### Semantic Theory Week 6 – Event semantics

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#### A problem with verbs and adjuncts

(1) The gardener killed the baron → kill<sub>1</sub>(g',b') kill<sub>1</sub> :: ⟨e,⟨e,t⟩⟩
(2) The gardener killed the baron in the park → kill<sub>2</sub>(g',b',p') kill<sub>2</sub> :: ⟨e,⟨e,t⟩⟩
(3) The gardener killed the baron at midnight → kill<sub>3</sub>(g',b',m') kill<sub>3</sub> :: ⟨e,⟨e,t⟩⟩
(4) The gardener killed the baron at midnight in the park → kill<sub>4</sub>(g',b',m',p') kill<sub>4</sub> :: ...
(5) The gardener killed the baron in the park at midnight → kill<sub>5</sub>(g',b',p',m') kill<sub>5</sub> :: ...



Q: How to explain the systematic logical entailment relations between the different uses of "kill"?

#### Davidson's solution: verbs introduce events.

Verbs expressing events have an additional event argument, which is not realised at linguistic surface:

• kill  $\mapsto \lambda y \lambda x \lambda e(kill'(e,x,y)) :: \langle e, \langle e, \langle e, t \rangle \rangle \rangle$  arity = n+1

Sentences denote sets of events:

•  $\lambda y \lambda x \lambda e(kill'(e,x,y))(b')(g') \Rightarrow^{\beta} \lambda e(kill'(e, g', b')) :: \langle e,t \rangle$ 

Existential closure turns sets of events into truth conditions

- $\lambda P \exists e(P(e)) :: \langle \langle e, t \rangle, t \rangle$
- $\lambda P \exists e(P(e))(\lambda e(kill'(e,g',b'))) \Rightarrow^{\beta} \exists e(kill'(e,g',b')) :: t$



#### Davisonian events and adjuncts

Adjuncts express two-place relations between events and the respective "circumstantial information": time, location, ...

- at midnight  $\mapsto \lambda P \lambda e(P(e) \land time(e,m')) :: \langle \langle e,t \rangle, \langle e,t \rangle \rangle$
- in the park  $\mapsto \lambda P \lambda e(P(e) \land Iocation(e,p')) :: \langle \langle e,t \rangle, \langle e,t \rangle \rangle$

The gardener killed the baron at midnight in the park

 $\mapsto \exists e (kill(e, g', b') \land time(e, m) \land location(e, p'))$   $\models \exists e (kill(e, g', b') \land time(e, m'))$   $\models \exists e (kill(e, g', b') \land location(e, p) \land time(e, m'))$   $\models \exists e (kill(e, g', b') \land location(e, p'))$   $\models \exists e (kill(e, g', b'))$ 

# Compositional derivation of event-semantic representations

the gardener killed the baron

 $\lambda x_e \lambda y_e \lambda e_e[ kill(e, y, x) ](b')(g') \Rightarrow^{\beta} \lambda e[ kill(e, g', b') ]$ 

... at midnight

 $\lambda F_{\langle e,t\rangle} \lambda e_e \left[ \begin{array}{c} F(e) \land time(e, m') \end{array} \right] (\lambda e_1 \left[ \begin{array}{c} kill(e_1, g', b') \end{array} \right] ) \Rightarrow^{\beta} \lambda e \left[ \begin{array}{c} kill(e, g, b) \land time(e, m') \end{array} \right]$ 

... in the park

$$\begin{split} \lambda F_{\langle e,t \rangle} \lambda e_e \left[ F(e) \land \text{location}(e, p') \right] \left( \lambda e_2 \left[ \text{kill}(e_2, g', b') \land \text{time}(e_2, m') \right] \right) \Rightarrow^{\beta} \\ \lambda e \left[ \text{kill}(e, g', b') \land \text{time}(e, m') \land \text{location}(e, p') \right] \end{split}$$

#### Existential closure

 $\lambda P_{\langle e,t \rangle} \exists e(P(e))(\lambda e'(K \land T \land L) \Rightarrow^{\beta} \exists e \ [ kill(e, g', b') \land time(e, m') \land location(e, p') ]$ 

#### Model structures with events

To interpret events, we need enriched ontological information

Ontology: The area of philosophy identifying and describing the basic "categories of being" and their relations.

A model structure with events is a triple  $M = \langle U, E, V \rangle$ , where

- U is a set of "standard individuals" or "objects"
- E is a set of events
- $U \cap E = \emptyset$ ,
- V is an interpretation function like in first order logic

#### Sorted (first-order) logic

A variable assignment g assigns individuals (of the correct sortspecific domain) to variables:

- $g(x) \in U$  for  $x \in VAR_U$   $VAR_U = \{x, y, z, \dots, x_1, x_2, \dots\}$  (Object variables)
- $g(e) \in E$  for  $e \in VAR_E$  VAR<sub>E</sub> = { e, e', e", ..., e<sub>1</sub>, e<sub>2</sub>, ... } (Event variables)

Quantification ranges over sort-specific domains:

- $[ \exists x \Phi ]]^{M,g} = 1$  iff there is an  $a \in U$  such that  $[ \Phi ]]^{M,g[x/a]} = 1$
- $\llbracket \exists e \Phi \rrbracket^{M,g} = 1$  iff there is an  $a \in E$  such that  $\llbracket \Phi \rrbracket^{M,g[e/a]} = 1$
- (universal quantification analogous)

#### Interpreting events

John kisses Mary → ∃e (kiss(e, j', m'))

$$\label{eq:main_set} \begin{split} & \text{[I]} \exists e \ (kiss(e,\,j',\,m')) \ ]\!]^{M,g} = 1 \\ & \textit{iff there is an } s \in E \ \text{such that} \ [I] \ kiss(e,\,j',\,m') \ ]\!]^{M,g[e/s]} = 1 \\ & \textit{iff there is an } s \in E \ \text{such that} \ \langle s, \ V_M(j'), \ V_M(m') \rangle \in V_M(kiss) \end{split}$$



#### Advantages of Davidsonian events

- ✓ Intuitive representation and semantic construction for adjuncts
- Uniform treatment of verb complements
- Uniform treatment of adjuncts and post-nominal modifiers
- Coherent treatment of tense information
- Highly compatible with analysis of semantic roles

#### Uniform treatment of verb complements

(1) Bill saw an elephant	
$\mapsto$ 3e 3x (see(e, b', x) $\land$ elephant(x))	see :: $\langle e, \langle e, \langle e, t \rangle \rangle$
(2) Bill saw an accident	
→ ∃e ∃e' (see(e, b, e') ∧ accident(e'))	see :: $\langle e, \langle e, \langle e, t \rangle \rangle$
(3) Bill saw the children play	
$\mapsto \exists e \exists e' (see(e, b, e') \land play(e', the-children))$	see :: <e,<e,<e,t>&gt;</e,<e,<e,t>

# Uniform treatment of adjuncts and post-nominal modifiers

Treatment of adjuncts as predicate modifiers, analogous to attributive adjectives:

- red  $\mapsto \lambda F \lambda x [F(x) \land red^{*}(x)]$   $\langle \langle e, t \rangle, \langle e, t \rangle \rangle$
- in the park  $\mapsto \lambda F \lambda e [F(e) \land Iocation(e, park)] \langle \langle e, t \rangle, \langle e, t \rangle \rangle$

(1) The murder in the park...

 $\mapsto \lambda F \lambda e[F(e) \land Iocation(e, park)] (\lambda e_1 [murder(e_1)])$ 

(2) The fountain in the park ....

 $\mapsto \lambda F \lambda x [F(x) \land location(x, park)] (\lambda y [fountain(y)])$ 

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## THE PAST. PRESENT & FUTURE WALK INTO A BAR IT WAS TENSE

WallingfordSign.com

#### Classical Tense Logic

- John walks walk(john)
- John walked P(walk(john))
- John will walk F(walk(john))

Syntax like in first-order logic, plus

 if Φ is a well-formed formula, then PΦ, FΦ, HΦ, GΦ are also well-formed formulae.

Φ happened in the past

 $\Phi$  has always

been the case

Φ will happen in the future

 $\Phi$  is always

going to be

the case

#### Classical Tense Logic (cont.)

Tense model structures are quadruples M =  $\langle U, T, \langle V \rangle$  where

- U is a non-empty set of individuals (the "universe")
- T is a non-empty sets of points in time
- $U \cap T = \emptyset$
- < is a linear order on T</li>
- V is a value assignment function, which assigns to every non-logical constant  $\alpha$  a function from T to appropriate denotations of  $\alpha$

 $\llbracket P\Phi \rrbracket^{M, t, g} = 1$  iff there is a t' < t such that  $\llbracket \Phi \rrbracket^{M, t', g} = 1$ 

 $\llbracket F\Phi \rrbracket^{M, t, g} = 1$  iff there is a t' > t such that  $\llbracket \Phi \rrbracket^{M, t', g} = 1$ 

#### **Temporal Relations and Events**

- (1) The door opened, and Mary entered the room.
- (2) John arrived. Then Mary left.
- (3) Mary left, before John arrived.
- (4) John arrived. Mary had left already.

Q: How to formalize temporal relations between events?

#### **Temporal Event Structure**

A model structure with events and temporal precedence is defined as M = (U, E , <, e\_u, V), where

- $U \cap E = \emptyset$ ,
- $< \subseteq E \times E$  is an asymmetric relation (temporal precedence)
- $e_u \in E$  is the utterance event
- V is an interpretation function like in standard FOL
- Overlapping events:  $e \cdot e'$  iff neither e < e' nor e' < e

### Tense in Semantic Construction

We can represent inflection as an abstract tense operator reflecting the temporal location of the reported event relative to the utterance event.

 $\mathsf{PAST} \mapsto \lambda \mathsf{P.3e} \left[ \mathsf{P}(e) \land e < e_u \right] : \langle \langle e, t \rangle, t \rangle$ 

 $\mathsf{PRES} \mapsto \lambda \mathsf{P.3e} \left[ \mathsf{P}(e) \land e \bullet e_u \right] : \langle \langle e, t \rangle, t \rangle$ 



#### Tense in Semantic Construction

Standard function application results in integration of temporal information and binding of the event variable (i.e., replacing E-CLOS):

- walk  $\mapsto \lambda x \lambda e$  [walk(e, x)]
- Bill walk  $\mapsto \lambda x \lambda e$  [walk(e, x)](b')  $\Rightarrow^{\beta} \lambda e$  [walk(e, b')]

Bill walk PAST
→ λE ∃e [E(e) ∧ e < e<sub>u</sub>](λe' [walk(e', b)])
⇒<sup>β</sup> ∃e [λe' [walk(e', b)](e) ∧ e < e<sub>u</sub>]
⇒<sup>β</sup> ∃e [walk(e, b) ∧ e < e<sub>u</sub>]



#### **Event Structure**

**Observation**: Events are generally constructs that consist of various (temporally ordered) sub-events

• E.g., "scripts": visit a restaurant or shopping in the supermarket

Idea: Induce structure into events universe



A **partial order** is a structure  $\langle A, \leq \rangle$  where  $\leq$  is a reflexive, transitive, and antisymmetric relation over A.

- The **join** of a and  $b \in A$  (Notation:  $a \sqcup b$ ) is the lowest upper bound for a and b.
- The **meet** of a and  $b \in A$  (Notation:  $a \sqcap b$ ) is the highest lower bound for a and b.

A **lattice** is a partial order  $\langle A, \leq \rangle$  that is closed under meet and join.

A join semi-lattice is a partial order  $\langle A, \leq \rangle$  that is closed under join

#### Model Structure with Sub-Events

We can change the structure of the events universe to represent sub-event relations:  $M = \langle U, \langle E, \leq_e \rangle, <, e_u, V \rangle$ , where:

- $U \cap E = \emptyset$ ,
- $< \subseteq E \times E$  is an asymmetric relation (temporal precedence)
- $e_u \in E$  is the utterance event
- $\langle E, \leq_e \rangle$  is a join semi-lattice
- V is an interpretation function

#### Model Structure with Sub-Events (cont.)

The model structure  $M = \langle U, \langle E, \leq_e \rangle$ ,  $\langle e_u, V \rangle$  must observe some additional constraints on  $\langle$  and  $\leq_e$ , for instance:

- If  $e_1 < e_2$  and  $e_1' \leq_e e_1$  and  $e_2' \leq_e e_2$ , then  $e_1' < e_2'$
- If  $e_1' \circ e_2'$  and  $e_1' \leq_e e_1$  and  $e_2' \leq_e e_2$ , then  $e_1 \circ e_2$

**Sidenote**: We could introduce a similar structuring of the universe of entities in order to capture *plurality* and other *composite entities* 

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### Verbal arguments; a related problem?

(1) John broke the window with a rock.

(2) A rock **broke** the window.

(3) The window broke.

And we're back to the same entailment issue:

```
\exists e(break_3(e, j, w, r)) \nvDash \exists e(break_2(e, r, w)) \nvDash \exists e(break_1(e, w))
```

#### Semantic/Thematic roles

 agent
 patient
 instrument

 (1) John broke the window with a rock

 ← ∃e [break(e) ∧ agent(e, j) ∧ patient(e, w) ∧ instrument(e, r)]

 (2) A rock broke the window.

 ← ∃e [break(e) ∧ patient(e, w) ∧ instrument(e, r)]

 (3) The window broke.

 ← ∃e [break(e) ∧ patient(e, w)]

In standard FOL: Thematic roles are implicitly represented by the canonical order of the arguments

In Davidsonian event semantics: Thematic roles are two-place relations between the event denoted by the verb, and an argument role filler.

#### Interpretation of events with thematic roles

John kisses Mary  $\mapsto \exists e (kiss(e) \land agent(e, j') \land patient(e, m'))$ 

$$\label{eq:main_set} \begin{split} & [\![ \exists e \ (kiss(e) \land agent(e, j') \land patient(e, m')) \ ]\!]^{M,g} = 1 \\ & \textit{iff there is an } s \in E \ \text{such that} \ [\![kiss(e)]\!]^{M,g[e/s]} = 1 \ \text{and} \ [\![ agent(e, j')]\!]^{M,g[e/s]} = 1 \\ & \text{and} \ [\![ patient(e, m')]\!]^{M,g[e/s]} = 1 \\ & \textit{iff there is an } s \in E \ \text{such that} \ s \in V_M(kiss) \ \text{and} \ \langle s, V_M(j') \rangle \in V_M(agent) \\ & \text{and} \ \langle s, V_M(m') \rangle \in V_M(patient) \end{split}$$



#### Thematic roles & verbal differences/similarities

Different verbs allow different thematic role configurations

- (1) a. John broke the window with a rock → agent, patient, instrument
   b. John smiled at Mary → agent, recipient
- (2) a. The window **broke** allows inanimate subject
  - *b.* \*The bread **cut** ----- does not allow inanimate subject

Thematic roles capture equivalences and entailment relations between different predicates

(3) a. Mary gave Peter the book
b. Peter received the book from Mary
∀e[give(e) ↔ receive(e)] ⊨ (3a) ↔ (3b)

#### Determining the role inventory

Fillmore (1968): "thematic roles form a small, closed, and universally applicable inventory conceptual argument types."

A typical role inventory might consist of the roles:

 Agent, Patient, Theme, Recipient, Instrument, Source, Goal, Beneficiary, Experiencer.

But... there are some difficult cases:

(1) Lufthansa is replacing its 737s with Airbus 320

(2) John sold the car to Bill for 3,000€

(3) Bill bought the car from John for 3,000€

#### Semantic corpora with thematic roles

- PropBank: includes a separate role inventory for every lemma
- FrameNet: "Frame-based" role inventories

Frames are structured schemata representing complex prototypical situations, events, and actions

(1) [Agent Lufthansa] is replacing Frame: REPLACING [Old its 737s] [New With Airbus A320s]

(2) [Agent Lufthansa] is substitutingFrame: REPLACING [New Airbus A320s] [Old for its 737s]

#### Semantic corpora with thematic roles (cont.)

**PropBank** (Palmer et al. 2005): Annotation of Penn TreeBank with predicate-argument structure.

- (1) [Arg0 Lufthansa] is replacing [Arg1 its 737s] [Arg2 with Airbus A320s]
- (2) [Arg0 Lufthansa] is substituting [Arg1 Airbus A320s] [Arg2 for its 737s]

**FrameNet** (Baker et al. 1998): A database of frames and a lexicon with frame information

- (3) [Agent Lufthansa] is replacingFrame: REPLACING [Old its 737s] [New with Airbus A320s]
- (4) [Agent Lufthansa] is substitutingFrame: REPLACING [New Airbus A320s] [Old for its 737s]

Pred	replace
ArgO	Lufthansa
Argl	its737s
Arg2	AirbusA320s
Prea	substitute
Arg0	substitute Lufthansa
Arg0 Arg1	substitute Lufthansa AirbusA320s

REPLACING
Lufthansa
its737s
AirbusA320s

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... but how does it combine with other semantic constructs?

### A problem with events and quantification

John kissed Mary

 $\mapsto \lambda P.P(j') [ \lambda P.P(m')(\lambda y \lambda x \lambda e [kiss(e) \land agent(e,x) \land patient(e,y)]) ]$ 

 $\Rightarrow^{\beta} \lambda e [kiss(e) \land agent(e,j') \land patient(e,m')]$ 

 $\Rightarrow^{\text{E-CLOS}} \exists e \text{ [kiss(e) } \land \text{ agent(e,j') } \land \text{ patient(e,m')]}$ 

John kissed every girl

 $\mapsto \lambda P.P(j') [\lambda P. \forall x(girl'(x) \rightarrow P(x))(\lambda y \lambda x \lambda e [kiss(e) \land agent(e, x) \land patient(e, y)])]$ 

 $\Rightarrow^{\beta} \lambda e [\forall x(girl'(x) \rightarrow kiss(e) \land agent(e,j') \land patient(e,x)]$ 

 $\Rightarrow^{E-CLOS} \exists e [\forall x(girl'(x) \rightarrow kiss(e) \land agent(e,j') \land patient(e,x)]$ 

### Two solutions to the event quantification problem

#### Solution I

Interpret sentences as generalized quantifiers over events:  $\langle \langle e,t \rangle, t \rangle$  instead of  $\langle e,t \rangle$  (E-CLOS part of lexical semantics) (Champollion, 2010; 2015)

```
kiss \mapsto \lambda F_{(v,t)}. \exists e \text{ (kiss(e) } \land F(e)\text{)} :: \langle \langle v,t \rangle, t \rangle \approx \{ F \mid F \cap K | SS \neq \emptyset \}
separate type for events!
```

Introduce separate types for regular NPs and quantified NPs, and restrict existential closure to regular NPs (Winter & Zwarts, 2011; de Groote & Winter, 2014)

john  $\mapsto$  j :: e every girl  $\mapsto \lambda Q. \forall x(girl(x) \rightarrow Q(x)) :: \langle \langle e, t \rangle, t \rangle$ kiss  $\mapsto \lambda x \lambda y \lambda e. kiss(e, x, y) :: \langle e, \langle e, \langle v, t \rangle \rangle \rangle$ e-clos  $\mapsto \lambda P. \exists e(P(e)) :: \langle \langle v, t \rangle, t \rangle$ separate type for events!

#### Solution I: Sentences as GQs over events



(Champollion, 2010; 2015)

#### Solution II: Type-restriction for existential closure

$$\frac{\vdash \text{EVERY} : N \to (NP \to S) \to S \qquad \vdash \text{GIRL} : N}{\vdash \text{EVERY GIRL} : (NP \to S) \to S}$$
(1)

$$\frac{\vdash \text{KISSED} : NP \to NP \to V \qquad x : NP \vdash x : NP}{x : NP \vdash \text{KISSED} \ x : NP \to V \qquad \qquad \vdash \text{JOHN} : NP}_{x : NP \vdash \text{KISSED} \ x JOHN} (2)$$

$$\frac{\vdash \text{E-CLOS} : V \to S \qquad x : NP \vdash \text{KISSED } x \text{ JOHN} : V}{x : NP \vdash \text{E-CLOS} (\text{KISSED } x \text{ JOHN}) : S} \xrightarrow[(3)]{} \\
\frac{\downarrow \text{A}x. \text{E-CLOS} (\text{KISSED } x \text{ JOHN}) : NP \to S}{(3)}$$

$$\begin{array}{c} \vdots & (1) \\ \vdash \text{ every Girl} : (NP \to S) \to S \\ \vdash \lambda x. \text{ e-clos} (\text{kissed } x \text{ john}) : NP \to S \\ \vdash \text{ every Girl} (\lambda x. \text{ e-clos} (\text{kissed } x \text{ john})) : S \end{array}$$

(Winter & Zwarts, 2011; de Groote & Winter, 2014)

#### Links

- Overview paper: Lasersohn (2012) Event-Based Semantics: <u>https://semanticsarchive.net/Archive/jFhNWM2M/</u> <u>eventbasedsemantics.pdf</u>
- PropBank: <u>http://propbank.github.io/</u>
- FrameNet: <a href="https://framenet.icsi.berkeley.edu/fndrupal/">https://framenet.icsi.berkeley.edu/fndrupal/</a>