FLST: Semantics II

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Semantics: The Logical Paradigm

- □ Validation of semantic representations via truthconditional interpretation
- Semantically controlled inference through entailment and deduction
- □A rigid model of compositionality



Deduction: A Question Answering Example



- Question: Which element is Thallium said to look like?
- □ Support passage: Thallium is a metallic element that resembles lead.
- Answer: Lead



Watson Again

We show that "lead" is a correct answer by deriving the representation of the question instantiated with "lead" (the "conclusion" or "hypothesis") from the representation of the answer passage (the "premiss").

Given:

metallic(thallium) ^ element(thallium) ^ resemble(thallium, lead)

U Wanted:

lement(lead) ^ look_like(thallium, lead)



More Ingredients for the Derivation

- ❑ We need some more deduction rules. These are justified by corresponding entailment relations: truth preserving transition from premises to conclusion (please, check!).
 - \Box A \land B \vdash A, A \land B \vdash B (Conjunction Elimination)
 - \Box A, B \vdash A \land B (Conjunction Introduction)
 - $\Box A, A \rightarrow B \vdash B$ (Modus Ponens)
 - \Box A \leftrightarrow B \vdash A \rightarrow B, A \leftrightarrow B \vdash B \rightarrow A (Equivalence Elimination)
 - \Box $\forall xA \vdash A[b/x]$ (Universal Instantiation)
- We need some extra bits of knowledge (axioms, taken e.g. from a lexical-semantic knowledge base):
 - □ element(lead)
 - $\Box \forall x \forall y (resemble(x,y) \leftrightarrow look_like(x,y))$



Example Derivation

- (1) metallic(th) ^ resemble(th, lead) Pre
- (2) resemble(th, lead)
- (3) $\forall x \forall y (resemble(x,y) \leftrightarrow look_like(x,y))$
- (4) $\forall y (resemble(th, y) \leftrightarrow look_like(th, y))$
- (5) resemble(th,lead) \Leftrightarrow look_like(th,lead)
- (6) resemble(th,lead) \rightarrow look_like(th,lead)
- (7) look_like(th,lead)
- (8) element(lead)
- (9) element(lead) < look_like(th, lead)

Premise 2x Conjunction Elim (1) Axiom Univ. Instantiation th/x, (3) Univ. Instantiation lead/y, (4) Equivalence Elim, (5) Modus Ponens (2), (6) Axiom Conjunction Intro (7), (8)



Word Meaning in the Logical Paradigm

Atomic predicates represent word senses, but are not very informative in themselves.

❑ Axioms express word-semantic information:
 ❑ semantic relations between different words:
 ∀x∀y(look_like(x, y) ↔ resemble(x, y))
 ❑ semantic properties of words:
 ∀x∀y(resemble(x, y) → resemble(y, x))

□ Where can we get these axioms from???

 \rightarrow Axioms can be read off lexical-semantic taxonomies like WordNet



WordNet Meaning Relations





Axioms Expressing Semantic Relations





WordNet Meaning Relations





Axioms Expressing Semantic Relations

□ B hypernym of A $\Rightarrow \forall x(A(x) \rightarrow B(x))$

 $\forall x (dolphin(x) \rightarrow toothed_whale(x)) \\ \forall x (toothed_whale(x) \rightarrow whale(x)) \\ \forall x (whale(x) \rightarrow mammal(x)) \\ \forall x (mammal(x) \rightarrow vertebrate(x)) \\ \forall x (vertebrate(x) \rightarrow animal(x)) \\ \end{cases}$

□ B hyponym of A $\Rightarrow \forall x(B(x) \rightarrow A(x))$





WordNet Meaning Relations





Axioms Expressing Semantic Relations





WordNet Meaning Relations





WordNet Meaning Relations





Axioms Expressing Semantic Relations

□ A and B cohyponyms $\Rightarrow \forall x(A(x) \rightarrow \neg B(x))$

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\forall x (mammal(x) \rightarrow \neg fish (x)) 
\forall x (fish(x) \rightarrow \neg bird (x)) 
\forall x (bird(x) \rightarrow \neg mammal(x))
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□ A and B synonyms $\Rightarrow \forall x(A(x) \leftrightarrow B(x))$



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Principle of Compositionality (Frege's Principle):

The meaning of a complex expression is uniquely determined by the meanings of its sub-expressions and its syntactic structure.



 \Box John likes Mary \Rightarrow like(john, mary)



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How do meanings of syntactic complements find their appropriate argument positions in the composition process?

 \Box The answer is: λ -Abstraction



FLST: Semantics II

Mary

λ -Abstraction

- □ student: a one-place predicate
- □ student(x): a formula containing a free variable
- \Box $\lambda x[student(x)]$: a one-place-predicate again: "to be a student"
- λx[student(x)](john): a formula: application of a one-place predicate (the λ-expression) to the individual constant "john",
- □ which is equivalent to student(john)



Interpretation of λ -expressions

$$\Box \ \llbracket \lambda x A \rrbracket^{M,g} = \{a \in U_M | \llbracket A \rrbracket^{M,g[x/a]} = 1\}$$

- $\Box \ [\lambda x[student(x)]]^{M,g}$
- = $\{a \in U_M | [student(x)]^{M,g[x/a]} = 1\}$ = $\{a \in U_M | a \in V_M (student)\}$ i.e., the set of individuals who are students,

that is V_{M} (student)

$$\begin{split} \square \ [[\lambda x[like(x, mary)]]^{M,g} &= \{a \in U_M | [[like(x, mary)]^{M,g[x/a]} = 1\} \\ &= \{a \in U_M | < a, V_M(mary) > \in V_M(like)\}, \\ &\text{ i.e., the set of individuals who like Mary.} \\ &\text{ This is not necessarily identical to the denotation of any predicate constant.} \end{split}$$



\Box John likes Mary \Rightarrow like(john, mary)





\Box John likes Mary \Rightarrow like(john, mary)





Application of λ -Expressions

 $John \Rightarrow john$

```
likes Mary \Rightarrow \lambda x[like(x, mary)]
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John likes Mary \Rightarrow \lambda x[like(x, mary)](john)
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```
⇔ like(john, mary)
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```
 \begin{split} & [\lambda x[like(x, mary)](john)]^{M,g} = 1 \\ & \text{iff } [john]^{M,g} \in [\lambda x[like(x, mary)]]^{M,g} \\ & \text{iff } V_M(john) \in \{a \in U_M | \leq a, V_M(mary) > \in V_M(like)\} \\ & \text{iff } \langle V_M(john), V_M(mary) > \in V_M(like) \end{split}
```

```
iff [[like(john, mary)]]<sup>M,g</sup> = 1
```



\Box John likes Mary \Rightarrow like(john, mary)





λ -Conversion

- □ $\lambda x[student(x)](john)$ and student(john) are equivalent, and so are $\lambda x[like(x, mary)](john)$ and like(john, mary).
- In general: \(\lambda x A(b) \⇔ A[x/b]\), where A[x/b] is the result of replacing all free occurrences of variable x in A with b. This equivalence holds independent of the choice of A and b.
- **□** Thus, we can rewrite any application of a λ -expression λ xA to an argument b by the result of substituting all free occurrences of the λ -variable x in A with b (without considering truth conditions).
- □ $\lambda xA(b) \Rightarrow A[x/b]$ as a rewrite rule is called the rule of λ -conversion or λ -reduction.



 \Box John likes Mary \Rightarrow like(john, mary)





 \Box John likes Mary \Rightarrow like(john, mary)





Semantic Composition: Lexical Information

 \Box John likes Mary \Rightarrow like(john, mary)





Semantic Composition: Projection

 \Box John likes Mary \Rightarrow like(john, mary)





Semantic Composition: Projection

 \Box John likes Mary \Rightarrow like(john, mary)





Semantic Composition: Application

 \Box John likes Mary \Rightarrow like(john, mary)





Semantic Composition: Reduction

 \Box John likes Mary \Rightarrow like(john, mary)





Semantic Composition: Application

\Box John likes Mary \Rightarrow like(john, mary)





Semantic Composition: Reduction

\Box John likes Mary \Rightarrow like(john, mary)





More λ -Expressions

- "to like Mary" λx[like(x, mary)]
- "to be liked by Mary" λx [like(mary, x)]
- "to like oneself" λx [like(x, x)]
- "to sing and dance" λx[sing(x)^dance(x)]

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"to be somebody, whom everyone likes" \lambda x[\forall y \text{ like}(y, x)]
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