

Semantic Theory: Scope

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Sentence Semantics

- Step 1: FOL Representations
- Step 2: Types and Higher-order Logic
 - Higher-order expressions (higher-order predicates, adjectives, degree modifiers)
 - Function application as basic operation for semantics construction
 - Unified, compositional semantics for noun phrases
- Step 3: λ -expressions and β -reduction
 - Higher-order expressions for semantic composition
 - Obtaining FOL sentence representations through β -reduction
 - Semantics Construction with Transitive Verbs
- Step 4: Treatment of scope variation
 - Cooper Storage
 - Underspecification



Semantics Construction: Basic rules

- Rule of functional application:

$$\begin{array}{c} A \\ / \quad \backslash \\ B \quad C \end{array} \quad \frac{B \Rightarrow \beta: \langle \sigma, \tau \rangle \quad C \Rightarrow \gamma: \sigma}{A \Rightarrow \beta(\gamma): \tau} \quad \text{or} \quad \frac{B \Rightarrow \beta: \sigma \quad C \Rightarrow \gamma: \langle \sigma, \tau \rangle}{A \Rightarrow \gamma(\beta): \tau}$$

- Rule of non-branching nodes:

$$\begin{array}{c} A \\ | \\ B \end{array} \quad \frac{B \Rightarrow \beta: \tau}{A \Rightarrow \beta: \tau}$$



Semantics Construction: Basic rules

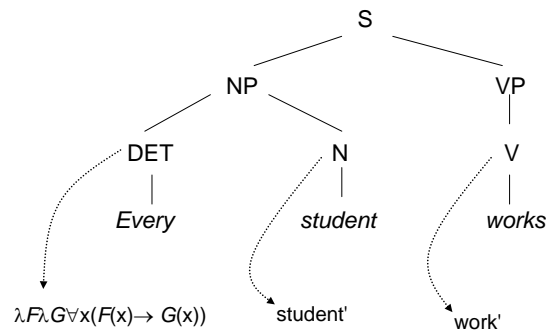
- Rule of lexical nodes:

$$\begin{array}{c} A \\ | \\ a \end{array} \quad \frac{}{A \Rightarrow \beta: \tau}$$

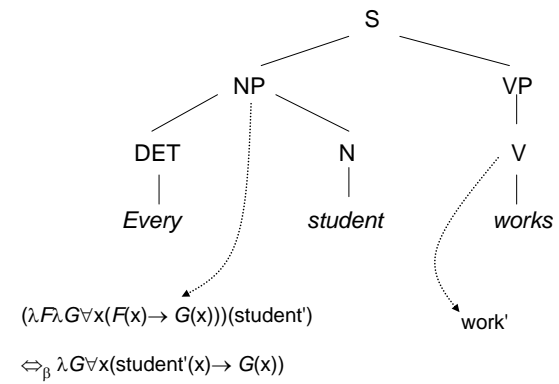
The semantic representation β for the word "a" is supplied by the lexicon.



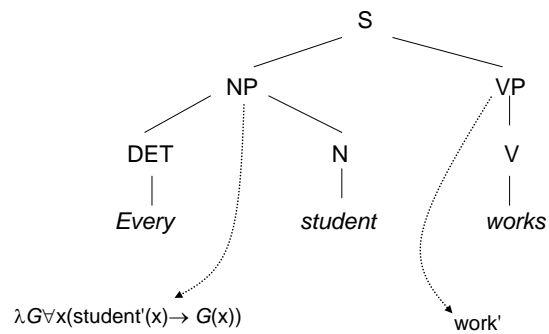
An example



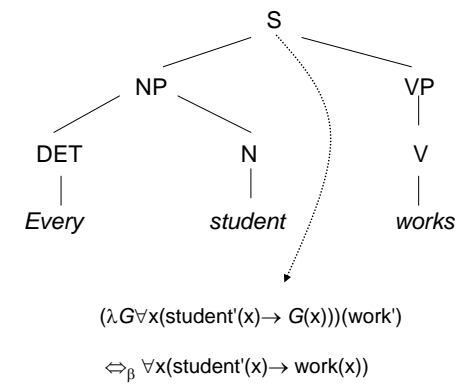
An example



An example



An example





Scope: Terminology

- Logic: **Quantifier** & **Scope**

$\forall x(\text{student}'(x) \rightarrow \text{work}(x))$

- NL Semantics
 - **Determiner**+**Restriction** form NP-Denotation („Generalized Quantifier“)
 - NP Denotation is applied to its **Nuclear Scope**

Every'(**student'**)(**work'**)



A Note on Notation

- Either: Use expanded notation from the beginning (e.g., $\lambda G \lambda G \forall x(F(x) \rightarrow G(x))$), and simplify (i.e., beta-reduce) as early as possible
- Or: Use abbreviations (every'), and expand them later:
 - Every'(student')(work')
 - $\lambda G \lambda G \forall x(F(x) \rightarrow G(x))(\text{student}')(\text{work}')$
- Or: Combine both in a sensible way
- But: Don't rewrite expanded forms, whenever you can avoid it



Variable NP Scope

- Every linguist speaks two languages*
- Our company has an expert for every problem*
- A search engine for every subject*



NPs and scope-sensitive operators

- Every student **didn't** pay attention*
- Every citizen **can** become president*
- During his visit to China, Helmut Kohl **intends** to visit a factory for CFC-free refrigerators*



The problem of scope variation

- The scope of noun phrases is not determined by the syntactic position in which they occur.
- Divergence between syntactic and semantic structure is a challenge for compositionality and semantics constructions.
- Scope variation may lead to a proliferation of readings



Scope ambiguity

Every student presents a paper.

(a) $\forall x[\textit{student}(x) \rightarrow \exists y[\textit{paper}(y) \wedge \textit{present}(x,y)]]]$

(b) $\exists y[\textit{paper}(y) \wedge \forall x[\textit{student}(x) \rightarrow \textit{present}(x,y)]]]$

Every student didn't pay attention.

(a) $\forall x[\textit{student}(x) \rightarrow \neg \textit{pay-attention}(x)]]$

(b) $\neg \forall x[\textit{student}(x) \rightarrow \textit{pay-attention}(x)]]$



Scope ambiguity

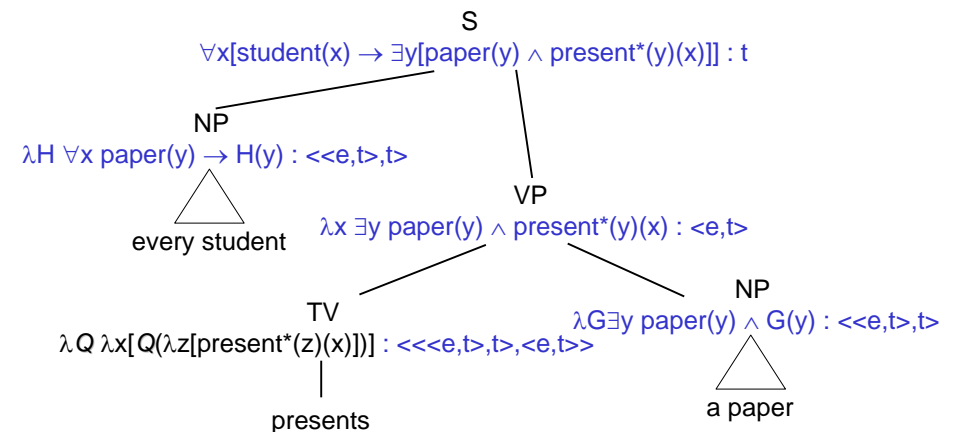
- Every researcher of a company saw some sample.

1. $\forall x(\textit{res}'(x) \wedge \exists y(\textit{cp}'(y) \wedge \textit{of}'(x,y)) \rightarrow \exists z(\textit{spl}'(z) \wedge \textit{see}'(x,z))$
2. $\exists z(\textit{spl}'(z) \wedge \forall x(\textit{res}'(x) \wedge \exists y(\textit{cp}'(y) \wedge \textit{of}'(x,y)) \rightarrow \textit{see}'(x,z))$
3. $\exists y(\textit{cp}'(y) \wedge \forall x(\textit{res}'(x) \wedge \textit{of}'(x,y)) \rightarrow \exists z(\textit{spl}'(z) \wedge \textit{see}'(x,z))$
4. $\exists y(\textit{cp}'(y) \wedge \exists z(\textit{spl}'(z) \wedge \forall x(\textit{res}'(x) \wedge \textit{of}'(x,y)) \rightarrow \textit{see}'(x,z))$
5. $\exists z(\textit{spl}'(z) \wedge \exists y(\textit{cp}'(y) \wedge \forall x(\textit{res}'(x) \wedge \textit{of}'(x,y)) \rightarrow \textit{see}'(x,z))$

Every researcher of a company saw some samples of most products.



So far, we get only one reading



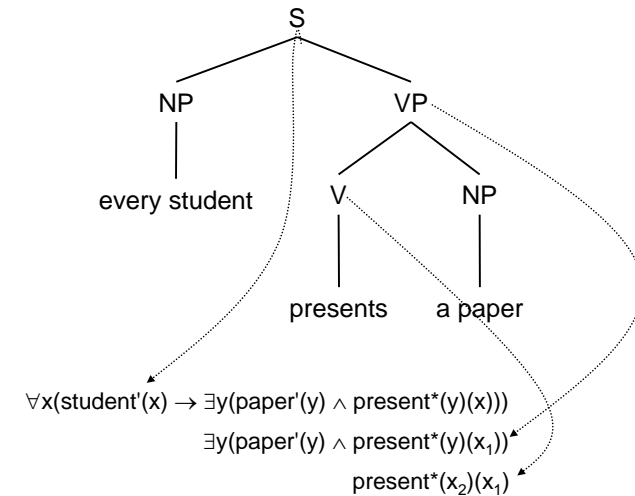


The problem with scope

- Sentences with scope ambiguities can have **multiple** semantic representations for a syntactic constituent.
- The order of the scope-bearing elements (quantifiers, negation, adverbs, ...) don't necessarily follow the order of the syntactic combination.
- But: With the approach we have so far, we can only derive a **single** semantic representation for each constituent.
- How can we solve this problem?



Example



The missing reading

- We get one reading of the sentence by deriving the following terms:

$$\begin{aligned} \forall x(\text{student}'(x) \rightarrow \exists y(\text{paper}'(y) \wedge \text{present}^*(y)(x))) \\ \exists y(\text{paper}'(y) \wedge \text{present}^*(y)(x_1)) \\ \text{present}^*(x_2)(x_1) \end{aligned}$$

- We should be able to construct the second reading correspondingly:

$$\begin{aligned} \exists y(\text{paper}'(y) \wedge \forall x(\text{student}'(x) \rightarrow \text{present}^*(y)(x))) \\ \forall x(\text{student}'(x) \rightarrow \text{present}^*(x_2)(x)) \\ \text{present}^*(x_2)(x_1) \end{aligned}$$



Solving the scope problem: Principles

- We can obtain the second reading by delaying the application of the inner noun phrase.
- To this purpose, we have to:
 - temporarily store the noun phrase denotation away
 - formally bind the object argument position by a variable
 - make sure that the correct argument position will be bound, when the „real“ noun phrase denotation is eventually applied



Using lambda abstraction („Quantifying-in“)

- Abstract over the correct variable and then apply the NP representation to the abstracted term.

$$\lambda F \forall x (\text{student}'(x) \rightarrow F(x)) (\lambda x_1. \lambda G \exists y (\text{paper}'(y) \wedge G(y)) (\lambda x_2. \text{present}^*(x_2)(x_1)))$$

$$\lambda G \exists y (\text{paper}'(y) \wedge G(y)) (\lambda x_2. \text{present}^*(x_2)(x_1))$$

$$\lambda G \exists y (\text{paper}'(y) \wedge G(y)) (\lambda x_2. \lambda F \forall x (\text{student}'(x) \rightarrow F(x)) (\lambda x_1. \text{present}^*(x_2)(x_1)))$$

$$\lambda F \forall x (\text{student}'(x) \rightarrow F(x)) (\lambda x_1. \text{present}^*(x_2)(x_1))$$

- Problem: How can we do this compositionally?



Nested Cooper Storage

- One algorithm for deriving such representations compositionally is Nested Cooper Storage (Keller 1988). It repairs some problems of the original Cooper Storage (Cooper 1975).
- Cooper Storage technique is used to compute the set of all semantic readings nondeterministically from a single syntactic analysis.



Nested Cooper Storage: Principles

- The semantic values of syntactic constituents are ordered pairs $\langle \alpha, \Delta \rangle$:
 - $\alpha \in WE_\tau$ is the **content**
 - Δ is the **quantifier store**: a set of NP representations that must still be applied.
- At NP nodes, we may **store** the content in Δ .
- At sentence nodes, we can **retrieve** NP representations from the store in arbitrary order and apply them to the appropriate argument positions.



Nested Cooper Storage: Storage

$$B \Rightarrow \langle \gamma, \Gamma \rangle \quad \text{B is an NP node}$$

$$B \Rightarrow \langle \lambda P.P(x_i), \{\langle \gamma, \Gamma \rangle_i\} \rangle \quad \text{where } i \in \mathbf{N} \text{ is a new index}$$

- Using this rule, we can assign more than one semantic value to an NP node.
- The content of the new semantic value is just a placeholder of type $\langle \langle e, t \rangle, t \rangle$, and the old value (including its store) is moved to the store.



Nested Cooper Storage: Old Rules Adjusted

- Rule of functional application:

$$\begin{array}{c} A \\ / \quad \backslash \\ B \quad C \end{array} \quad \frac{B \Rightarrow \langle \beta, \Delta \rangle \quad C \Rightarrow \langle \gamma, \Gamma \rangle}{A \Rightarrow \langle \beta(\gamma), \Delta \cup \Gamma \rangle} \quad \text{or} \quad \frac{B \Rightarrow \langle \beta, \Delta \rangle \quad C \Rightarrow \langle \gamma, \Gamma \rangle}{A \Rightarrow \langle \gamma(\beta), \Delta \cup \Gamma \rangle}$$

- Rule of non-branching nodes:

$$\begin{array}{c} A \\ | \\ B \end{array} \quad \frac{B \Rightarrow \langle \beta, \Delta \rangle}{A \Rightarrow \langle \beta, \Delta \rangle}$$

- Rule of lexical nodes:

$$\begin{array}{c} A \\ | \\ a \end{array} \quad \frac{}{A \Rightarrow \langle \beta, \emptyset \rangle}$$



Nested Cooper Storage: Principles

- A syntactic constituent may be associated with multiple semantic values of this form.
- A lambda term M counts as a semantic representation for the entire sentence iff we can derive $\langle M, \emptyset \rangle$ as a value for the root of the syntax tree.
- Hence, there may be more than one valid semantic representation for the complete sentence.



Nested Cooper Storage: Retrieval

$$A \Rightarrow \langle \alpha, \Delta \cup \{ \langle \gamma, \Gamma \rangle_i \} \rangle \quad \text{A is any sentence node}$$

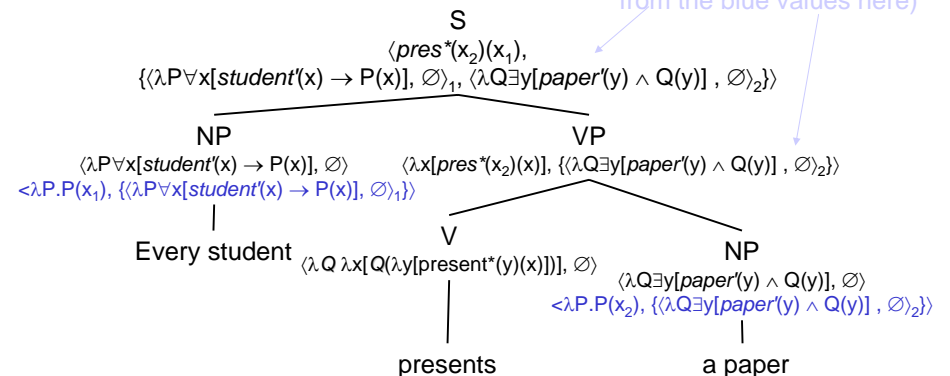
$$\frac{}{A \Rightarrow \langle \gamma(\lambda x_i \alpha), \Delta \cup \Gamma \rangle}$$

- Using this rule, we can apply a stored NP.
- At this point, the correct λ -abstraction for the variable associated with the stored element is introduced.
- The old store Γ is released into the store for A.



Nested Cooper Storage: Example

Every student presents a paper. only showing the results from the blue values here





Retrieval: Reading 1

- By applying the Retrieval rule, we can derive the following representation for the S node:

$$\begin{aligned} & \langle pres^*(x_2)(x_1), \{ \langle \lambda P \forall x [student(x) \rightarrow P(x)], \emptyset \rangle_1, \\ & \langle \lambda Q \exists y [paper(y) \wedge Q(y)], \emptyset \rangle_2 \} \rangle \Rightarrow_R \langle \lambda Q \exists y [paper(y) \wedge \\ & Q(y)] (\lambda x_2. pres^*(x_2)(x_1)), \\ & \{ \langle \lambda P \forall x [student(x) \rightarrow P(x)], \emptyset \rangle_1 \} \rangle \\ & \Rightarrow_\beta \langle \exists y [paper(y) \wedge pres^*(y)(x_1)], \{ \langle \lambda P \forall x [student(x) \rightarrow \\ & P(x)], \emptyset \rangle_1 \} \rangle \\ & \Rightarrow_R \langle \lambda P \forall x [student(x) \rightarrow P(x)] (\lambda x_1. \exists y [paper(y) \wedge \\ & pres^*(y)(x_1)]), \emptyset \rangle \\ & \Rightarrow_\beta \langle \forall x [student(x) \rightarrow \exists y [paper(y) \wedge pres^*(y)(x)]], \emptyset \rangle \end{aligned}$$



Retrieval: Reading 2

$$\begin{aligned} & \langle pres^*(x_2)(x_1), \{ \langle \lambda P \forall x [student(x) \rightarrow P(x)], \emptyset \rangle_1, \\ & \langle \lambda Q \exists y [paper(y) \wedge Q(y)], \emptyset \rangle_2 \} \rangle \Rightarrow_R \langle \lambda P \forall x [student(x) \rightarrow \\ & P(x)] (\lambda x_1. pres^*(x_2)(x_1)), \\ & \{ \langle \lambda Q \exists y [paper(y) \wedge Q(y)], \emptyset \rangle_2 \} \rangle \\ & \Rightarrow_\beta \langle \forall x [student(x) \rightarrow pres^*(x_2)(x)], \{ \langle \lambda Q \exists y [paper(y) \wedge \\ & Q(y)], \emptyset \rangle_2 \} \rangle \\ & \Rightarrow_R \langle \lambda Q \exists y [paper(y) \wedge Q(y)] (\lambda x_2. \forall x [student(x) \rightarrow \\ & pres^*(x_2)(x)]), \emptyset \rangle \\ & \Rightarrow_\beta \langle \exists y [paper(y) \wedge \forall x [student(x) \rightarrow pres^*(y)(x)]], \emptyset \rangle \end{aligned}$$



Compositionality

- The Compositionality Principle as stated earlier:
The meaning of a complex expression is uniquely determined by the meaning of its subexpressions and its syntactic structure.
- Nested Cooper Storage shows: We can maintain this principle even in the face of semantic (scope) ambiguity, if we use a relaxed concept of „meaning“.



Compositionality and NCS

- Two versions of the Compositionality Principle:
 - on the level of denotations
 - on the level of semantic representations
- Nested Cooper Storage is clearly compositional on the level of semantic representations - but in a less straightforward way than last week's construction algorithm.
- Compositional on the level of denotations: only in a very indirect sense.



Scope islands

- Nested Cooper Storage makes the simplifying assumption that NPs can be retrieved at all sentence nodes.
- This is not true in general because sentence-embedding verbs create **scope islands**:
 - John said that he saw every girl. (1 reading)
- Quantifiers may not be lifted across the S node of the embedded clause; the sentence cannot mean "for every girl x, John said that he saw x".



Scope ambiguities in real-world texts

- Some large-scale grammars (e.g. the English Resource Grammar) compute semantic representations with scope.
- The ERG analyses all NPs as scope bearers to keep the grammar simple. (This is not necessarily correct: proper names, definites, etc.)
- Median number of scope readings in the Rondane corpus: 55.
(But: The median number of semantic equivalence classes is only 3!)



Conclusion

- Last week's type-driven semantics construction is a nice first step.
- But it is fundamentally unable to deal with semantically ambiguous sentences.
- Scope ambiguity: Application order of NP representations can be different from syntactic structure.
- Nested Cooper Storage: Equip semantic representations with a quantifier store to allow flexible application of quantifiers; multiple semantic representations per syntactic constituents allowed.