

# Semantic Theory

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Modal and Tense logic

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## Two levels of interpretation

- Semantic analysis of a NL expression in a logical framework is a two-step process – construction of a semantic representation and its truth-conditional interpretation. Accordingly, the Compositionality Principle has two version:
- The semantic representation of a NL expression is uniquely determined by the semantic representation of its sub-expressions, and the way they are syntactically combined.
- The denotation of a semantic representation is uniquely determined by the denotation of its sub-expressions (and their syntactic combination).

## Substitutability

- From the denotational version of the Principle of Compositionality, a substitution principle follows:
- If A is sub-expression in sentence C, and A and B have identical denotation, then A can be replaced by B in C without affecting C's truth value ("salva veritate" substitutability).

*George W. Bush is married to Laura Bush.*

*George W. Bush is the American president.*

Therefore: *The American president is married to Laura Bush*

## Substitutability ?

*George W. Bush has always been married to Laura Bush.*

*George W. Bush is the American president.*

Therefore: *The American president has always been married to Laura Bush ???*

*By constitution, the American president is the Supreme Commander of the Armed Forces.*

*George W. Bush is the American president.*

Therefore: *By constitution, George W. Bush is the Supreme Commander of the Armed Forces. ???*

## Substitutability ?

Let the following sentences be true:

- The weather is nice
- Bill is working
- $2+2=4$

It is not the case that ...

Necessarily ...

Yesterday, it was the case that ...

## Extensions and Intensions

- Meanings of expressions are incompletely modelled by interpretation via (extensional) models.
- Moreover, we need the missing parts of the meaning to interpret important types of expressions and constructions:
- Tensed sentences (past, future, ...), temporal adverbs (sometimes, always, lately, tomorrow) and connectives (before, during) require to look into prior and future states of the world.
- Modal adverbs (necessarily, perhaps), modal verbs (can, may, must, ...), counterfactual conditionals require to look into alternative states of the world.
- For a more complete account of meaning information, we define denotations as intensions: functions from points in time / possible worlds to extensions (i.e., truth values, individuals, sets etc.)

## Propositional Modal Logic

- Formulas of propositional modal logic: The smallest set such that:
  - Propositional constants are in For
  - If A, B are in For, so are  $\neg A$ ,  $(A \wedge B)$ ,  $(A \vee B)$ ,  $(A \rightarrow B)$ ,  $(A \leftrightarrow B)$ ,  $\Box A$ ,  $\Diamond A$

## Model Structure

- Model structure for propositional modal logic:  
 $M = \langle W, V \rangle$ 
  - $W$  is a non-empty set (set of possible worlds)
  - $V$  is value assignment function, which assigns each propositional constant a function  $W \rightarrow \{0,1\}$   
For  $V(p)(w)$  we also write  $V_w(p) - V_{M,w}(p)$  respectively

## Interpretation

- Interpretation of formulas (with respect to model structure M and possible world w):

$[[p]]^{M,w} = V_M(p)(w)$ , if p propositional constant

$[[\neg\phi]]^{M,w} = 1$                       iff             $[[\phi]]^{M,w} = 0$

$[[\phi \wedge \psi]]^{M,w} = 1$       iff             $[[\phi]]^{M,w} = 1$  and  $[[\psi]]^{M,w} = 1$

$[[\phi \vee \psi]]^{M,w} = 1$       iff             $[[\phi]]^{M,w} = 1$  or  $[[\psi]]^{M,w} = 1$

$[[\phi \rightarrow \psi]]^{M,w} = 1$     iff             $[[\phi]]^{M,w} = 0$  or  $[[\psi]]^{M,w} = 1$

$[[\phi \leftrightarrow \psi]]^{M,w} = 1$     iff             $[[\phi]]^{M,w} = [[\psi]]^{M,w}$

$[[\diamond \phi]]^{M,w} = 1$             iff             $[[\phi]]^{M,w'} = 1$  for at least one  $w' \in W$

$[[\Box \phi]]^{M,w} = 1$             iff             $[[\phi]]^{M,w'} = 1$  for all  $w' \in W$

## Substitutability ?

Let the following sentences be true:

- The weather is nice
- Bill is working
- $2+2=4$

It is not the case that ...

Necessarily ...

Yesterday, it was the case that ...

## Validity, Entailment, Deduction

- According to the semantics of propositional modal logic, there are additional valid propositions and entailment relations, which again can be cast into deduction calculi.
- Examples:  $\models \Box A \rightarrow A$  ( $\Box A \models A$ ,  $\Box A \models \neg A$ )  
 $\models A \rightarrow \Diamond A$   
 $\models \Diamond \Box A \rightarrow A$   
 $\models \Box(A \vee \neg A)$

## Propositional Tense Logic

- Formulas of propositional tense logic: The smallest set such that:
  - Propositional constants are in For
  - If A, B are in For, so are  $\neg A$ ,  $(A \wedge B)$ ,  $(A \vee B)$ ,  $(A \rightarrow B)$ ,  $(A \leftrightarrow B)$ , **FA**, **GA**, **PA**, **HA**

**FA** – "it will at some stage be the case that A"

**GA** – "it is always going to be the case that A"

**PA** – "it was at some stage the case that A"

**HA** – "it always has been the case that A"

## Model Structure

- Model structure for propositional tense logic:  
 $M = \langle T, <, V \rangle$ 
  - $T$  is non-empty set (set of points in time)
  - $<$  is a strict ordering relation on  $T$
  - $V$  is value assignment function, which assigns each propositional constant a function  $T \rightarrow \{0,1\}$   
 For  $V(p)(t)$  we also write  $V_t(p) - V_{M,t}(p)$  respectively

## Interpretation

- Interpretation of formulas (with respect to model structure  $M$  and time  $t$ ):  
 $[[p]]^{M,t} = V_M(p)(w)$ , if  $p$  propositional constant
- $[[\neg\varphi]]^{M,t} = 1$       iff       $[[\varphi]]^{M,t} = 0$
- $[[\varphi \wedge \psi]]^{M,t} = 1$       iff       $[[\varphi]]^{M,t} = 1$  and  $[[\psi]]^{M,t} = 1$
- $[[\varphi \vee \psi]]^{M,t} = 1$       iff       $[[\varphi]]^{M,t} = 1$  or  $[[\psi]]^{M,t} = 1$
- $[[\varphi \rightarrow \psi]]^{M,t} = 1$       iff       $[[\varphi]]^{M,t} = 0$  or  $[[\psi]]^{M,t} = 1$
- $[[\varphi \leftrightarrow \psi]]^{M,t} = 1$       iff       $[[\varphi]]^{M,t} = [[\psi]]^{M,t}$
- $[[F\varphi]]^{M,t} = 1$  iff       $[[\varphi]]^{M,t'} = 1$  for at least one  $t' > t$
- $[[G\varphi]]^{M,t} = 1$  iff       $[[\varphi]]^{M,t'} = 1$  for all  $t' > t$
- $[[P\varphi]]^{M,t} = 1$  iff       $[[\varphi]]^{M,t'} = 1$  for at least one  $t' < t$
- $[[H\varphi]]^{M,t} = 1$  iff       $[[\varphi]]^{M,t'} = 1$  for all  $t' < t$

## Propositional Logic with Tense and Modality

- Syntax: Tense + modal operators
- Model structure:  $M = \langle W, T, <, V \rangle$  with  
 $V(p): W \times T \rightarrow \{0,1\}$   
alternative notation:  $V_{M,w,t}(p)$
- Interpretation with respect to  $M, w$  and  $t$ .

## Semantics for FOL with tense and modalities

- Model structure:  $M = \langle U, W, T, <, V \rangle$ 
  - $V (V_M)$  is value assignment function for non-logical constants, which assigns
    - individuals ( $\in U_M$ ) to individual constants
    - functions  $W \times T \rightarrow U^n$  to n-place relational constants
- Assignment function for variables  $g: IV \rightarrow U_M$



## Interpretation of Terms

- Interpretation of terms (with respect to model structure M and variable assignment g):

$$[[\alpha]]^{M,g,w,t} = V_M(\alpha), \text{ if } \alpha \text{ individual constant}$$

$$[[\alpha]]^{M,g,w,t} = g(\alpha), \text{ if } \alpha \text{ variable}$$

## Interpretation

- Interpretation of formulas (with respect to model structure M, variable assignment g, world w and time t):

$$[[R(t_1, \dots, t_n)]]^{M,g,w,t} = 1 \text{ iff } \langle [[t_1]]^{M,g,w,t}, \dots, [[t_n]]^{M,g,w,t} \rangle \in V_M(R)(w,t)$$

$$[[s=t]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[s]]^{M,g,w,t} = [[t]]^{M,g,w,t}$$

$$[[\neg\phi]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[\phi]]^{M,g,w,t} = 0$$

$$[[\phi \wedge \psi]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[\phi]]^{M,g,w,t} = 1 \text{ and } [[\psi]]^{M,g,w,t} = 1$$

etc.

$$[[\exists x\phi]]^{M,g,w,t} = 1 \quad \text{iff} \quad \text{there is } a \in U_M \text{ such that } [[\phi]]^{M,g[x/a],w,t} = 1$$

$$[[\forall x\phi]]^{M,g,w,t} = 1 \quad \text{iff} \quad \text{for all } a \in U_M : [[\phi]]^{M,g[x/a],w,t} = 1$$

$$[[F\phi]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[\phi]]^{M,g,w,t'} = 1 \text{ for at least one } t' > t$$

etc.

$$[[\in A]]^{M,g,w,t} = 1 \quad \text{iff} \quad [[\phi]]^{M,g,w,t} = 1 \text{ for all } w' \in W, \text{ etc.}$$

## *De dicto/de re readings*

*De dicto/de re readings:*

- *The American president has always been married to Laura Bush*

Only *de re* reading is true.

- *By constitution, the American president is the Supreme Commander of the Armed Forces.*

Only *de dicto* reading is true.