

Semantic Theory

Lecture 6 – Discourse Semantics

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(Slides by Manfred Pinkal)

A simple context theory

- Some natural-language expressions vary their meaning with context:
 - *I, you, here, this, now, ...*
- Model contexts as vectors: sequences of semantically relevant context data with fixed arity.
- Model meanings as functions from contexts to denotations – more specifically, as functions from specific context components to denotations.

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An Example

- Context $c = \langle a, b, l, t, r \rangle$
 - *a* speaker
 - *b* addressee
 - *l* utterance location
 - *t* utterance time
 - *r* referred object

$\llbracket I \rrbracket^{M,g,c} = \text{utt}(c) = a$
 $\llbracket you \rrbracket^{M,g,c} = \text{adr}(c) = b$
 $\llbracket here \rrbracket^{M,g,c} = \text{loc}(c) = l$
 $\llbracket now \rrbracket^{M,g,c} = \text{time}(c) = t$
 $\llbracket this \rrbracket^{M,g,c} = \text{ref}(c) = r$

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Type-theoretic context semantics

- **Model structure:** $M = \langle U, C, V \rangle$
 - U - model universe
 - C - context set
 - V - value assignment function that assigns non-logical constants functions from contexts to denotations of appropriate type.
- **Interpretation:**
 - $\llbracket \alpha \rrbracket^{M,h,c} = V(\alpha)(c)$, if α is a non-logical constant
 - $\llbracket \alpha \rrbracket^{M,h,c} = h(\alpha)$, if α is a variable
 - $\llbracket \alpha(\beta) \rrbracket^{M,h,c} = \llbracket \alpha \rrbracket^{M,h,c}(\llbracket \beta \rrbracket^{M,h,c})$
 - etc.

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An example

- $I \text{ am reading this book} \Rightarrow \text{read}'(\text{this-book}')(I')$
- $\llbracket \text{read}'(\text{this-book}')(I') \rrbracket^{M,h,c} = 1$
 - iff $\llbracket \text{read}' \rrbracket^{M,h,c}(\llbracket \text{this-book}' \rrbracket^{M,h,c})(\llbracket I' \rrbracket^{M,h,c}) = 1$
 - iff $V(\text{read}')(\text{ref}(c))(\text{utt}(c)) = 1$
- Context-invariant expressions are constant functions:
 - $V(\text{read}')(\text{c}) = V(\text{read}')(\text{c}')$ for all $c, c' \in C$

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Two basic classes of context dependent expressions

- **Deictic expressions** depend on the physical utterance situation:
 - *I, you, now, here, this, ...*
- **Anaphoric expressions** refer to the linguistic context / previous discourse:
 - *he, she, it, then, ...*
- But there is more ...

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More context-dependent expressions

Semantic context dependence is a pervasive property of natural language:

- (1) **Every student** must be familiar with the basic properties of first-order logic
- (2) It is hot and sunny **everywhere**
- (3) John **always** is late
- (4) Bill has bought an **expensive** car
- (5) **Another one**, please!
- (6) **The student** is working

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Definite NPs: Type-theoretic analysis

Standard type-theoretic representation of definite NPs:

the $\Rightarrow \lambda F \lambda G \exists y [\forall x (F(x) \leftrightarrow x=y) \wedge G(y)]$

the sun $\Rightarrow \lambda G \exists y [\forall x [\text{sun}'(x) \leftrightarrow x=y] \wedge G(y)]$

the sun is shining $\Rightarrow \exists y [\forall x [\text{sun}'(x) \leftrightarrow x=y] \wedge \text{shine}'(y)]$

the student is working $\Rightarrow \exists y [\forall x [\text{stud}'(x) \leftrightarrow x=y] \wedge \text{work}'(y)]$

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Context-dependence of definite NPs

- Definite NPs pick an appropriate object from context.
 - *The student* is working
 - $\exists y [\forall x [\text{student}'(x) \leftrightarrow x=y] \wedge \text{work}'(y)]$ (??)
- Utterances typically contain several noun phrases referring to different objects:
 - **The student** is reading **the book** in **the library**
- Noun phrases may refer to different objects of the same type, in one utterance situation:
 - *the book*
 - *the blue book*
 - *the blue book about discourse semantics*

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Indefinite Noun Phrases

Standard type-theoretic analysis of *a student is working*

$$a \Rightarrow \lambda P \lambda Q \exists x [P(x) \wedge Q(x)]$$

$$\text{a student} \Rightarrow \lambda Q \exists x [\text{student}'(x) \wedge Q(x)]$$

$$\text{a student is working} \Rightarrow \exists x [\text{student}'(x) \wedge \text{work}'(x)]$$

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Indefinite Noun Phrases

- *A student is working*
 - $\Rightarrow \exists x [\text{student}'(x) \wedge \text{work}'(x)]$
- Indefinite noun phrases establish the context for later reference, they introduce new reference objects:
 - **A student is working. She is successful.**
- Type-theoretic semantics cannot model this effect:
 - $\text{she} \Rightarrow \lambda P.P(x)$
 - *She is successful* $\Rightarrow \text{successful}'(x)$
 - *A student is working. She is successful.*
 $\Rightarrow \exists x [\text{student}'(x) \wedge \text{work}'(x)] \wedge \text{successful}'(x)$

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The interaction of definite and indefinite noun phrases

- Natural-language meaning and context interact in two ways:
 - Context determines the utterance meaning.
 - The semantics of the utterance changes the context.
- The „context change potential“ is part of the meaning of natural-language expressions.

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Context dependence: Wrap up [1/3]

- The interpretation of most context-dependent expressions, e.g., definite noun phrases, is determined by context in a complex way.

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Context dependence: Wrap up [2/3]

- Some types of expressions, like indefinite noun phrases, introduce new context information, which is available at a later stage of discourse for anaphoric reference.
- Modelling this kind of context change potential is outside the reach of standard type-theoretic semantics, with or without context-semantic extension.

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Context dependence: Wrap up [3/3]

Reference objects established in discourse need not be real objects:

- (1) *Once upon a time there was a king, who had a beautiful daughter.*
- (2) *Someone - whoever that may be - will eventually find out. That person will tell others, and everyone will be terribly upset.*

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Discourse Semantics

Discourse Semantics

- The basic idea:
 - Meaning = **Context Change Potential**
- Focus on anaphoric use of noun phrases:
 - definite and indefinite, full noun phrases and pronouns.
- Meaning representation uses **discourse referents** in addition to formulas encoding truth conditions.
- “Division of labor” between definite and indefinite NPs:
 - Indefinite NPs introduce new discourse referents
 - Definite NPs refer to “old” or “familiar” discourse referents (which are already part of the meaning representation)

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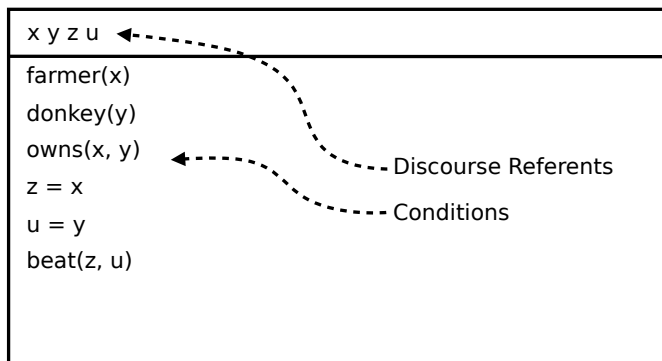
Discourse Semantics

- Discourse Representation Theory: Hans Kamp (1981), Irene Heim (1980)
- Reading: Hans Kamp and Uwe Reyle: From Discourse to Logic, Kluwer: Dordrecht 1993.

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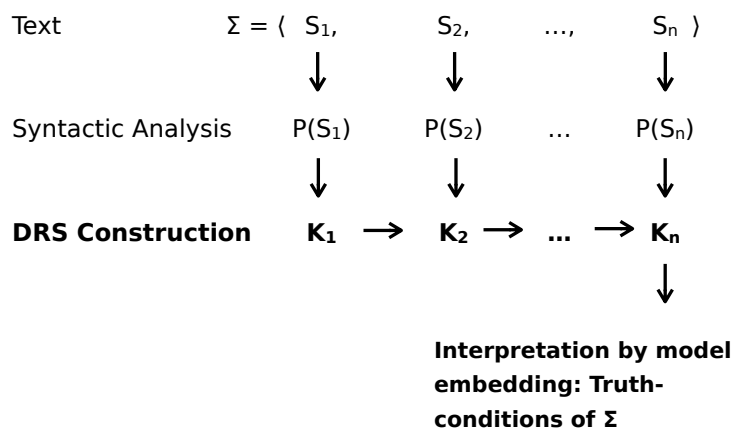
A simple example

- A farmer owns a donkey. He beats it.



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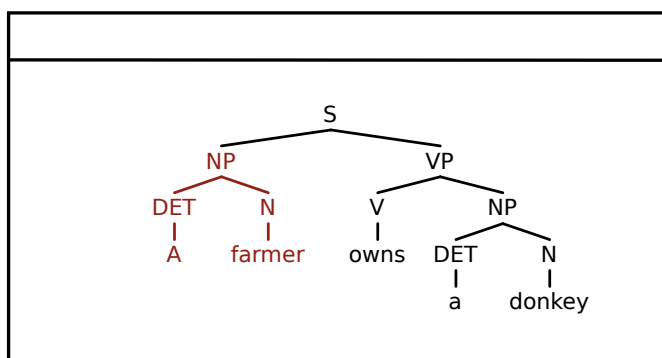
Discourse Representation Theory



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An Example

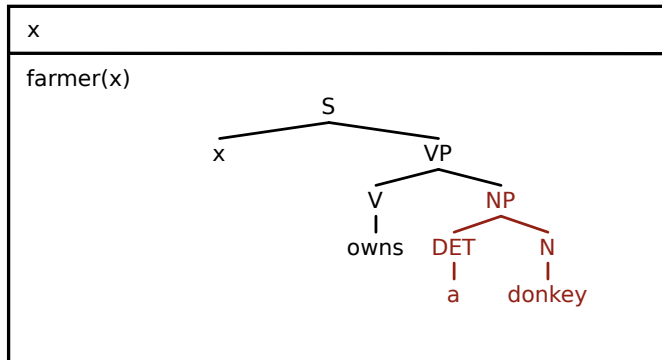
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An Example

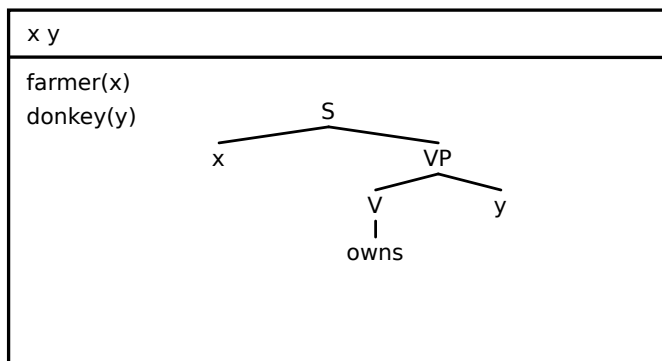
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An Example

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An Example

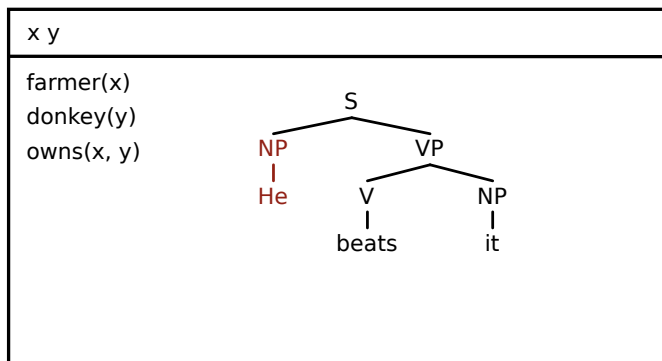
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An Example

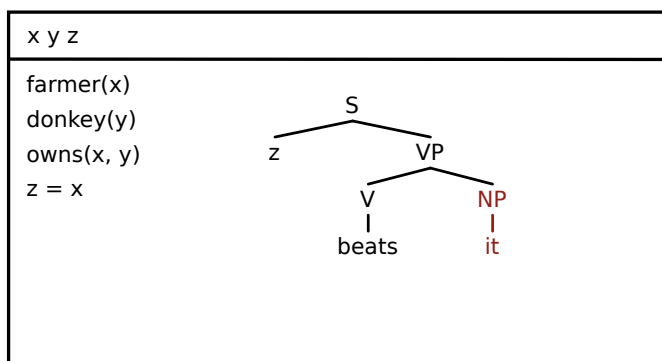
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An Example

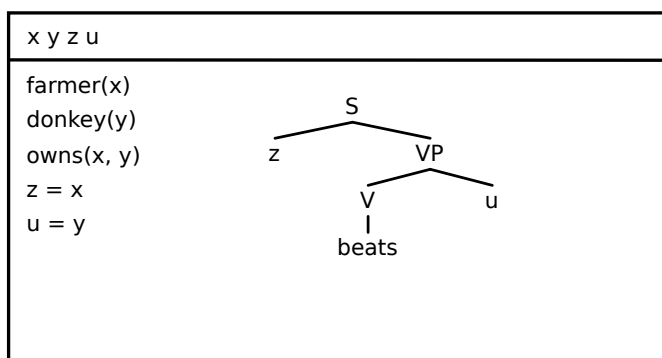
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An Example

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An Example

- A farmer owns a donkey. He beats it.

x y z u
farmer(x) donkey(y) owns(x, y) z = x u = y beat(z, u)

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DRS (Basic Version)

- **A discourse representation structure (DRS) K** is a pair $\langle U_K, C_K \rangle$, where
 - U_K is a set of discourse referents
 - C_K is a set of (reduced or reducible) conditions
- **Reduced conditions:**
 - $R(u_1, \dots, u_n)$ R an n -place relation, $u_i \in U_K$
 - $u = v$ $u, v \in U_K$
 - $u = a$ $u \in U_K, a$ is proper name
- **Reducible conditions:**
 - Conditions of form α or $\alpha(x_1, \dots, x_n)$, where α is a context-free parse tree.

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DRS (Basic Version)

- A discourse referent (DR) **u is free in DRS $K = \langle U_K, C_K \rangle$**
 - if u is free in one of K 's conditions,
 - and $u \notin U_K$.
- A DRS **K is closed** iff no DR occurs free in K .
- A reducible (fully reduced) DRS is a DRS which contains (does not contain) reducible conditions.

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DRS Construction Algorithm

- Input:
 - a text $\Sigma = \langle S_1, \dots, S_n \rangle$
 - a DRS $K_0 (= \langle \emptyset, \emptyset \rangle)$, by default
- Repeat for $i = 1, \dots, n$:
 - Add parse tree $P(S_i)$ to the conditions of K_{i-1} .
 - Apply DRS construction rules to reducible conditions of K_{i-1} , until no reduction steps are possible any more.
 - The resulting DRS is K_i , the discourse representation of text $\langle S_1, \dots, S_i \rangle$.

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Construction Rule for Indefinite NPs

- **Triggering Configuration:**
 - α is reducible condition in DRS K , containing $[s [NP \beta] [VP \gamma]]$ or $[VP [v \gamma] [NP \beta]]$ as a substructure.
 - β is $\varepsilon\delta$, ε indefinite article
- **Action:**
 - Add a new DR x to U_K .
 - Replace β in α by x .
 - Add $\delta(x)$ to C_K .

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Construction Rule for Personal Pronouns

- **Triggering Configuration:**
 - α is reducible condition in DRS K , containing $[s [NP \beta] [VP \gamma]]$ or $[VP [v \gamma] [NP \beta]]$ as substructure.
 - β is a personal pronoun.
- **Action:**
 - Add a new DR x to U_K .
 - Replace β in α by x .
 - Select an appropriate DR $y \in U_K$, and add $x = y$ to C_K .

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Construction Rule for Proper Names

- **Triggering Configuration:**

- α is reducible condition in DRS K , containing $[s [NP \beta] [VP \gamma]]$ or $[VP [v \gamma] [NP \beta]]$ as substructure.
- β is a proper name.

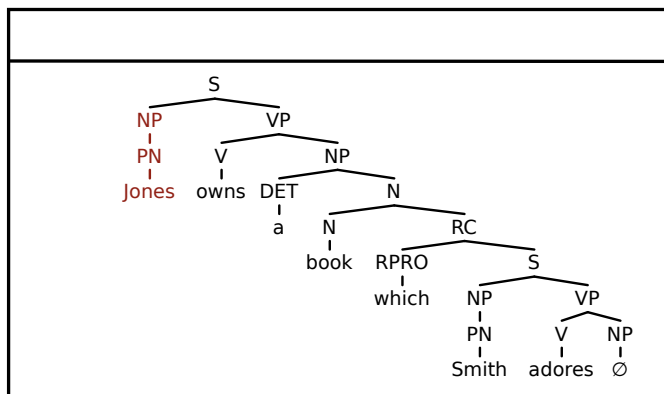
- **Action:**

- Add a new DR x to U_k .
- Replace β in α by x .
- Add $x = \beta$ to C_k .

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Relative Clauses

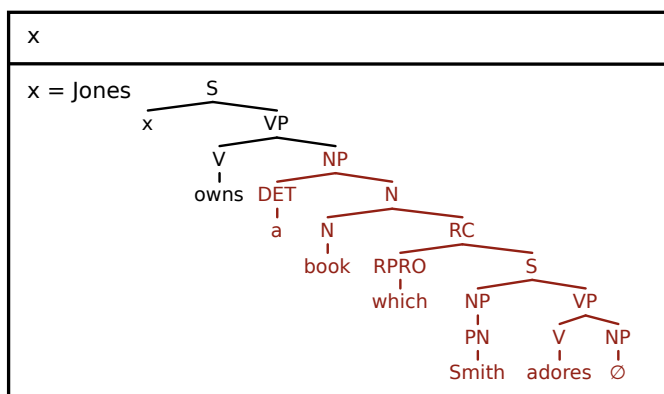
- Jones owns a book which Smith adores.



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Relative Clauses

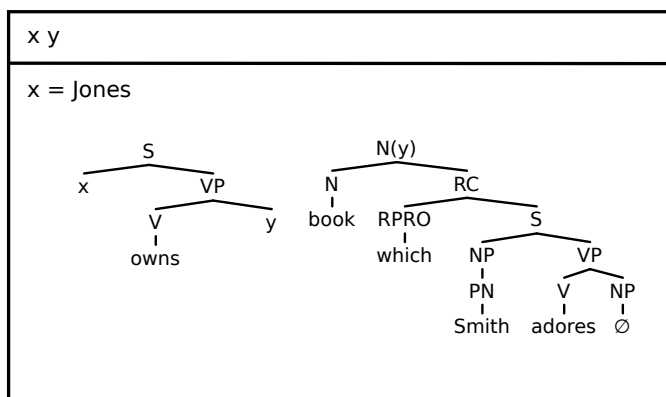
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Relative Clauses

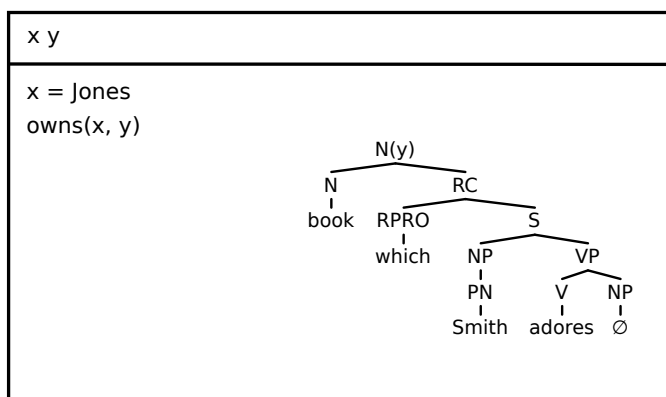
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Relative Clauses

- Jones owns a book which Smith adores.



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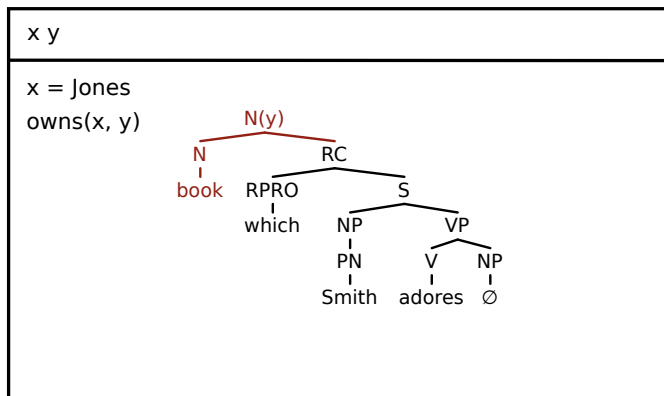
Construction Rule for Relative Clauses

- **Triggering configuration:**
 - $\alpha(x)$ is reducible condition in DRS K, containing $[N' [N' \beta] [RC \gamma]]$ as a substructure
 - γ is relative clause of the form $\delta \epsilon$, where δ is a relative pronoun and ϵ a sentence with an NP gap, δ and t are co-indexed.
- **Actions:**
 - Remove $\alpha(x)$ from CK.
 - Add $\beta(x)$ to CK .
 - Replace the NP gap in ϵ by x , and add the resulting structure to CK.

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Relative Clauses

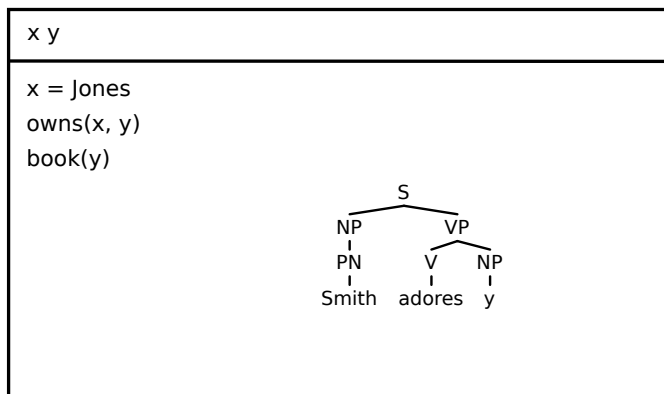
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Relative Clauses

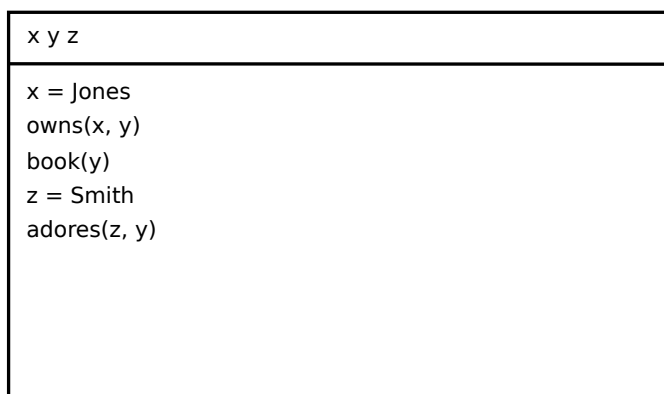
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Relative Clauses

- Jones owns a book which Smith adores.



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A constraint on the DRS construction algorithm

A problem: The basic DRS construction algorithm can derive DRSES for both of the following sentences, with the indicated anaphoric binding:

(1) *[A professor]_i recommends a book that she_i likes*

(2) **She_i recommends a book that [a professor]_i likes*

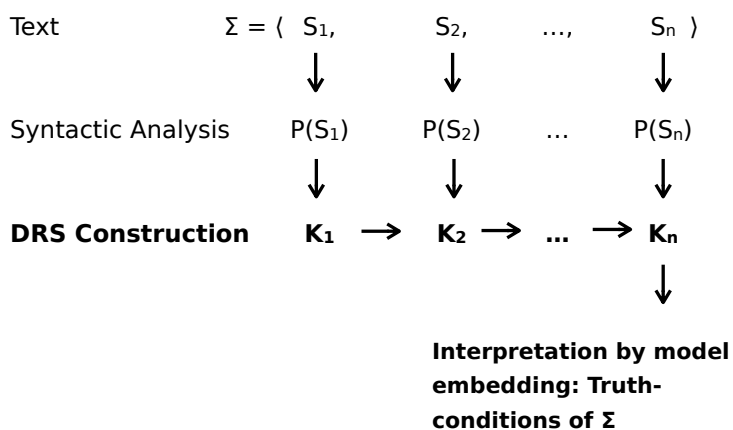
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The Highest Triggering Configuration Constraint

- If two triggering configurations of one or two different DRS construction rules occur in a reducible condition, then first apply the construction rule to the highest triggering configuration.
- The highest triggering configuration is the one whose top node dominates the top nodes of all other triggering configurations.

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Discourse Representation Theory



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Denotational Interpretation

- Let
 - U_D a set of discourse referents,
 - $K = \langle U_K, C_K \rangle$ a DRS with $U_K \subseteq U_D$,
 - $M = \langle U_M, V_M \rangle$ a FOL model structure appropriate for K , i.e. a model structure that provides interpretations for all relation symbols occurring in K
- An embedding of K into M is a (partial) function f from U_D to U_M such that $U_K \subseteq \text{Dom}(f)$.

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Verifying embedding

- **An embedding** f of K in M **verifies K in M** iff f verifies every condition $\alpha \in C_K$
 - Notation: $f \models_M K$
- **f verifies condition α in M** ($f \models_M \alpha$):
 - $f \models_M R(x_1, \dots, x_n)$ iff $\langle f(x_1), \dots, f(x_n) \rangle \in V_M(R)$
 - $f \models_M x = a$ iff $f(x) = V_M(a)$
 - $f \models_M x = y$ iff $f(x) = f(y)$

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Example Computation

- [BEISPIEL]

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Truth

- Let K be a closed DRS and M be an appropriate model structure for K .
- K is true in M iff there is a verifying embedding f of K in M such that $\text{Dom}(f) = U_K$

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Basic features of DRT

- DRT models linguistic meaning as anaphoric potential (through DRS construction) plus truth conditions (through model embedding).
- In particular, DRT explains the ambivalent character of indefinite noun phrases:

Expressions that introduce new reference objects into the context, and are truth conditionally equivalent to existential quantifiers.

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Translation of DRSs to FOL

- A DRS $K = \langle \{x_1, \dots, x_n\}, \{c_1, \dots, c_k\} \rangle$

$x_1 \dots x_n$
c_1
\vdots
c_n

is truth-conditionally equivalent to the following FOL formula:

- $\exists x_1 \dots \exists x_n [c_1 \wedge \dots \wedge c_k]$

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Literature

- Hans Kamp (1981). A Theory of Truth and Semantic Representation.
- Hans Kamp & Uwe Reyle (1993). From Discourse to Logic. Introduction to Modeltheoretic Semantics of Natural Language, Formal Logic and Discourse Representation Theory. Kluwer, Dordrecht.