

Linguistic Inference and Textual Entailment

Modelling Inference
through Logical Deduction
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Model-theoretic Interpretation

- Formula A is **true in the model structure M** iff $\llbracket A \rrbracket^{M,g} = 1$ for every variable assignment g .
- A model structure M **satisfies** a set of formulas Γ (or: M is a **model** of Γ) iff every formula $A \in \Gamma$ is true in M .



Levels of Logical Method

- Model-theoretic interpretation
- Deduction calculi
- Deduction procedures
- Implemented deduction systems



Central semantic concepts

- A formula A is **valid** ($\models A$) iff A is true in every model structure.
- A set of formulas Γ **entails** formula A ($\Gamma \models A$) iff A is true in every model of Γ (i.e., in every model structure that satisfies Γ).
- A set of formulas Γ is **satisfiable** iff Γ has a model (i.e., there is a model structure that satisfies Γ).



Important Theorems

... actually, **metatheorems** (see below):

- Validity and entailment:

$A \models B$ iff $\models A \rightarrow B$, more general:

$\{A_1, \dots, A_n\} \models B$ iff $\models A_1 \wedge \dots \wedge A_n \rightarrow B$

- Entailment and satisfiability:

$\Gamma \models A$ iff $\Gamma \cup \{\neg A\}$ is unsatisfiable.



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Deduction Calculi

- Computing entailment and other logical concepts through semantic interpretation inefficient and in many cases infeasible.
- Deduction calculi (or **proof theoretic systems**) provide a strictly syntactic way of checking logical concepts and relations, through symbol manipulation/ rewrite of logical formulas.



Axioms and Deduction Rules

- Deduction calculi are typically made up of (1) axioms and (2) deduction rules.
- Example for a frequently used axiom:
 - $A \vee \neg A$ (“Tertium non datur”)
- Example for a frequently used deduction rule (“Modus Ponens”)

$$\frac{A \rightarrow B, A}{B}$$



Semantic and Deductive Concepts

- There is a correspondence between basic semantic and deductive/ proof-theoretic concepts:

Validity	Provability
Entailment	Derivability/Deducibility
Satisfiability	Consistency



Soundness and Completeness

- **Soundness:** If $\Gamma \vdash A$, then $\Gamma \models A$.
- **Completeness:** If $\Gamma \models A$, then $\Gamma \vdash A$.



Central proof-theoretic concepts

- Formula A is derivable (deducible) from a set of formulas Γ ($\Gamma \vdash A$) iff there is a sequence of formulas A_1, \dots, A_n such that $A_n = A$ and for all members A_i of the sequence: either
 - A_i is an (instantiation of an) axiom, or
 - $A_i \in \Gamma$, or
 - A_i is the result of the application of a deduction rule, whose conclusion is A_i , and whose premisses all occur in the sequence before A_i
- A formula A is **provable** ($\vdash A$) iff $\emptyset \vdash A$
- A set of formulas Γ is **inconsistent** iff there is a formula A such that $\Gamma \vdash A$ and $\Gamma \vdash \neg A$
- A set of formulas Γ is **consistent** iff it is not inconsistent.



Important Metatheorems

- **Derivability and Provability:**
 $\{A_1, \dots, A_n\} \vdash B$ iff $\vdash A_1 \wedge \dots \wedge A_n \rightarrow B$
- **Derivability and Consistency:**
 $\Gamma \models A$ iff $\Gamma \cup \{\neg A\}$ is inconsistent.
- **Validity and Provability:**
 $\models A$ iff $\vdash A$
- **Satisfiability and Consistency:**
 Γ is satisfiable iff Γ is consistent.



Deduction Calculi

- There is one model-theoretic interpretation (for standard predicate logic).
- There is a wide variety of deduction calculi, e.g.:
 - Hilbert calculus
 - Semantic tableau calculus
 - Calculus of natural deduction (Gentzen calculus)
 - Resolution



Semantic Tableau Calculus

- Derivation and proofs through the generation of tableau trees via decomposition rules.
- Semantic Tableaus use rewrite on formulas, so it is a deduction calculus.
- They are called “semantic tableaus” because there is an affinity to semantics.



Semantic Tableau Rules

	Affirmed	Negated
$A \wedge B$	$\{A, B\}$	$\{\neg A\}, \{\neg B\}$
$A \vee B$	$\{A\}, \{B\}$	$\{\neg A, \neg B\}$
$A \rightarrow B$	$\{\neg A\}, \{B\}$	$\{A, \neg B\}$
$A \leftrightarrow B$	$\{A \rightarrow B, B \rightarrow A\}$	$\{\neg(A \rightarrow B)\}, \{\neg(B \rightarrow A)\}$
$\forall xA$	$A[a/x]$ for arbitrary a	$\neg A[a/x]$ for a new a
$\exists xA$	$A[a/x]$ for a new a	$\neg A[a/x]$ for arbitrary a



Semantic Tableau Calculus

- A subtableau is closed, iff it contains A and $\neg A$
- A tableau is closed iff all subtableaus are closed.
- $\Gamma \models A$ iff the decomposition rules result in a closed tableau for $\Gamma \cup \{\neg A\}$.

Refutation proof: To prove A from premisses Γ , add its negation and show that the result is inconsistent.



Deduction procedures

- Deduction procedure = deduction calculus + algorithm

Tractability:

- Propositional calculus is NP-complete (it requires exponential time)
- FOL is undecidable (provable /valid formulas are recursively enumerable)
- To arrive at efficient systems, heuristic knowledge and a lot of fine-tuning is required.



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Implemented deduction systems

We distinguish:

- **Theorem provers**, typically with
 - Refutation proofs
 - **Resolution** proof procedure
 - Input: Set of formulas (premisses)
 - Output: Yes, if proof successful.
 - Examples: Vampire, SPASS, BLIKSEM, OTTER
- Interactive theorem provers (“**proof assistants**”)
 - Provide information about proof steps
 - Ask for guidance
 - Are typically based on more intuitive calculi (e.g. Gentzen calculus)
 - Example: OMEGA
- **Model generators**
 - Check consistency
 - Using tableau techniques
 - Output is Yes, if the hypothesis is consistent with the premisses
 - Plus a model for $\Gamma \cup \{A\}$ as an important side effect.
 - Examples: MACE, KIMBA



Problems: Efficiency

- Combination of Theorem Provers (and Model Builders), Distributed Theorem Proving
- Optimization for specific tasks (e.g., mathematical vs. linguistic applications)
- Restriction to a FOL fragment, with
 - Horn Clause Logic (Prolog) and
 - **Description Logics** (e.g., RACER) as prominent examples



- Required input is logical formulas
- Available linguistic input is text
 - Grammatical analysis, semantic construction
 - Disambiguation, Underspecification
 - Discourse analysis (e.g., coreference resolution)
- Additional input required comprises
 - Lexical semantic information (e.g., WordNet, FrameNet)
 - Extralinguistic Knowledge