

Agenda

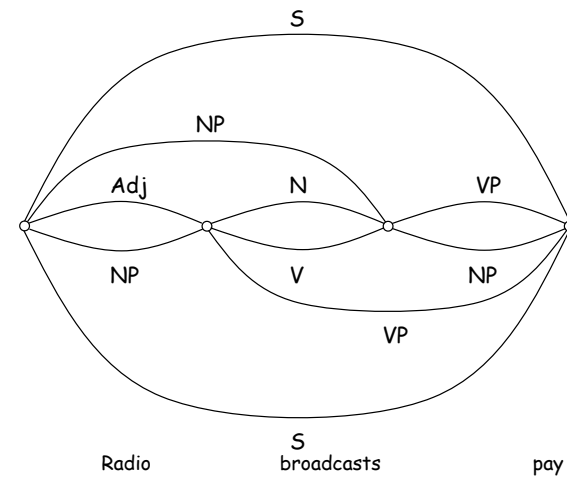
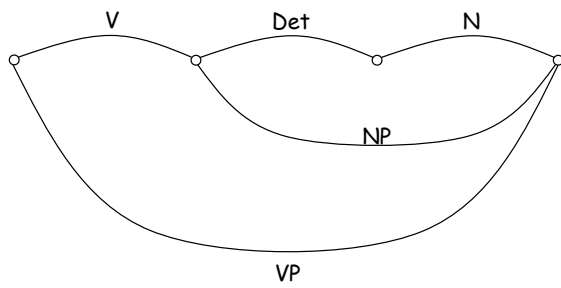
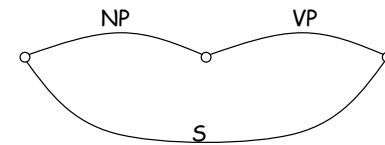
and

Chart

Martin Kay

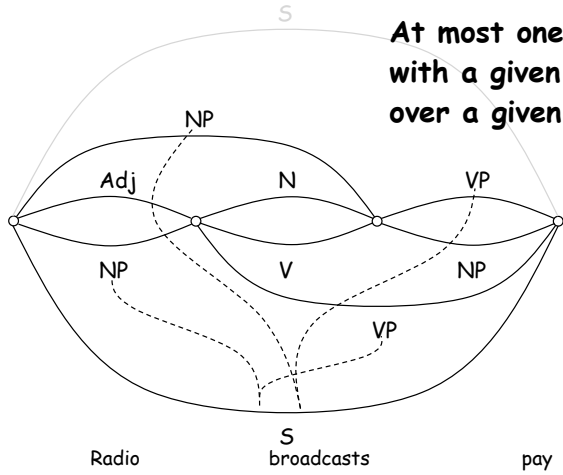
Stanford University and
the University of the Saarland

Charts

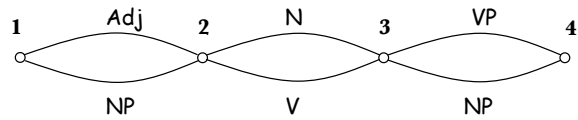
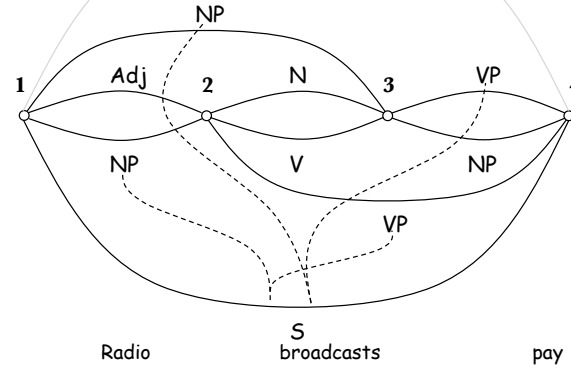


Structure

At most one edge with a given label over a given span



S



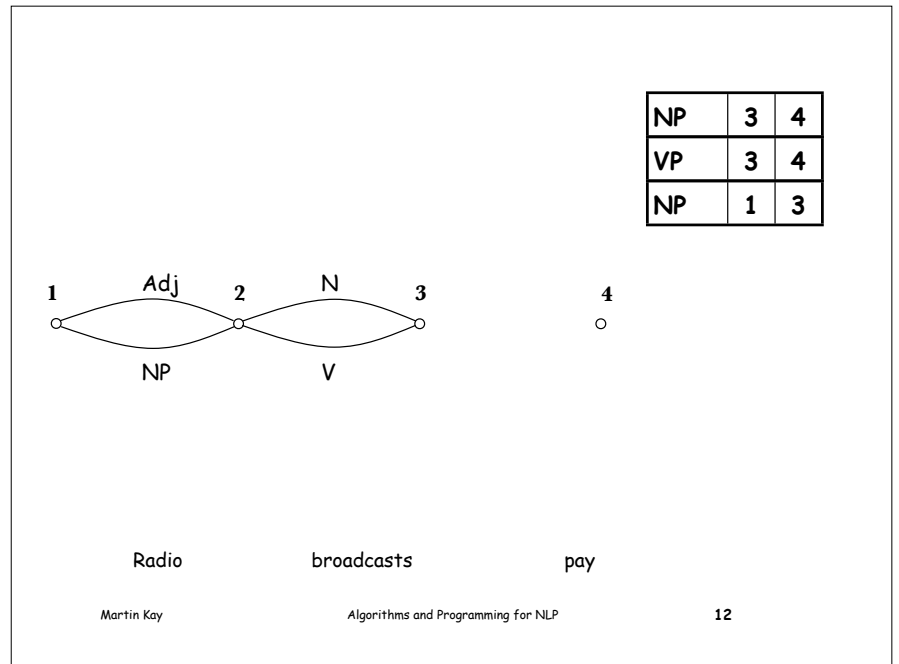
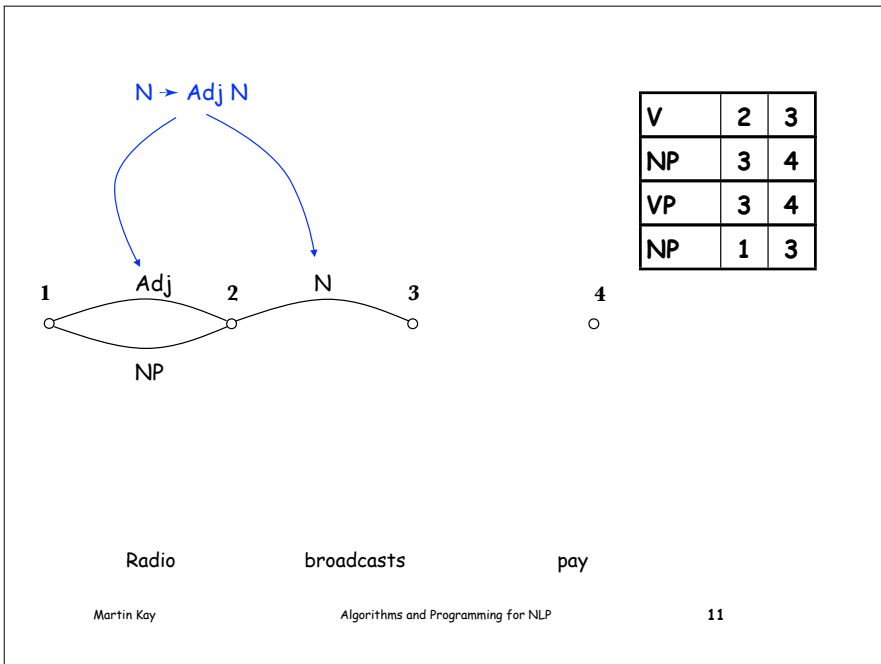
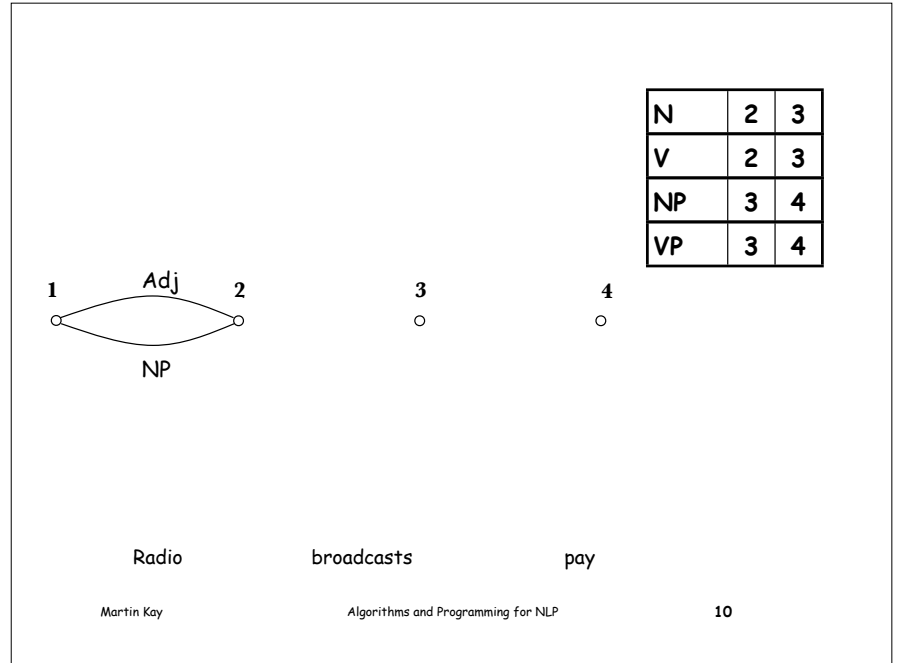
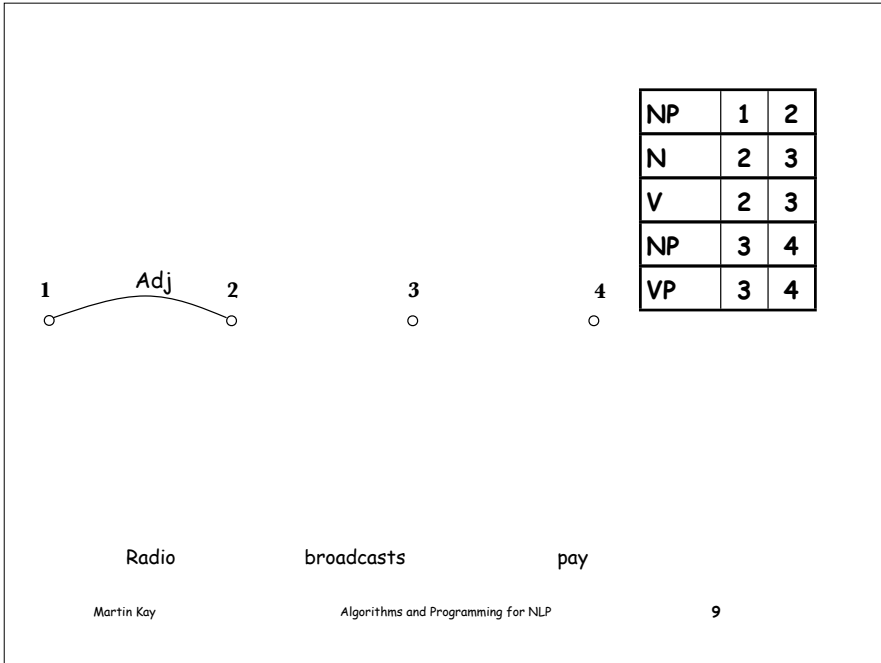
Radio broadcasts pay

Adj	1	2
NP	1	2
N	2	3
V	2	3
NP	3	4
VP	3	4

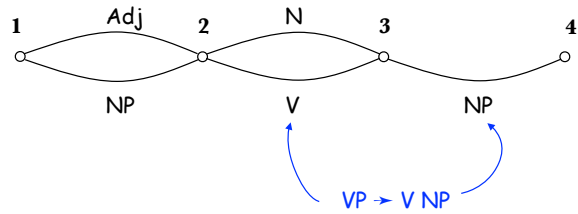


Radio broadcasts pay

Adj	1	2
NP	1	2
N	2	3
V	2	3
NP	3	4
VP	3	4

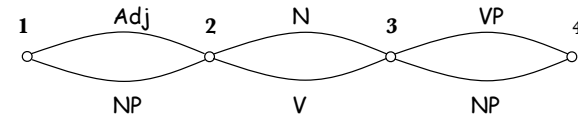


VP	3	4
NP	1	3
VP	2	4



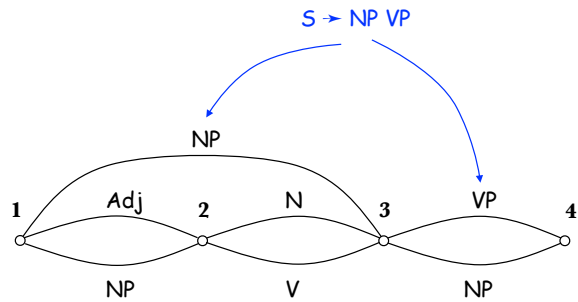
Radio broadcasts pay

NP	1	3
VP	2	4



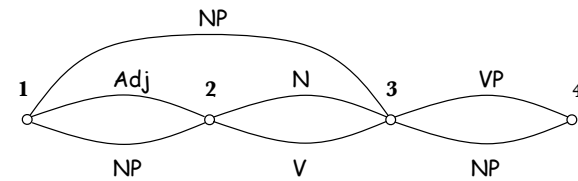
Radio broadcasts pay

VP	2	4
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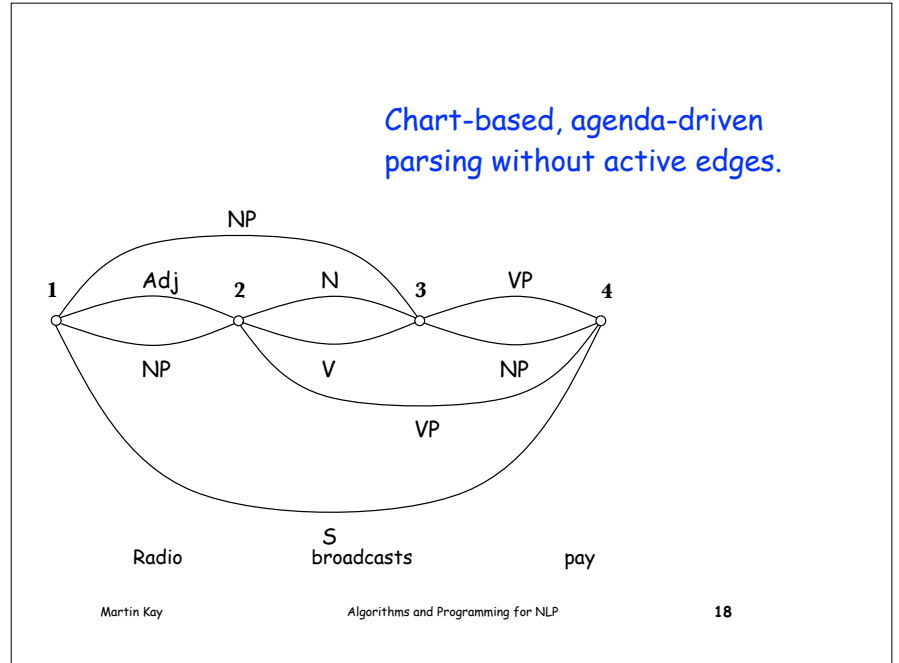
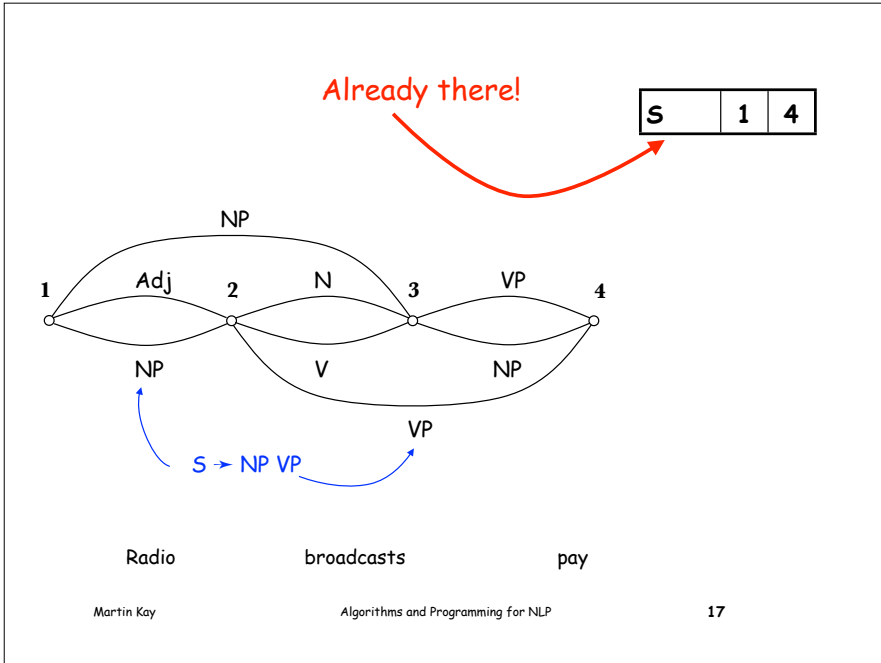


Radio broadcasts pay

VP	2	4
S	1	4



Radio broadcasts pay



Pseudo CKY

- **When a new edge is entered in the chart, apply**
 - apply all eligible rules and the edges to its left and right.
 - apply unary rules
- **This is pseudo-CKY because**
 - the order of events is not constrained to be length-first, left to right, etc.
 - CKY does not provide for unary rules (though it requires only a small addition)

Martin Kay Algorithms and Programming for NLP 19

<u>Grammar</u>			<u>Agenda</u>
$S \rightarrow NP VP$	$NP \rightarrow they$	$N \rightarrow planes$	0-1 they
$NP \rightarrow N$	$V \rightarrow are$	$Be \rightarrow are$	1-2 are
$N \rightarrow Adj N$	$Adj \rightarrow flying$	$PrPt \rightarrow flying$	2-3 flying
$VP \rightarrow V NP$			3-4 planes
$V \rightarrow Be PrPt$			

Chart

Martin Kay Algorithms and Programming for NLP 20

Grammar

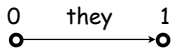
NP → they
 S → NP VP
 NP → N
 N → Adj N
 VP → V NP
 V → Be PrPt

PrPt → flying
 Be → are
 V → are
 N → planes

Agenda

1-2 are
 2-3 flying
 3-4 planes

Chart



Grammar

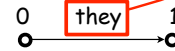
NP → they
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 N → Adj N
 VP → V NP
 V → Be PrPt

PrPt → flying
 Be → are
 V → are
 N → planes

Agenda

1-2 are
 2-3 flying
 3-4 planes
 0-1 NP

Chart



Grammar

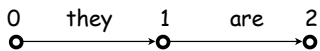
NP → they
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 N → Adj N
 VP → V NP
 V → Be PrPt

PrPt → flying
 Be → are
 V → are
 N → planes

Agenda

2-3 flying
 3-4 planes
 0-1 NP

Chart



Grammar

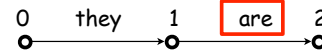
NP → they
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 N → Adj N
 VP → V NP
 V → Be PrPt

PrPt → flying
 Be → are
 V → are
 N → planes

Agenda

2-3 flying
 3-4 planes
 0-1 NP
 1-2 V
 1-2 Be

Chart



Grammar

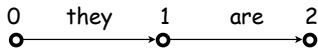
NP → they
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 N → Adj N
 VP → V NP
 V → Be PrPt

NP → they
 N → planes
 V → are
 Be → are
 Adj → flying
 PrPt → flying

Agenda

2-3 flying
 3-4 planes
 0-1 NP
 1-2 V
 1-2 Be

Chart



Grammar

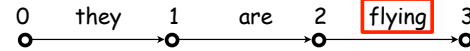
NP → they
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 VP → V NP
 V → Be PrPt

NP → they
 N → planes
 V → are
 Be → are
 Adj → flying
 PrPt → flying

Agenda

3-4 planes
 0-1 NP
 1-2 V
 1-2 Be
 2-3 Adj
 2-3 PrPt

Chart



Grammar

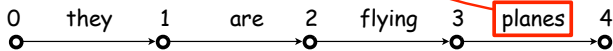
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NP → they
 N → planes
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Agenda

0-1 NP
 1-2 V
 1-2 Be
 2-3 Adj
 2-3 PrPt
 3-4 N

Chart



Grammar

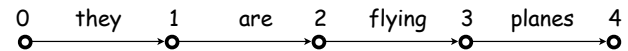
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NP → they
 N → planes
 V → are
 Be → are
 Adj → flying
 PrPt → flying

Agenda

0-1 NP
 1-2 V
 1-2 Be
 2-3 Adj
 2-3 PrPt
 3-4 N

Chart



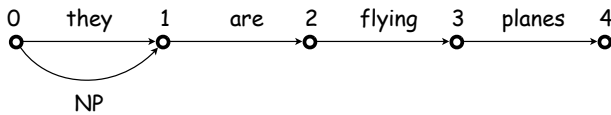
Grammar

NP → they
 S → NP VP
 N → planes
 NP → N
 V → are
 N → Adj N
 Be → are
 VP → V NP
 Adj → flying
 V → Be PrPt
 PrPt → flying

Agenda

1-2 V
 1-2 Be
 2-3 Adj
 2-3 PrPt
 3-4 N

Chart



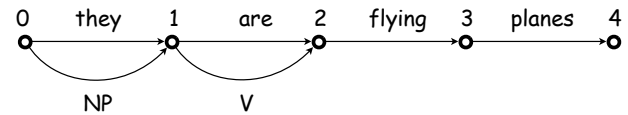
Grammar

NP → they
 S → NP VP
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 NP → N
 V → are
 N → Adj N
 Be → are
 VP → V NP
 Adj → flying
 V → Be PrPt
 PrPt → flying

Agenda

1-2 Be
 2-3 Adj
 2-3 PrPt
 3-4 N

Chart



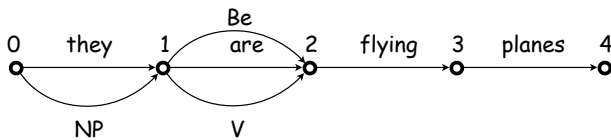
Grammar

NP → they
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 N → planes
 NP → N
 V → are
 N → Adj N
 Be → are
 VP → V NP
 Adj → flying
 V → Be PrPt
 PrPt → flying

Agenda

2-3 Adj
 2-3 PrPt
 3-4 N

Chart



No interactions

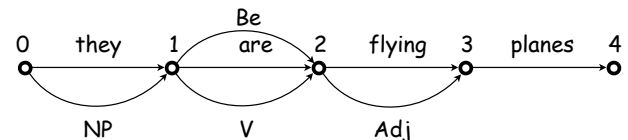
Grammar

NP → they
 S → NP VP
 N → planes
 NP → N
 V → are
 N → Adj N
 Be → are
 VP → V NP
 Adj → flying
 V → Be PrPt
 PrPt → flying

Agenda

2-3 PrPt
 3-4 N

Chart



No interactions

Grammar NP → they

S → NP VP N → planes

NP → N V → are

N → Adj N Be → are

VP → V NP Adj → flying

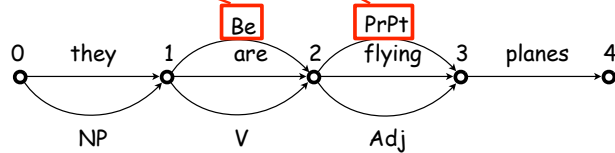
V → Be PrPt PrPt → flying

Agenda

3-4 N

1-3 V

Chart



Grammar NP → they

S → NP VP N → planes

NP → N V → are

N → Adj N Be → are

VP → V NP Adj → flying

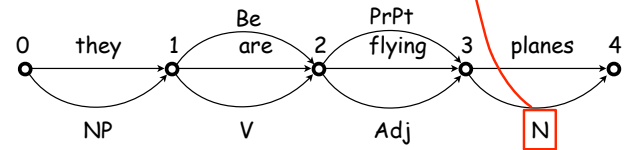
V → Be PrPt PrPt → flying

Agenda

1-3 V

3-4 NP

Chart



Grammar NP → they

S → NP VP N → planes

NP → N V → are

N → Adj N Be → are

VP → V NP Adj → flying

V → Be PrPt PrPt → flying

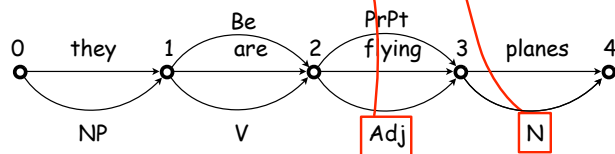
Agenda

1-3 V

3-4 NP

2-4 N

Chart



Grammar NP → they

S → NP VP N → planes

NP → N V → are

N → Adj N Be → are

VP → V NP Adj → flying

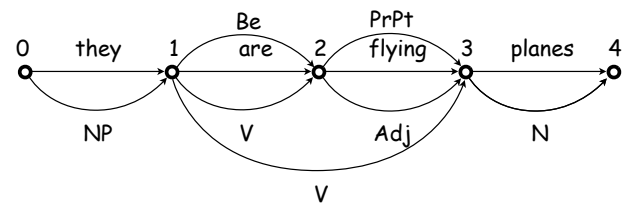
V → Be PrPt PrPt → flying

Agenda

3-4 NP

2-4 N

Chart



Grammar NP → they

S → NP VP

NP → N

N → Adj N

VP → V NP

V → Be PrPt

N → planes

V → are

Be → are

Adj → flying

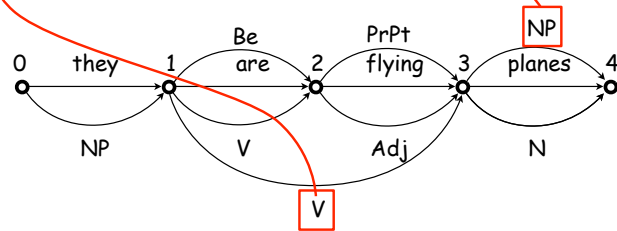
PrPt → flying

Agenda

2-4 N

1-4 VP

Chart



Grammar NP → they

S → NP VP

NP → N

N → Adj N

VP → V NP

V → Be PrPt

N → planes

V → are

Be → are

Adj → flying

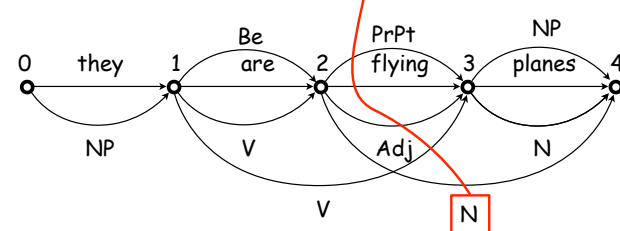
PrPt → flying

Agenda

1-4 VP

2-4 NP

Chart



Grammar NP → they

S → NP VP

NP → N

N → Adj N

VP → V NP

V → Be PrPt

N → planes

V → are

Be → are

Adj → flying

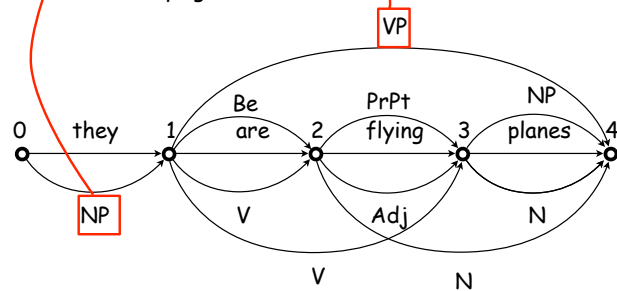
PrPt → flying

Agenda

2-4 NP

0-4 S

Chart



Grammar NP → they

S → NP VP

NP → N

N → Adj N

VP → V NP

V → Be PrPt

N → planes

V → are

Be → are

Adj → flying

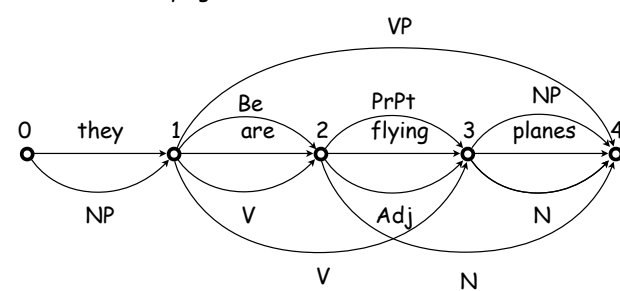
PrPt → flying

Agenda

2-4 NP

0-4 S

Chart



Grammar NP → they

S → NP VP N → planes

NP → N V → are

N → Adj N Be → are

VP → V NP Adj → flying

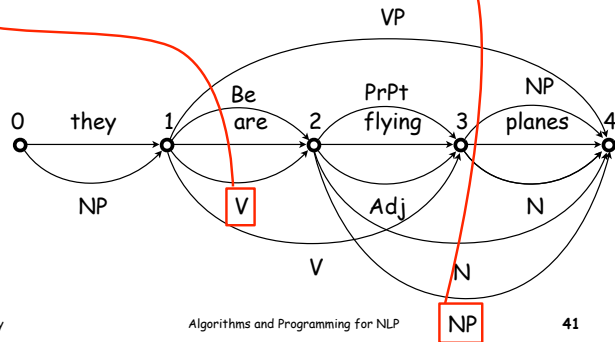
V → Be PrPt PrPt → flying

Agenda

0-4 S

1-4 VP

Chart



Grammar NP → they

S → NP VP N → planes

NP → N V → are

N → Adj N Be → are

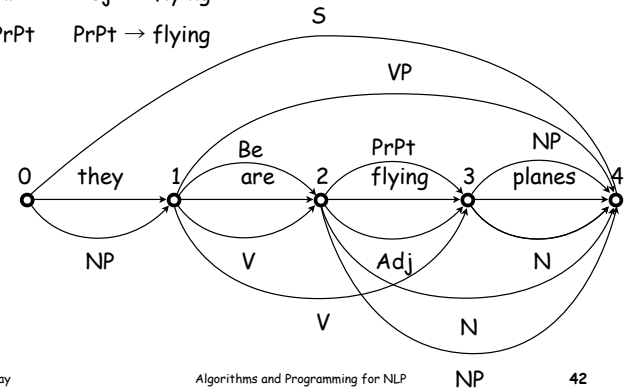
VP → V NP Adj → flying

V → Be PrPt PrPt → flying

Agenda

1-4 VP

Chart



Grammar NP → they

S → NP VP N → planes

NP → N V → are

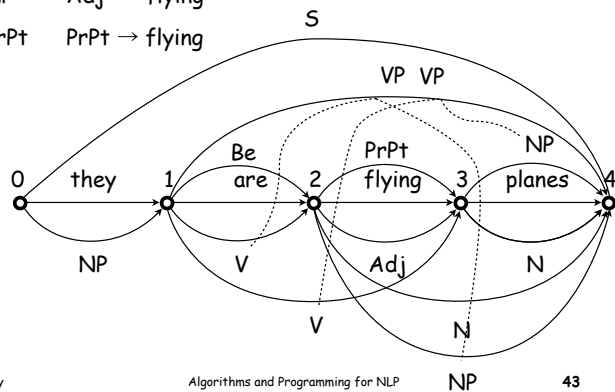
N → Adj N Be → are

VP → V NP Adj → flying

V → Be PrPt PrPt → flying

Agenda

Chart



Active Edges

The Chart Invariant— Active and Inactive



Every $\forall \alpha$ and β where

- α is active and β is inactive, and
- α is incident to, and β from, the same vertex, α is applied to β , and any resulting vertices added to the agenda

The Fundamental Rule

Grammar

$NP \rightarrow they$	$NP \rightarrow they$
$S \rightarrow NP VP$	$N \rightarrow planes$
$NP \rightarrow N$	$V \rightarrow are$
$N \rightarrow Adj N$	$Be \rightarrow are$
$VP \rightarrow V NP$	$Adj \rightarrow flying$
$V \rightarrow Be PrPt$	$PrPt \rightarrow flying$

Dotted Rules

Agenda

- 0-0 $S \rightarrow \cdot NP VP$
- 0-1 they
- 1-2 are
- 2-3 flying
- 3-4 planes

Chart



Grammar

$NP \rightarrow they$	$NP \rightarrow they$
$S \rightarrow NP VP$	$N \rightarrow planes$
$NP \rightarrow N$	$V \rightarrow are$
$N \rightarrow Adj N$	$Be \rightarrow are$
$VP \rightarrow V NP$	$Adj \rightarrow flying$
$V \rightarrow Be PrPt$	$PrPt \rightarrow flying$

Agenda

- 0-1 they
- 1-2 are
- 2-3 flying
- 3-4 planes

Chart $S \rightarrow \cdot NP VP$



Grammar

$NP \rightarrow they$	$NP \rightarrow they$
$S \rightarrow NP VP$	$N \rightarrow planes$
$NP \rightarrow N$	$V \rightarrow are$
$N \rightarrow Adj N$	$Be \rightarrow are$
$VP \rightarrow V NP$	$Adj \rightarrow flying$
$V \rightarrow Be PrPt$	$PrPt \rightarrow flying$

Agenda

- 0-0 $NP \rightarrow \cdot N$
- 0-1 they
- 1-2 are
- 2-3 flying
- 3-4 planes

Chart $S \rightarrow \cdot NP VP$

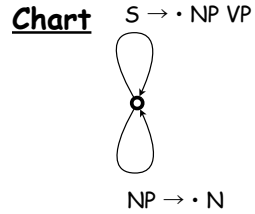


Grammar

NP → they
 S → NP VP N → planes
 NP → N V → are
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 VP → V NP Adj → flying
 V → Be PrPt PrPt → flying

Agenda

0-1 they
 1-2 are
 2-3 flying
 3-4 planes

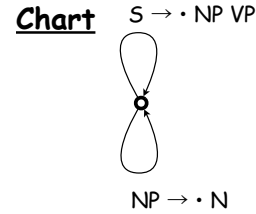


Grammar

NP → they
 S → NP VP N → planes
 NP → N V → are
 N → Adj N Be → are
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Agenda

0-0 planes ...
 0-1 they
 1-2 are
 2-3 flying
 3-4 planes



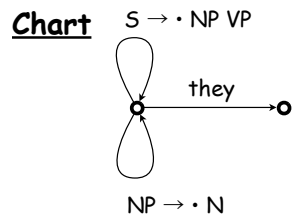
Oops! Can't handle lexical items top-down

Grammar

NP → they
 S → NP VP N → planes
 NP → N V → are
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 VP → V NP Adj → flying
 V → Be PrPt PrPt → flying

Agenda

1-2 are
 2-3 flying
 3-4 planes

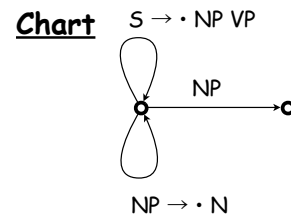


Grammar

NP → they
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Agenda

1-2 are
 2-3 flying
 3-4 planes

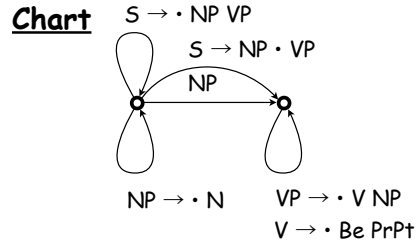


Grammar

NP → they
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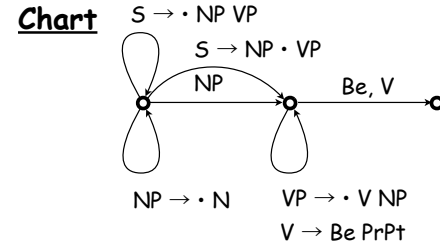
Agenda

1-2 are
 2-3 flying
 3-4 planes



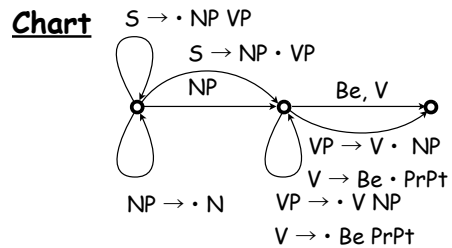
Grammar

NP → they
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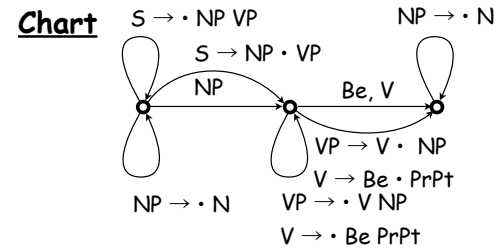
Grammar

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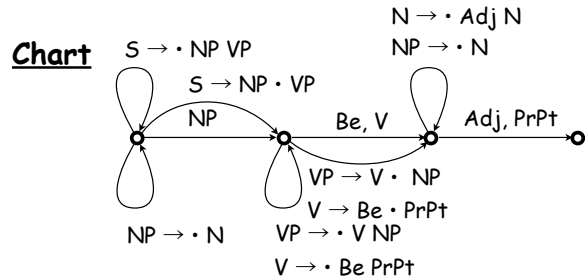


Grammar

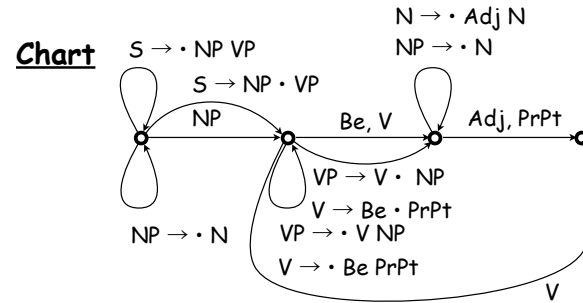
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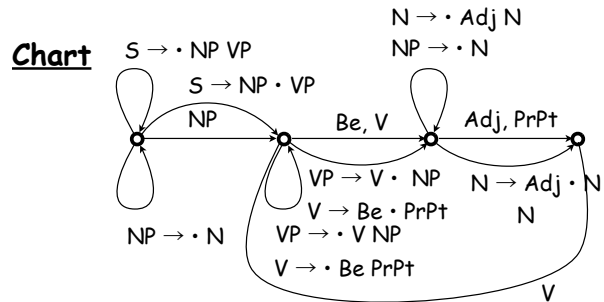
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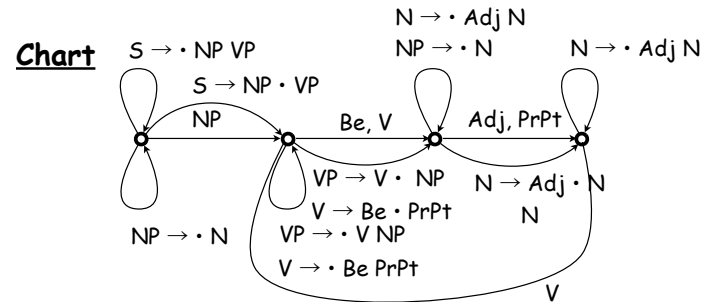
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 S → NP VP N → planes
 NP → N V → are
 N → Adj N Be → are
 VP → V NP Adj → flying
 V → Be PrPt PrPt → flying



Grammar NP → they
 S → NP VP N → planes
 NP → N V → are
 N → Adj N Be → are
 VP → V NP Adj → flying
 V → Be PrPt PrPt → flying



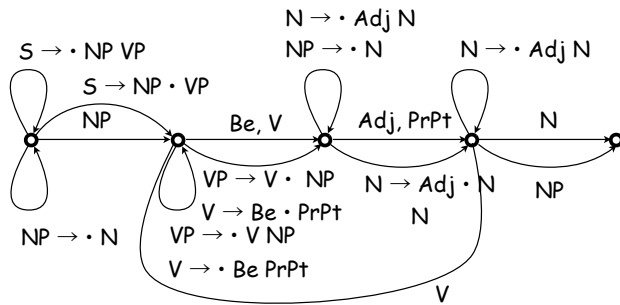
Grammar NP → they
 S → NP VP N → planes
 NP → N V → are
 N → Adj N Be → are
 VP → V NP Adj → flying
 V → Be PrPt PrPt → flying



Grammar

NP → they
 S → NP VP N → planes
 NP → N V → are
 N → Adj N Be → are
 VP → V NP Adj → flying
 V → Be PrPt PrPt → flying

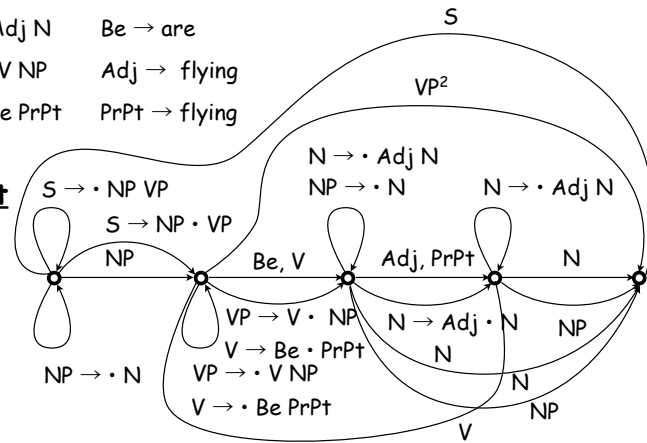
Chart



Grammar

NP → they
 S → NP VP N → planes
 NP → N V → are
 N → Adj N Be → are
 VP → V NP Adj → flying
 V → Be PrPt PrPt → flying

Chart



Predictor

predicts more than necessary

Use reachability table to
 prune predictions?

Representation of Edges

- Square array. Could use one side of diagonal for active and the other side for inactive edges—there are no inactive edges on the diagonal.
- Triangular array. For n-word sentence, store it in $n(n+1)/2$ consecutive locations.
 $index(i, j) = i(i+1)/2 + j$.

Representation of Edges

- List(s) of edges for each vertex.
- Active edges need to be accessed by their right-hand ends and inactive edges by their left hand ends. If active edges are not used, then edges need to be accessed by both ends.

The agenda

- Duplicates have to be sought in both chart and agenda.
- Indexing by vertex is useful for this search, but is otherwise not required in the agenda.
- \therefore Put edges in chart immediately, but mark them as 'pending' until removed from the agenda. Advantage: search only one indexed data structure to check for duplicates.

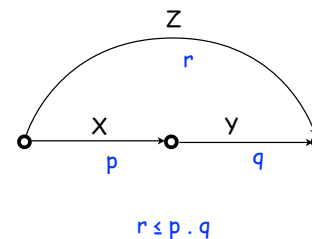
Stack Regime

Queue \Rightarrow Breadth first search

Stack \Rightarrow Depth first search

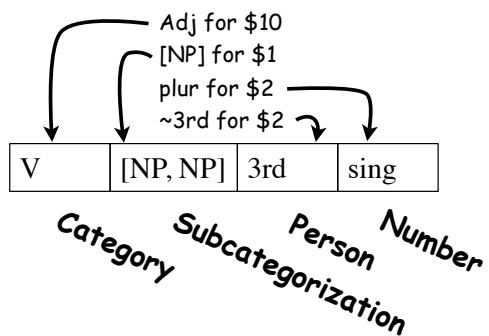
Priority sort \Rightarrow Best first

Priority invariant



Optimality

Constraints on what may be combined to form a phrase are *defeasible*. Optimality vectors:



Center-Embedding

This is the rat that ate the cheese.

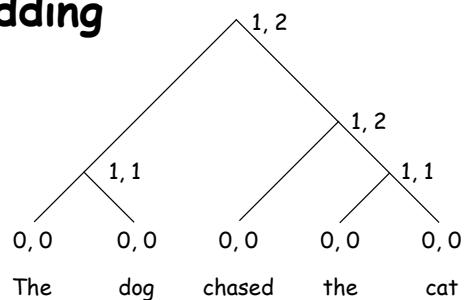
This is the cheese that the rat ate.

This is the cat that bit the rat that ate the cheese.

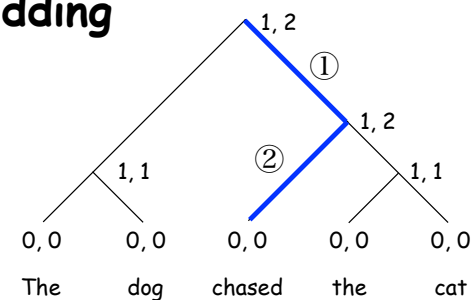
This is the cheese that the rat that the cat bit ate.

This is the cheese that [the rat that [the cat that [the dog chased] bit] ate].

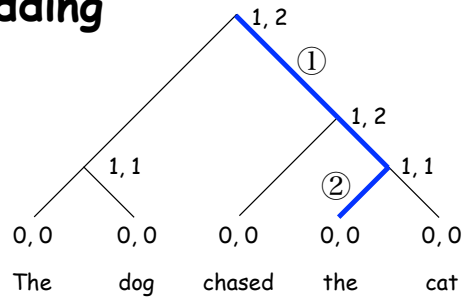
Center Embedding



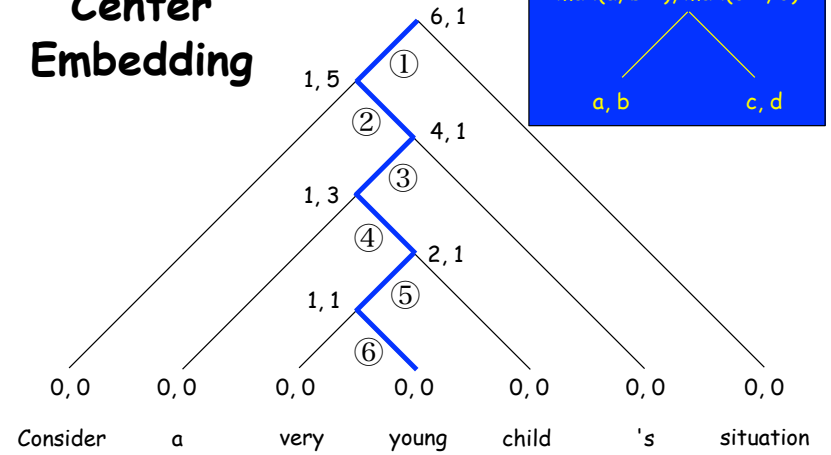
Center Embedding



Center Embedding



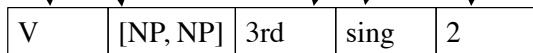
Center Embedding



Optimality

Center Embedding
 0 or 1 for \$0
 2 for \$2
 3 for \$10
 >3 for \$50

Adj for \$10
 [NP] for \$1
 plur for \$2
 ~3rd for \$2



Category
 Subcategorization
 Person
 Number

Best First

approximates breadth first because longer phrases are much less probable and probabilities of phrases of different lengths are incomparable. Likewise for most other well-behaved figures of merit.

Merging Right-Hand Sides

Grammar might have rules

$$X \rightarrow A G H P$$
$$X \rightarrow B G H P$$

Could end up with both of these in chart:

$$(2, X \rightarrow A \cdot G H P)$$
$$(2, X \rightarrow B \cdot G H P)$$

But these are now interchangeable: if one produces X then so will the other

To avoid this redundancy, can always use dotted rules of the form: $X \rightarrow \dots G H P$

Merging Right-Hand Sides

Similarly, grammar might have rules

$$X \rightarrow A G H P$$
$$X \rightarrow A G H Q$$

Could end up with both of these in chart:

$$(2, X \rightarrow A \cdot G H P) \text{ in column 5}$$
$$(2, X \rightarrow A \cdot G H Q) \text{ in column 5}$$

Not interchangeable, but we'll be processing them in parallel for a while ...

Solution: write grammar as $X \rightarrow A G H (P|Q)$

Merging Right-Hand Sides

Combining the two previous cases:

$$X \rightarrow A G H P$$
$$X \rightarrow A G H Q$$
$$X \rightarrow B G H P$$
$$X \rightarrow B G H Q$$

becomes

$$X \rightarrow (A | B) G H (P | Q)$$

And often nice to write stuff like

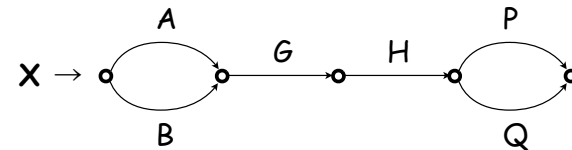
$$NP \rightarrow (\text{Det} | \varepsilon) \text{Adj}^* N$$

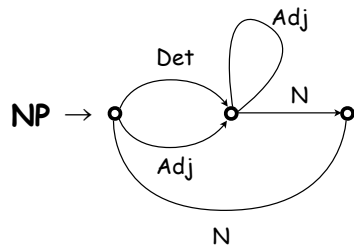
Merging Right-Hand Sides

$$X \rightarrow (A | B) G H (P | Q)$$
$$NP \rightarrow (\text{Det} | \varepsilon) \text{Adj}^* N$$

These are regular expressions!

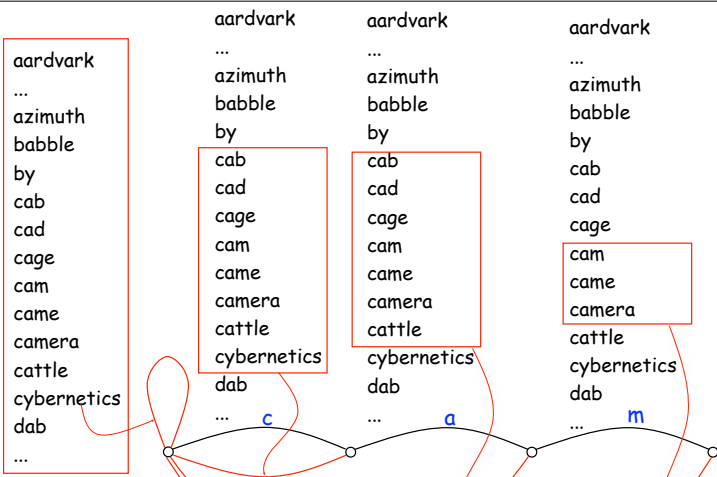
Build their minimal DFAs:





**Indeed, all NP → rules can
be combined into a single
DFA!**

- NP → ADJP ADJP JJ JJ NN NNS
 - NP → ADJP DT NN
 - NP → ADJP JJ NN
 - NP → ADJP JJ NN NNS
 - NP → ADJP JJ NNS
 - NP → ADJP NN
 - NP → ADJP NN NN
 - NP → ADJP NN NNS
 - NP → ADJP NNS
 - NP → ADJP NPR
 - NP → ADJP NPRS
 - NP → DT
 - NP → DT ADJP
 - NP → DT ADJP , JJ NN
 - NP → DT ADJP ADJP NN
 - NP → DT ADJP JJ JJ NN
 - NP → DT ADJP JJ NN
 - NP → DT ADJP JJ NN NN
- etc.**



```

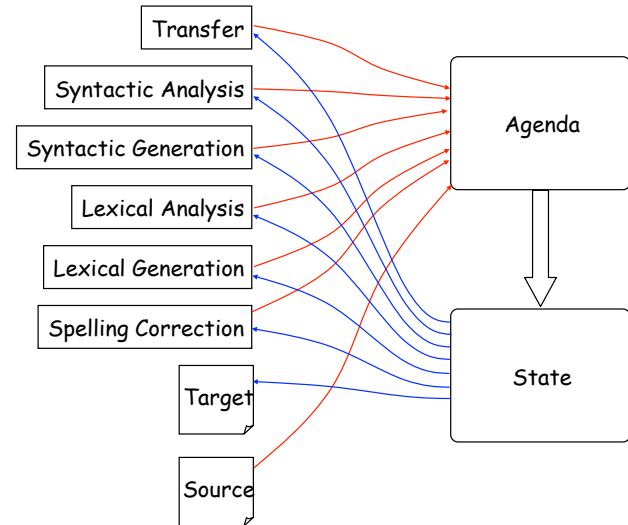
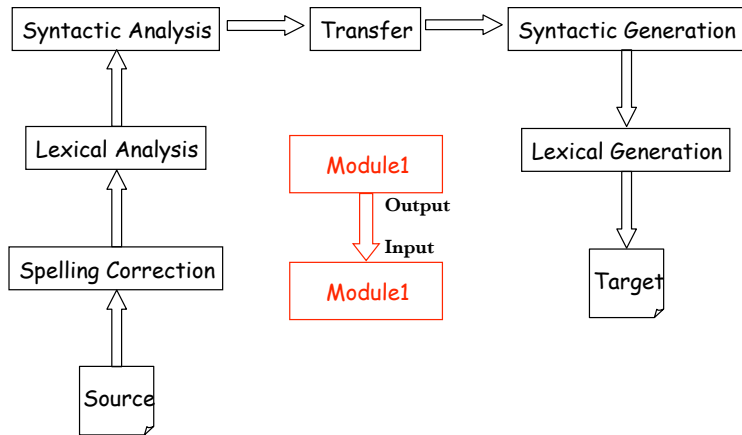
class Chart

    def initialize(agenda, monitor, filters=[])
    def reset
    def get_vertex(name)
    def schedule(task)
    def work_agenda
    def apply(active, inactive, new_pos)
    def each_vertex
    def sorted_vertices
    def show

end

```

Modular Design—Translation



```

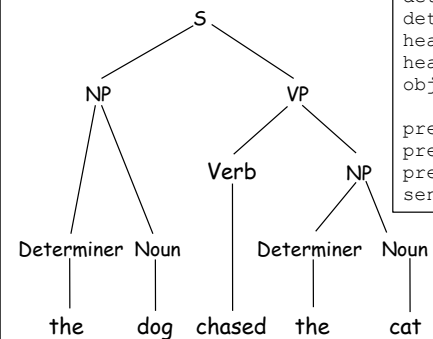
class Node
  attr_reader :name
  def initialize(name, daughters=[])
    @name = name
    @daughters = daughters
  end
end
  
```

```

det = Node.new('the')
det = Node.new('Determiner', [det])
head = Node.new('dog')
head = Node.new('Noun', [head])
subj = Node.new('NP', [det, head])

det = Node.new('the')
det = Node.new('Determiner', [det])
head = Node.new('cat')
head = Node.new('Noun', [head])
obj = Node.new('NP', [det, head])

pred = Node.new('chased')
pred = Node.new('Verb', [pred])
pred = Node.new('VP', [pred, obj])
sent = Node.new('S', [subj, pred])
  
```



```

S
NP
  Determiner
  the
  Noun
  dog
VP
  Verb
  chased
NP
  Determiner
  the
  Noun
  cat

```

```

def scan(level=0)
  yield level, self
  for d in @daughters
    d.scan(level+1) {|l, n| yield l, n}
  end
end

```

```
sent.scan {|level, node| puts "#{ ' '*level}#{@node.name}"}
```

```

class Task
  agenda = [Task.new(sent)]
  while task = agenda.shift
    task.exec {|t| agenda.unshift(t)}
  end

  def initialize(node, level=0)
    @node = node
    @level = level
  end

  def exec
    puts "#{ ' '*@level}#{@node.name}"
    @node.each_daughter{|d| yield Task.new(d, @level+1)}
  end
end

```

```

S
VP
  NP
  Noun
  dog
  Determiner
  the
  Verb
  chased
NP
  Noun
  dog
  Determiner
  the

```

```

class Task
  def initialize(node, level=[])
    @node = node
    @level = level
  end

  def exec
    puts "#{ ' '*@level.length}#{@node.name}"
    i=0
    @node.each_daughter do |d|
      yield Task.new(d, @level+[i])
      i+=1
    end
  end
end

```

```

S
NP
  Determiner
  the
  Noun
  dog
VP
  Verb
  chased
NP
  Determiner
  the
  Noun
  cat

```

```

agenda = [Task.new(sent)]
while task = agenda.shift
  task.exec do |new_task|
    (0..agenda.length).each do |i|
      agenda[i], new_task = new_task, agenda[i] \
        if (new_task.level <=> agenda[i].level) == -1
    end
    agenda.push << new_task
  end
end

```

A* (A-star)

- * begins at a selected node. Applied to this node are the "cost" of entering this node (usually zero for the initial node). A* then estimates the distance to the goal node from the current node. This estimate and the cost added together are the heuristic which is assigned to the path leading to this node. The node is then added to a priority queue, often called "open".
- The algorithm then removes the next node from the priority queue (because of the way a priority queue works, the node removed will have the lowest heuristic). If the queue is empty, there is no path from the initial node to the goal node and the algorithm stops. If the node is the goal node, A* reconstructs and outputs the successful path and stops. This path reconstruction from the stored closed nodes (see below) means it is not necessary to store the path-so-far within each node.

A* (A-star)

- If the node is not the goal node, new nodes are created for all admissible adjoining nodes; the exact way of doing this depends on the problem at hand. For each successive node, A* calculates the "cost" of entering the node and saves it with the node. This cost is calculated from the cumulative sum of costs stored with its ancestors, plus the cost of the operation which reached this new node. The algorithm also maintains a "closed" list of nodes which have been checked. If a newly generated node is already in this list with an equal or lower cost, no further processing is done on that node or with the path associated with it. If a node in the closed list matches the new one, but has been stored with a higher cost, it is removed from the closed list, and processing continues on the new node.

A* (A-star)

- Next, an estimate of the new node's distance to the goal is added to the cost to form the heuristic for that node. This is then added to the "open" priority queue, unless an identical node with lesser or equal heuristic is found there. Once the above three steps have been repeated for each new adjoining node, the original node taken from the priority queue is added to the "closed" list. The next node is then popped from the priority queue and the process is repeated.

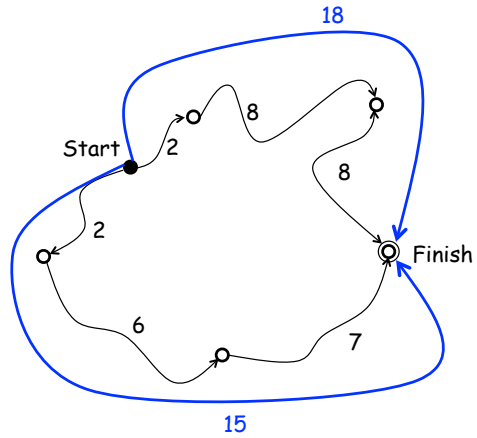
A*—Admissibility

- There is an intuitive explanation of why A* is both admissible and considers fewer nodes than any other admissible search algorithm. The intuition rests on the recognition that A* has an "optimistic" estimate of the cost of a path through every node that it considers -- optimistic in that the true cost of a path through that node to the goal will be at least as great as the estimate. But, critically, as far as A* "knows", that optimistic estimate might be achievable.
- When A* terminates its search, it by definition has found a path whose actual cost is lower than the estimated cost of any path through any open node. But since those estimates are optimistic, A* can safely ignore those nodes. In other words, A* will never overlook the possibility of a lower-cost path and so is admissible.
- Suppose now that some other search algorithm A terminates its search with a path whose actual cost is not less than the estimated cost of a path through some open node. Algorithm A cannot rule out the possibility, based on the heuristic information it has, that a path through that node might have a lower cost. So while A might consider fewer nodes than A*, it cannot be admissible. Accordingly, A* considers the fewest nodes of any admissible search algorithm that uses a no more accurate heuristic estimate.

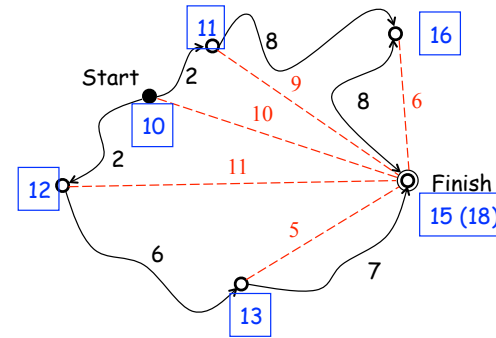
A*—Monotonicity

- If it can be guaranteed that the first path generated to any node is the shortest one, then the "closed list" is no longer necessary. Only one list of visited nodes needs to be maintained, so that nodes are not unnecessarily revisited. This guarantee can be made if the heuristic is not only admissible but monotonic, meaning that the difference between the heuristics of any two connected nodes does not overestimate the actual distance between those nodes. This is a form of the [triangle inequality](#). Not all admissible heuristics are monotonic, but straight-line distance on a map is.

A* Search



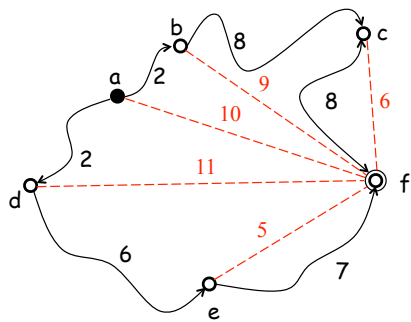
A* Search



A* Search

Open

→ a 0 10 10

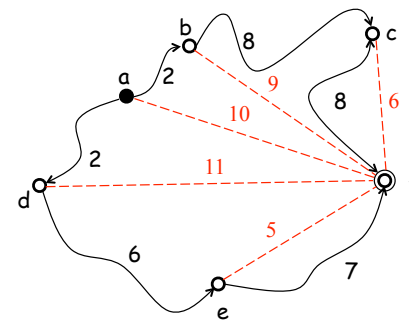


Closed

A* Search

Open

→ b 2 9 11
d 2 11 13



Closed

a 0 10 10

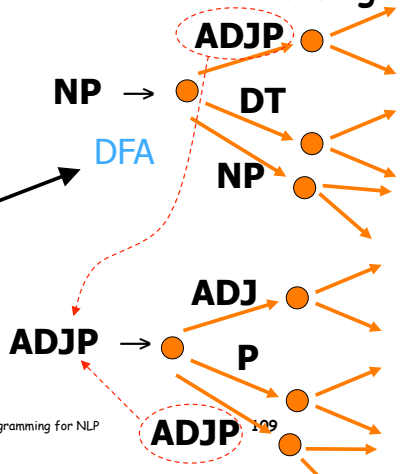
Merging Right-Hand Sides

Indeed, all NP → rules can be unioned into a single

DFA! ADJP ADJP JJ JJ NN NNS

- | ADJP DT NN
- | ADJP JJ NN
- | ADJP JJ NN NNS
- | ADJP JJ NNS
- | ADJP NN
- | ADJP NN NN
- | ADJP NN NNS
- | ADJP NNS
- | ADJP NPR
- | ADJP NPRS
- | DT
- | DT ADJP
- | DT ADJP , JJ NN
- | DT ADJP ADJP NN
- | DT ADJP JJ JJ NN
- | DT ADJP JJ NN
- | DT ADJP JJ NN NN

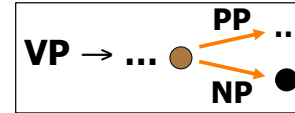
regular expression →



etc.

Earley's Algorithm on DFAs

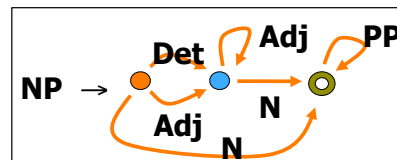
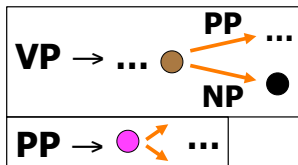
• What does Earley's algorithm now look like?



Column 4
...
(2, ●) predict

Earley's Algorithm on DFAs

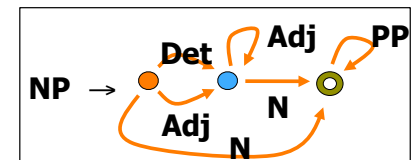
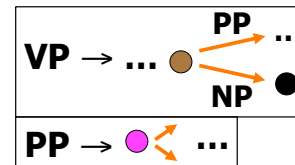
• What does Earley's algorithm now look like?



Column 4
...
(2, ●) predict
(4, ●)
(4, ○)

Earley's Algorithm on DFAs

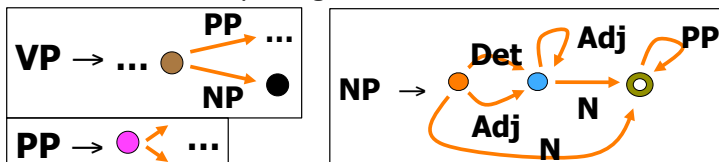
• What does Earley's algorithm now look like?



Column 4	Column 5	...	Column 7
...	...		
(2, ●)			(4, ○) predict or attach?
(4, ●)			
(4, ○) →	(4, ●)		

Earley's Algorithm on DFAs

- What does Earley's algorithm now look like?



Column 4	Column 5	...	Column 7
...	...		
(2, ●)			(4, ○) predict or attach?
(4, ●)			(7, ●) Both!
(4, ●) → (4, ●)	(4, ●) → (4, ●)		(2, ●)

Martin Kay, Algorithms and Programming for NLP 113

Preprocessing

- First "tag" the input with parts of speech:
 - Guess the correct preterminal for each word, using faster methods we'll learn later
 - Now only allow one part of speech per word
 - This eliminates a lot of crazy constituents!
 - But if you tagged wrong you could be hosed

• Raise the stakes.

Merging Right-Hand Sides

Indeed, all NP → rules can be unioned into a single

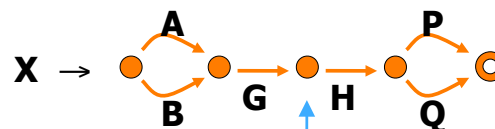
- DFA!**
- ADJP ADJP JJ JJ NN NNS
 - NP → ADJP DT NN
 - NP → ADJP JJ NN
 - NP → ADJP JJ NN NNS
 - NP → ADJP JJ NNS
 - NP → ADJP NN
 - NP → ADJP NN NN
 - NP → ADJP NN NNS
 - NP → ADJP NNS
 - NP → ADJP NPR
 - NP → ADJP NPRS
 - NP → DT
 - NP → DT ADJP
 - NP → DT ADJP, JJ NN
 - NP → DT ADJP ADJP NN
 - NP → DT ADJP JJ JJ NN
 - NP → DT ADJP JJ NN
 - NP → DT ADJP JJ NN NN

Merging Right-Hand Sides

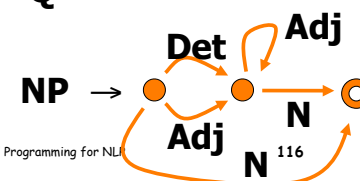
$X \rightarrow (A \mid B) G H (P \mid Q)$

$NP \rightarrow (Det \mid \epsilon) Adj^* N$

- These are regular expressions!
- Build their minimal DFAs:



- Automaton states replace dotted rules ($X \rightarrow A G \cdot H P$)



Merging Right-Hand Sides

- Combining the two previous cases:

$X \rightarrow A G H P$

$X \rightarrow A G H Q$

$X \rightarrow B G H P$

$X \rightarrow B G H Q$

becomes

$X \rightarrow (A | B) G H (P | Q)$

- And often nice to write stuff like

$NP \rightarrow (Det | \epsilon) Adj^* N$

Merging Right-Hand Sides

- Similarly, grammar might have rules

$X \rightarrow A G H P$

$X \rightarrow A G H Q$

- Could end up with both of these in chart:

(2, $X \rightarrow A . G H P$) in column 5

(2, $X \rightarrow A . G H Q$) in column 5

- Not interchangeable, but we'll be processing them in parallel for a while ...

- Solution: write grammar as $X \rightarrow A G H (P|Q)$

Merging Right-Hand Sides

- Grammar might have rules

$X \rightarrow A G H P$

$X \rightarrow B G H P$

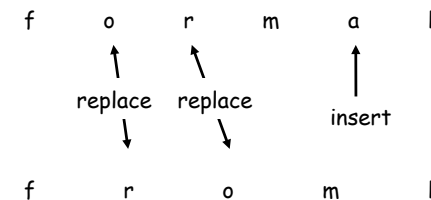
- Could end up with both of these in chart:

(2, $X \rightarrow A . G H P$) in column 5

(2, $X \rightarrow B . G H P$) in column 5

- But these are now interchangeable: if one produces X then so will the other

- To avoid this redundancy, can always use dotted rules of this form: $X \rightarrow \dots G H P$



A* Edit Distance

to goal

from start

	f	r	o	m	l		
f	1	0	2	1			
o	0	1	1	0	2	1	
r	1	2	0	2	1	1	2
m		1	3	0	2	1	2
a					0	3	1
l						0	3

Make initial chart and agenda.

Repeat until agenda is empty:

Take first arc from agenda.

Add arc to chart. (Only do this if edge is not already on the chart!)

Use the fundamental rule to combine this arc with arcs from the chart. Any edges obtained in this way should be added to the agenda.

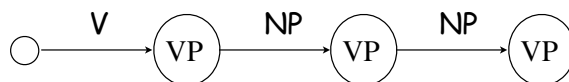
Make hypotheses (i.e., active edges) about new constituents

Representation of Rules

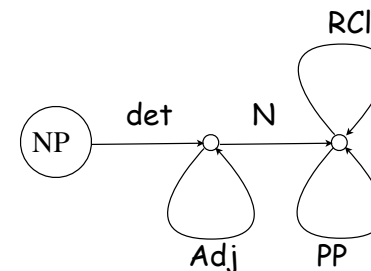
$VP \rightarrow V NP NP$

$VP \rightarrow V NP$

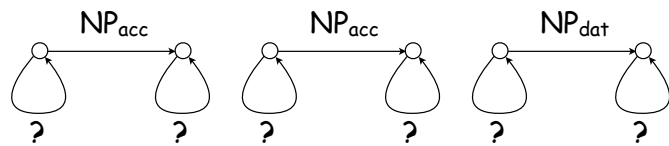
$VP \rightarrow V$



Representation of Rules



Representation of Rules



NP_{-} V_{fin} NP_{-} NP_{-} NP_{-} V_{nonfin}