

Suffix Trees

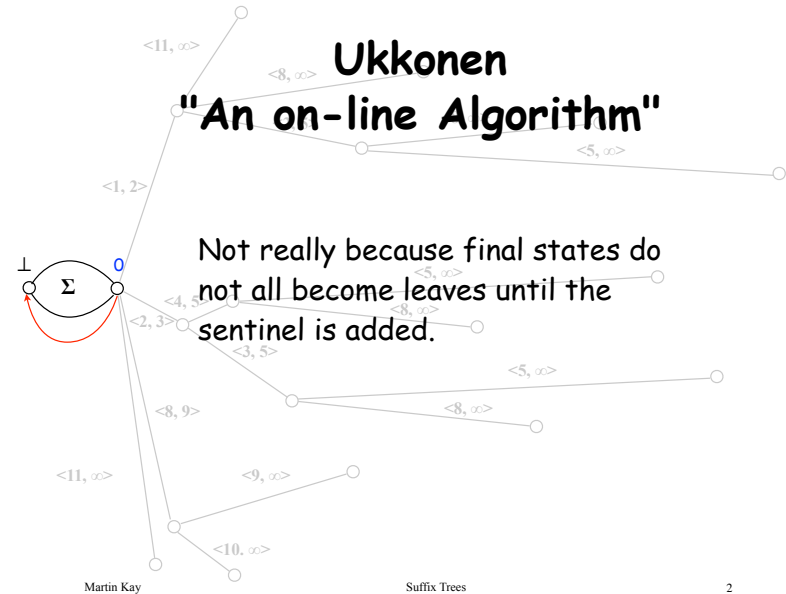
Ukkonen's Algorithm

Martin Kay

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and the University of the Saarland

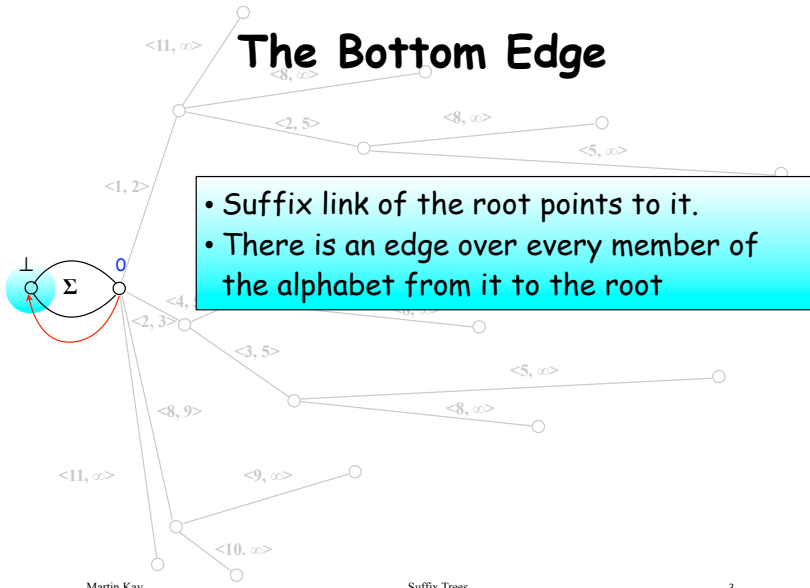
Ukkonen "An on-line Algorithm"

Not really because final states do not all become leaves until the sentinel is added.



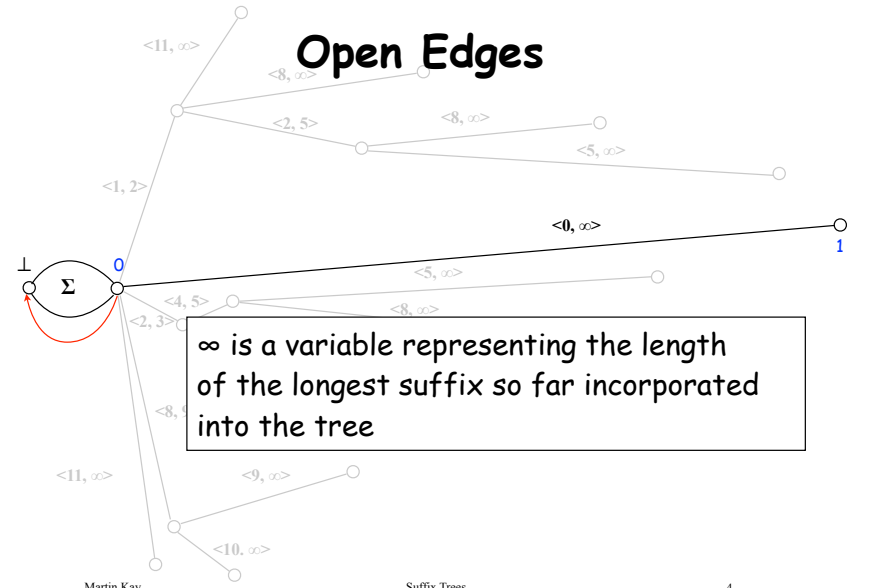
The Bottom Edge

- Suffix link of the root points to it.
- There is an edge over every member of the alphabet from it to the root



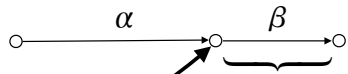
Open Edges

∞ is a variable representing the length of the longest suffix so far incorporated into the tree



The Active Point

The longest string (α) that is in the tree both as a suffix and elsewhere.



Additions in here are implicit in open edges

So we can look for places to split edges working back from here.

The End Point

If the current character already begins the label on some edge on the boundary path, then it will begin the label on every subsequent edge on the boundary path because, if x follows a α , it also follows α .

So we can stop looking for places to split once we encounter such a node.

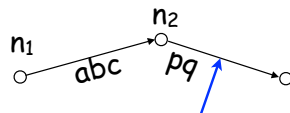
Reference Pairs

References to real and virtual nodes.

(n, s) where n is a real node and s is a string.

If $s = \epsilon$, then (n, s) is a real node.

The reverse does not hold:



$$(n_1 \text{ abcpq}) = (n_2 \text{ pq})$$

Canonical form: $(n_2 \text{ pq})$ — shortest string

The Active point

