

Semantic Theory

Lecture 4: Further topics in sentence semantics

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Semantics of λ -expressions

- If $\alpha \in WE_\tau$, $v \in Var_\sigma$, then $[[\lambda v \alpha]]^{M,g}$ is that function $f : D_\sigma \rightarrow D_\tau$ such that for all $a \in D_\sigma$, $f(a) = [[\alpha]]^{M,g[v/a]}$
- Notice that of course $f \in D_{\langle \sigma, \tau \rangle}$.
- In general: $[[(\lambda v \alpha) (\beta)]]^{M,g} = [[\alpha]]^{M,g[v/[[\beta]]^{M,g}]}$

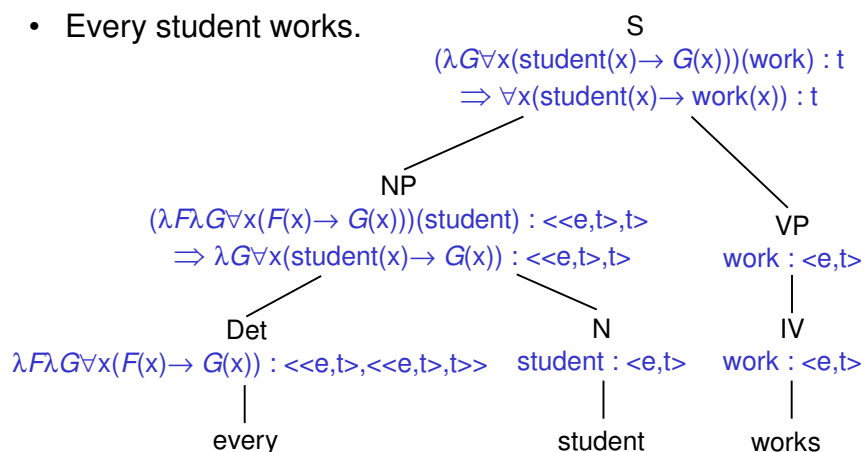
Conversion rules in the λ -calculus

- β -conversion:
 $\lambda v \alpha(\beta) \Leftrightarrow [\beta/v] \alpha$, if all free variables in β are free for v in α .
- α -conversion:
 $\lambda v \alpha \Leftrightarrow \lambda v' [\beta'/v] \alpha$, if v' is free for v in α .
- η -conversion:
 $\lambda v \alpha(v) \Leftrightarrow \alpha$

The rule which we will use most in semantics construction is β -conversion in the left-to-right direction (β -reduction), which allows us to simplify representations.

Semantics construction

- Every student works.



This week: Conservative extensions

- Semantics construction for further constructions:
 - adjectives
 - transitive verbs
- Extensions to the logic:
 - intensionality
 - tense and modality

The problem with adjective semantics

John is a blond criminal

$\text{criminal}(j) \wedge \text{blond}(j)$

John is an honest criminal

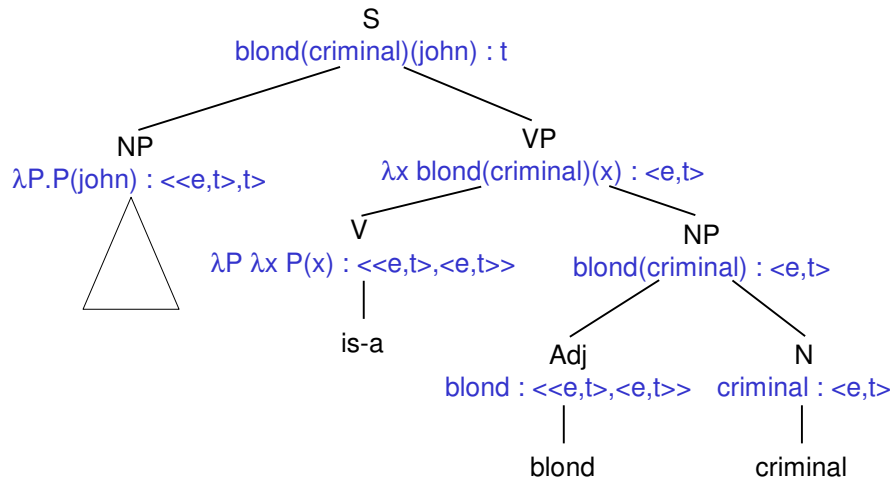
$\text{criminal}(j) \wedge \text{honest}(j)$?

John is an alleged criminal

$\text{criminal}(j) \wedge \text{alleged}(j)$??

... but this one is
actually ok!

Another problem with adjective semantics



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Adjective classes

- Adjectives can be classified with respect to the way that they combine the adjective and noun meanings:
 - intersective adjectives (blond, carnivorous, ...):
 $|| \text{blond N} || = || \text{blond} || \cap || \text{N} ||$
 - subsective adjectives (skillful, typical, ...):
 $|| \text{skillful N} || \subseteq || \text{N} ||$
 - privative adjectives (past, fake, ...):
 $|| \text{past N} || \cap || \text{N} || = \emptyset$
 - there are also other non-subsective adjectives that are not privative (alleged, ...)

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A new problem with adjectives

- We want the best of both worlds:
 - compositional semantics construction
 - explicit and meaningful final semantic representations
- We don't have this yet for intersective adjectives.
- We can get this in two different ways:
 - use meaning postulates
 - use more explicit lambda terms

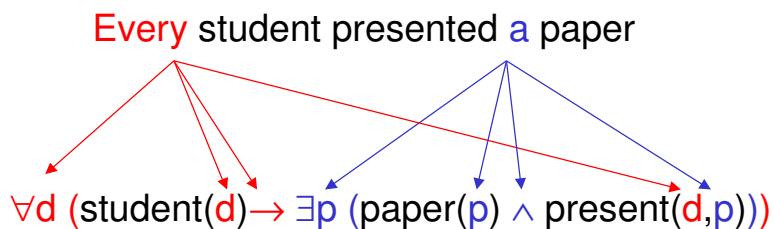
Meaning postulates

- Characterise the meaning of a word by using logical axioms.
- Meaning postulate for intersective adjectives:
 - $\exists P \forall Q \forall x \text{ blond}(Q)(x) \leftrightarrow P(x) \wedge Q(x)$
- These axioms would be part of our background knowledge.
- For example, we could infer "criminal(john)" from "blond(criminal)(john)" and this axiom.
- More generally applicable for other words.

More explicit lambda terms

- In the special case of intersective adjectives, we can also do it by assigning the word a more elaborate lambda term:
 - $\lambda P \lambda x (P(x) \wedge \text{blond}^*(x))$
- This will beta-reduce to the formula we want.
- Note that the symbol "blond*" has type $\langle e, t \rangle$ here (and should denote the set of blond individuals in the universe), but the entire semantic representation of the word "blond" has type $\langle \langle e, t \rangle, \langle e, t \rangle \rangle$.

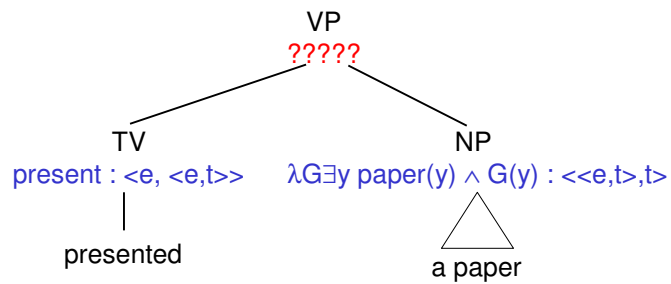
Back to nouns and verbs



Quantificational NPs and transitive verbs

A composition problem:

- *every student* $\Rightarrow \lambda F \forall x (\text{student}(x) \rightarrow F(x)) : \langle \langle e, t \rangle, t \rangle$
- *a paper* $\Rightarrow \lambda G \exists y (\text{paper}(y) \wedge G(y)) : \langle \langle e, t \rangle, t \rangle$
- *presented* $\Rightarrow \text{present} : \langle e, \langle e, t \rangle \rangle$



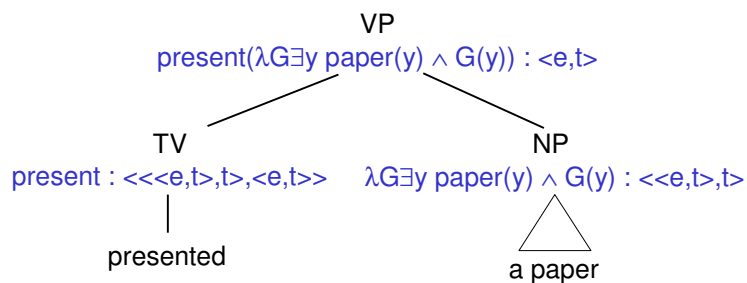
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An attempt at a solution

Raise the type of the first-order relation:

present: $\langle \langle \langle e, t \rangle, t \rangle, \langle e, t \rangle \rangle$



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Spelling out the meaning of a transitive verb

- But now our semantic representation no longer beta-reduces to a FOL formula!

$$\forall x \text{ student}(x) \rightarrow \text{present}(\lambda G \exists y \text{ paper}(y) \wedge G(y))(x)$$

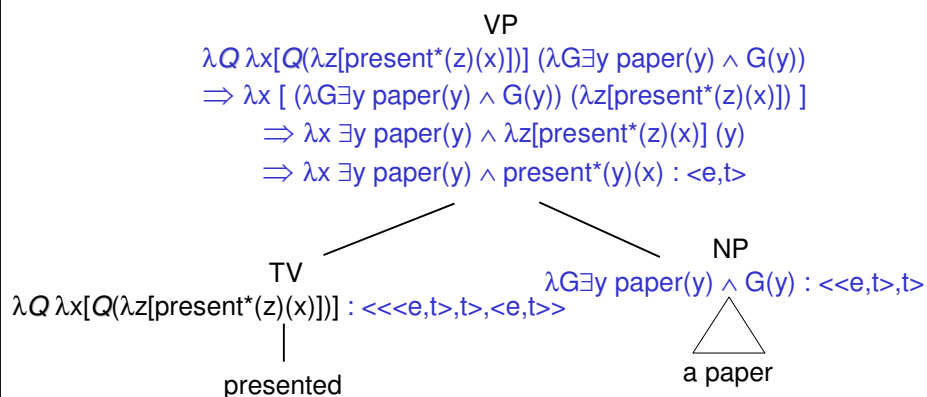
- Same problem as above, same solution.
- Represent transitive verbs like "present" as follows:

$$\lambda Q \lambda x [Q(\lambda y [\text{present}^*(y)(x)])] : \langle \langle e, t \rangle, t \rangle, \langle e, t \rangle \rangle,$$

where $\text{present}^* : \langle e, \langle e, t \rangle \rangle$

- This is hard to read, but it does the trick.

Spelling out the meaning of a transitive verb



Non-referential arguments

- John finds a unicorn $\models \exists x \text{ unicorn}'(x)$
- John seeks a unicorn $\not\models \exists x \text{ unicorn}'(x)$

- Subject position of verbs is always **referential**.
- Direct object position of some verbs is referential, of some other verbs isn't.
- That is, not all transitive verbs can be spelled out in the same way as "present" or "find".
- I.e. there are linguistic reasons why the monster type for transitive verbs makes sense.

Let's look back at compositionality

- Principle of Compositionality:
 - The meaning of a complex expression is uniquely determined by the meanings of its sub-expressions and its syntactic structure.
- But if we assume this, we should be allowed to exchange sub-expressions with the same interpretation without changing the meaning ("salva veritate" substitutability).

Substitutability

- From the denotational version of the Principle of Compositionality, a substitution principle follows:
- If A is sub-expression in sentence C, and A and B have identical denotation, then A can be replaced by B in C without affecting C's truth value.

George W. Bush is married to Laura Bush.

"George W. Bush" = "the American president"

Therefore:

The American president is married to Laura Bush.

Substitutability ?

In 1977, George W. Bush married Laura Bush.

"George W. Bush" = "the American president"

Therefore:

In 1977, the American president married Laura Bush.

??

Substitutability ?

In 1977, the American president was a Democrat.

"George W. Bush" = "the American president"

Therefore:

In 1977, George W. Bush was a Democrat. ???

Substitutability ?

*By constitution, the American president is the Supreme
Commander of the Armed Forces.*

"George W. Bush" = "the American president"

Therefore:

*By constitution, George W. Bush is the Supreme
Commander of the Armed Forces. ???*

Substitutability ?

The following sentences are true:

- The weather is bad
- A semantics lecture is taking place
- $2 + 2 = 4$

How about these?

- It is not the case that ...
- Necessarily ...
- Yesterday, it was the case that ...
- John believes that ...

The lesson

- There are expressions that our semantic representations assign the same interpretation, but which cannot always be exchanged for each other without changing the meaning of the sentence.
- Two possibilities:
 - we have to give up compositionality
 - there are meaning distinctions that our representations don't capture

Extensions vs. intensions

- Two concepts have the same **extension** if they have the same interpretations:
 - "semantics lecture taking place" and " $2 + 2 = 4$ " are both true right now
 - "George W. Bush" and "the US president" refer to the same individual
- However, extensionally equal concepts may still have different "senses": General truths vs. statements that may become false; can believe in one but not the other...
- These senses are also called **intensions**.

Intensions are everywhere

- If we ignore intensions, we will get into trouble with substitutability in a lot of contexts:
 - propositional attitudes: verbs like "believe", "know", "doubt", "desire", ...
 - verbs of saying: "say", "claim", ...
 - tensed sentences (past, future, ...), temporal adverbs (sometimes, always, lately, tomorrow) and connectives (before, during)
 - modal adverbs (necessarily, perhaps), modal verbs (can, may, must, ...), counterfactual conditionals

Modelling intensions

- In order to capture the meaning of a NL expression completely, we must extend the logic to talk about intensions.
- Standard technique:
 - Introduce the concept of a "possible world";
 - define the extension of a term in each possible world;
 - the intension is the mapping of possible worlds to extensions.

Formal models of intensionality

- Intensional Logic (IL): Extend type theory with
 - mechanisms for talking about possible worlds (modal logic)
 - mechanisms for talking about time (temporal logic)
 - mechanisms for abstracting over possible worlds
- Montague Grammar: use IL in semantics construction (basic ideas as presented here)
- We will now look into modal and temporal aspects in some more detail.

Propositional Modal Logic

- Formulas of propositional modal logic: The smallest set such that:
 - Propositional constants are in For
 - If A, B are in For, so are $\neg A$, $(A \wedge B)$, $(A \vee B)$, $(A \rightarrow B)$, $(A \leftrightarrow B)$, $\Box A$, $\Diamond A$

Model Structure

- Model structure for propositional modal logic:
 $M = \langle W, V \rangle$
 - W is a non-empty set (set of possible worlds)
 - V is value assignment function, which assigns each propositional constant a function $W \rightarrow \{0,1\}$
For $V(p)(w)$ we also write $V_w(p)$ or $V_{M,w}(p)$.

Interpretation

- Interpretation of formulas (with respect to model structure M and possible world w):

$[[p]]^{M,w} = V_M(p)(w)$, if p propositional constant

$[[\neg\varphi]]^{M,w} = 1$ iff $[[\varphi]]^{M,w} = 0$

$[[\varphi \wedge \psi]]^{M,w} = 1$ iff $[[\varphi]]^{M,w} = 1$ and $[[\psi]]^{M,w} = 1$

$[[\varphi \vee \psi]]^{M,w} = 1$ iff $[[\varphi]]^{M,w} = 1$ or $[[\psi]]^{M,w} = 1$

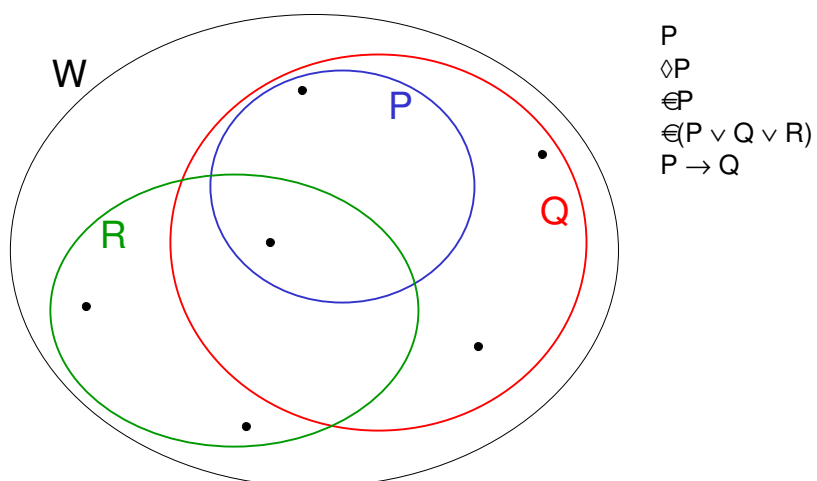
$[[\varphi \rightarrow \psi]]^{M,w} = 1$ iff $[[\varphi]]^{M,w} = 0$ or $[[\psi]]^{M,w} = 1$

$[[\varphi \leftrightarrow \psi]]^{M,w} = 1$ iff $[[\varphi]]^{M,w} = [[\psi]]^{M,w}$

$[[\diamond\varphi]]^{M,w} = 1$ iff $[[\varphi]]^{M,w'} = 1$ for at least one $w' \in W$

$[[\Box\varphi]]^{M,w} = 1$ iff $[[\varphi]]^{M,w'} = 1$ for all $w' \in W$

Propositional modal logic



Substitutability ?

Let the following sentences be true:

- The weather is bad
- A semantics lecture is taking place.
- $2+2=4$

It is not the case that ...

Necessarily ...

Yesterday, it was the case that ...

John believes that ...

Propositional Tense Logic

- Formulas of propositional tense logic: The smallest set such that:
 - Propositional constants are in For
 - If A, B are in For, so are $\neg A$, $(A \wedge B)$, $(A \vee B)$, $(A \rightarrow B)$, $(A \leftrightarrow B)$, **FA**, **GA**, **PA**, **HA**

FA – "it will at some stage be the case that A"

GA – "it is always going to be the case that A"

PA – "it was at some stage the case that A"

HA – "it always has been the case that A"

Model Structure

- Model structure for propositional tense logic (with linear time):

$M = \langle T, <, V \rangle$

- T is non-empty set (set of points in time)
- $<$ is a strict ordering relation on T
- V is value assignment function, which assigns each propositional constant a function $T \rightarrow \{0, 1\}$

For $V(p)(t)$ we also write $V_t(p)$ or $V_{M,t}(p)$

Interpretation

- Interpretation of formulas (with respect to model structure M and time t):

$[[p]]^{M,t} = V_{M,t}(p)$, if p propositional constant

$[[\neg\varphi]]^{M,t} = 1$ iff $[[\varphi]]^{M,t} = 0$

$[[\varphi \wedge \psi]]^{M,t} = 1$ iff $[[\varphi]]^{M,t} = 1$ and $[[\psi]]^{M,t} = 1$

$[[\varphi \vee \psi]]^{M,t} = 1$ iff $[[\varphi]]^{M,t} = 1$ or $[[\psi]]^{M,t} = 1$

$[[\varphi \rightarrow \psi]]^{M,t} = 1$ iff $[[\varphi]]^{M,t} = 0$ or $[[\psi]]^{M,t} = 1$

$[[\varphi \leftrightarrow \psi]]^{M,t} = 1$ iff $[[\varphi]]^{M,t} = [[\psi]]^{M,t}$

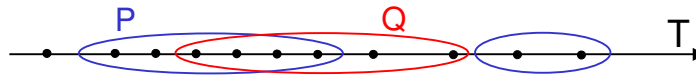
$[[F\varphi]]^{M,t} = 1$ iff $[[\varphi]]^{M,t'} = 1$ for at least one $t' > t$

$[[G\varphi]]^{M,t} = 1$ iff $[[\varphi]]^{M,t'} = 1$ for all $t' > t$

$[[P\varphi]]^{M,t} = 1$ iff $[[\varphi]]^{M,t'} = 1$ for at least one $t' < t$

$[[H\varphi]]^{M,t} = 1$ iff $[[\varphi]]^{M,t'} = 1$ for all $t' < t$

Propositional tense logic



FP
HQ
 $P(P \wedge FQ)$
 $FP \vee P \vee PP$

Substitutability ?

Let the following sentences be true:

- The weather is bad
- A semantics lecture is taking place.
- $2+2=4$

It is not the case that ...

Necessarily ...

Yesterday, it was the case that ...

John believes that ...

Propositional Logic with Tense and Modality

- Syntax: Tense + modal operators
- Model structure: $M = \langle W, T, <, V \rangle$ with
 $V(p): W \times T \rightarrow \{0,1\}$
alternative notation: $V_{M,w,t}(p)$
- Interpretation with respect to M, w and t .

Semantics for FOL with tense and modalities

- Model structure: $M = \langle U, W, T, <, V \rangle$
 - V (or V_M) is value assignment function for non-logical constants, which assigns
 - individuals ($\in U_M$) to individual constants
 - functions $W \times T \rightarrow U^n$ to n -place relational constants
- Assignment function for variables $g: IV \rightarrow U_M$

Interpretation of Terms

- Interpretation of terms (with respect to model structure M and variable assignment g):

$$[[\alpha]]^{M,g,w,t} = V_M(\alpha), \text{ if } \alpha \text{ individual constant}$$

$$[[\alpha]]^{M,g,w,t} = g(\alpha), \text{ if } \alpha \text{ variable}$$

- Notice: Interpretation of terms doesn't depend on the world and time.

Interpretation

- Interpretation of formulas (with respect to model structure M , variable assignment g , world w and time t):

$$[[R(t_1, \dots, t_n)]]^{M,g,w,t} = 1 \text{ iff } \langle [[t_1]]^{M,g,w,t}, \dots, [[t_n]]^{M,g,w,t} \rangle \in V_M(R)(w,t)$$

$$[[s=t]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[s]]^{M,g,w,t} = [[t]]^{M,g,w,t}$$

$$[[\neg\phi]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[\phi]]^{M,g,w,t} = 0$$

$$[[\phi \wedge \psi]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[\phi]]^{M,g,w,t} = 1 \text{ and } [[\psi]]^{M,g,w,t} = 1$$

etc.

$$[[\exists x\phi]]^{M,g,w,t} = 1 \quad \text{iff} \quad \text{there is } a \in U_M \text{ such that } [[\phi]]^{M,g[x/a],w,t} = 1$$

$$[[\forall x\phi]]^{M,g,w,t} = 1 \quad \text{iff} \quad \text{for all } a \in U_M : [[\phi]]^{M,g[x/a],w,t} = 1$$

$$[[F\phi]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[\phi]]^{M,g,w,t'} = 1 \text{ for at least one } t' > t$$

etc.

$$[[\in A]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[\phi]]^{M,g,w,t} = 1 \text{ for all } w' \in W, \text{ etc.}$$

Non-Substitutability: Explained!

In 1977, George W. Bush married Laura Bush.

"George W. Bush" = "the American president"

*Therefore: In 1977, the American president married
Laura Bush. ??*

*George W. Bush has always been married to Laura
Bush.*

George W. Bush is the American president.

*Therefore: The American president has always been
married to Laura Bush ???*

Non-Substitutability: Explained!

*By constitution, the American president is the Supreme
Commander of the Armed Forces.*

George W. Bush is the American president.

*Therefore: By constitution, George W. Bush is the
Supreme Commander of the Armed Forces. ???*

Conclusion

- We extended our type theoretical semantics construction algorithm in two ways:
 - semantics construction for some new semantic phenomena (adjectives, transitive verbs)
 - extension of FOL with intensional constructs (modality, tense)
- Type theory + tense + modality + intensional abstractions = Intensional Logic (not discussed here).