{l} \mbox{higher threshold} \\ \Rightarrow \mbox{ higher precision, greater delays} \\ \Rightarrow \mbox{ trade-off between precision and delay} \end{array}$

Kato et al.

Summary

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Baumann et. al

use right context and message smoothing to reduce/avoid wrong hypotheses in incremental ASR Summary

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delay of output ⇒ Trade-off between incrementality/speed and output quality (revisions) Introduction

Baumann et al.

Kato et al.

Summary

Thank you!

Thanks for your attention!

References

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