

Natural Language Parsing Technology

Foundations of Language Science and Technology (WS 2009/2010)

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Outline

- 1 Overview
- 2 Basic Parsing Algorithms
 - Parsing Strategies
 - CYK Algorithm
 - Earley's Algorithm
- 3 Parsing with Probabilistic Context-Free Grammar
 - PCFG
 - Inside-Outside Algorithm
- 4 Recent Advances in Parsing Technology

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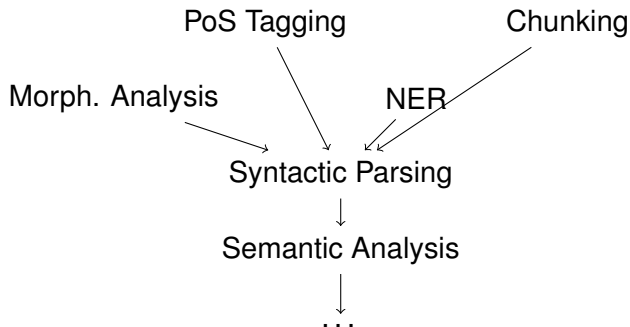
Language & Grammar

- Language
 - Structural
 - Productive
 - Ambiguous, yet efficient in human-human communication
- Grammar
 - Generalization of regularities in language structures
 - Morphology & syntax, often complemented by phonetics, phonology, semantics, and pragmatics

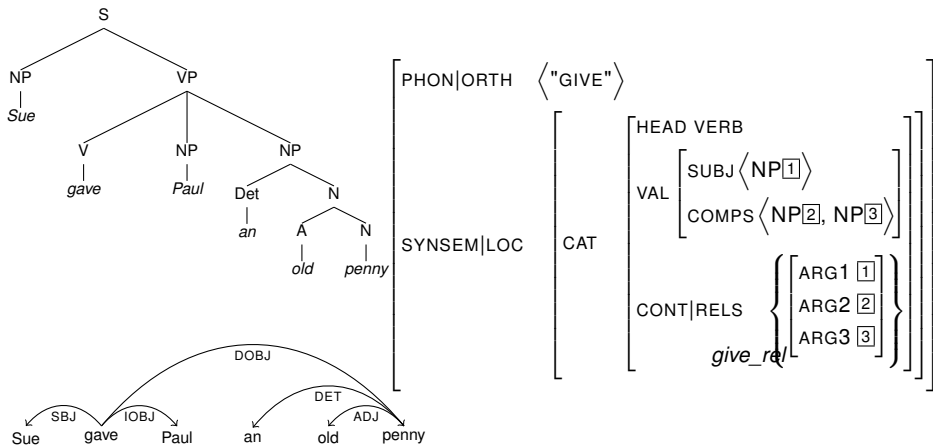
Ambiguity

- Human languages are ambiguous on almost every layer
- Grammar frameworks are designed to represent necessary ambiguities, and eliminate unnecessary ones
- Parsing models are responsible for retrieving valid analyses according to the grammar

Syntactic Parser as NLP Component



Trees (or not)



Chomsky Hierarchy

- Type 0 (unrestricted rewriting system)

$$\alpha \rightarrow \beta \quad \alpha, \beta \in (V_N \cup V_T)^*$$

- Type 1 (context sensitive grammars)

$$\phi A \omega \rightarrow \phi \beta \omega \quad A \in V_N, \beta, \phi, \omega \in (V_N \cup V_T)^*$$

- Type 2 (context free grammars)

$$A \rightarrow \beta \quad A \in V_N, \beta \in (V_N \cup V_T)^*$$

- Type 3 (regular grammars)

$$A \rightarrow xB \vee A \rightarrow x \quad A, B \in V_N, x \in V_T$$

Context-Free Grammar

A CFG is a quadruple: $\langle V_T, V_N, \mathcal{P}, S \rangle$

- V_T : terminal symbols
- V_N : non-terminal symbols
- \mathcal{P} : context-free productions

$$A \rightarrow \beta \quad A \in V_N, \beta \in (V_N \cup V_T)^*$$

- S : start symbol

Context-Free Phrase Structure Grammar

- $S \rightarrow NP VP$
- $NP \rightarrow Det N$
- $NP \rightarrow Adj NP$
- $VP \rightarrow V$
- $VP \rightarrow V NP$
- $VP \rightarrow Adv VP$
- $N \rightarrow dog|cat$
- $Det \rightarrow the|a$
- $V \rightarrow chases|sleeps$
- $Adj \rightarrow gray|lazy$
- $Adv \rightarrow fiercely$

CFG Derivation

- If $\phi = \beta A \gamma$, $\omega = \beta \alpha \gamma$ and $A \rightarrow \alpha \in \mathcal{P}$
then ω follows ϕ , $\phi \Rightarrow \omega$
- If a sequence of strings $\phi_1, \phi_2, \dots, \phi_m$ where for all i
($1 \leq i \leq m - 1$), $\phi_i \Rightarrow \phi_{i+1}$
then $\phi_1, \phi_2, \dots, \phi_m$ is a derivation from ϕ_1 to ϕ_m
- **“Derivable”** relation: transitive, reflexive

$$\phi_1 \xRightarrow{*} \phi_m$$

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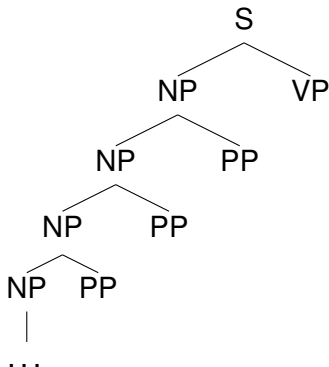
Parsing Strategies

- Top-down: start from the start symbol, and expand the tree with grammar rules (e.g. replace LHS symbol with RHS sequences of CFG productions)
- Bottom-up: start from the input sequence, and apply grammar rules to build trees upwards (e.g. reducing RHS sequence into LHS symbols)

Top-Down Parsing

- Goal-directed search
- Waste time on trees that do not match input sentence
- Pure top-down (left-first) approach cannot parse (left-)recursion grammars

- 1 $S \rightarrow NP VP$
- 2 $NP \rightarrow NP PP$
- 3 ...



Bottom-Up Parsing

- Use the input to guide the search (data-driven)
- Waste time on trees that don't result in S
- Recursive unary rules still create an infinite parse forest for a finite length sentence

$$① \quad A \rightarrow B|a$$

$$② \quad B \rightarrow A$$

$$③ \quad \dots$$

\dots
 $|$
 B
 $|$
 A
 $|$
 B
 $|$
 A
 $|$
 a

Problems

- Left-recursion $NP \rightarrow NP PP$
- Ambiguity
- Repeated parsing of subtrees

Dynamic Programming (DP)

- Divisibility: the optimal solution of a sub problem is part of the optimal solution of the whole problem
- Memorization: solve small problems only once and remember the answers

Example

Calculating Fibonacci numbers:

$$F_n = F_{n-1} + F_{n-2} \quad (F_0 = 0, F_1 = 1)$$

CYK Algorithm

- Cocke-Younger-Kasami, also known as CKY algorithm
- Essentially a bottom-up chart parsing algorithm using dynamic programming
- CFG is in Chomsky Normal Form (CNF)
 - $A \rightarrow BC$
 - $A \rightarrow a$
 - $S \rightarrow \epsilon$
 - $A, B, C \in V_N, \quad a \in V_T, \quad B, C \neq S$
- Fill in a two-dimension array: $\mathbb{C}[i][j]$ contains all the possible syntactic interpretations of the substring $w_{i+1} \dots w_j$
- Complexity $O(n^3)$

CYK Algorithm

```
1: for all  $i, j \quad 0 \leq i < j \leq n$  do
2:    $C[i][j] \leftarrow \phi$ 
3: end for
4: for all  $A \rightarrow w_i \in \mathcal{P}$  do
5:    $C[i-1][i] \leftarrow \{A\} \cup C[i-1][i]$ 
6: end for
7: for  $s = \langle 2 \dots n \rangle$  do
8:   for all  $A \rightarrow B C \in \mathcal{P}, i, k : 0 \leq i < k < i + s$  do
9:     if  $B \in C[i][k] \wedge C \in C[k][i + s]$  then
10:       $C[i][i + s] \leftarrow \{A\} \cup C[i][i + s]$ 
11:    end if
12:  end for
13: end for
```

CYK Chart: Example

	The	man	saw	a	girl	with	a	telescope
1	Det	N	V	Det	N	P	Det	N
2	NP			NP			NP	
3			VP			PP		
4								
5	S			NP				
6			VP					
7								
8	S							

- 1 $S \rightarrow NP VP$
- 2 $NP \rightarrow Det N$
- 3 $NP \rightarrow NP PP$
- 4 $VP \rightarrow V NP$
- 5 $VP \rightarrow VP PP$
- 6 $PP \rightarrow P NP$
- 7 $V \rightarrow saw$
- 8 $N \rightarrow man|girl|telescope$
- 9 $Det \rightarrow a|the$
- 10 $P \rightarrow with$

CYK Chart: Example

	The	man	saw	a	girl	with	a	telescope
1	Det	N	V	Det	N	P	Det	N
2	NP			NP			NP	
3			VP			PP		
4								
5	S			NP				
6			VP					
7								
8	S							

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Earley's Algorithm

- Use dynamic programming to do top-down search
- Chart: a set of items $\langle h, i, A \rightarrow \alpha \cdot \beta \rangle$
 - h, i : positions in the input $0 \leq h \leq i \leq n$
 - $A \rightarrow \alpha \cdot \beta$: dotted rule ($A \rightarrow \alpha\beta \in \mathcal{P}$)
 - α : RHS prefix that has already been applied to input from h to i
 - β : RHS suffix yet to be found

Earley's Algorithm

- **Initialize**

foreach $S \rightarrow \alpha \in \mathcal{P}$
 $\mathbb{C} \leftarrow \langle 0, 0, S \rightarrow \cdot \alpha \rangle$

- **Scan(i)**

if $w_i = a \wedge \langle h, i-1, A \rightarrow \alpha \cdot a \beta \rangle \in \mathbb{C}$
 $\mathbb{C} \leftarrow \langle h, i, A \rightarrow \alpha a \cdot \beta \rangle$

- **Complete(i)**

foreach $\langle h, i, A \rightarrow \alpha \cdot \rangle \in \mathbb{C}$
 foreach $\langle k, h, B \rightarrow \beta \cdot A \gamma \rangle \in \mathbb{C}$
 $\mathbb{C} \leftarrow \langle k, i, B \rightarrow \beta A \cdot \gamma \rangle$

- **Predict(i)**

foreach $\langle h, i, A \rightarrow \alpha \cdot B \beta \rangle \in \mathbb{C}$
 foreach $B \rightarrow \gamma \in \mathcal{P}$
 $\mathbb{C} \leftarrow \langle i, i, B \rightarrow \cdot \gamma \rangle$

- **Parse**

Initialize
 for $i = \langle 1 \dots n \rangle$
 Predict($i-1$)
 Scan(i)
 Complete(i)
 if $\exists \langle 0, n, S \rightarrow \alpha \cdot \rangle \in \mathbb{C}$
 return *success*
 else
 return *failed*

Earley Chart: Example

0 the/det 1 dog/n 2 chases/v 3 a/det 4 cat/n 5

	0	1	2	3	4	5
0	$S \rightarrow \cdot NP VP$ $NP \rightarrow \cdot det n$					
1	$NP \rightarrow det \cdot n$					
2	$NP \rightarrow det n \cdot$ $S \rightarrow NP \cdot VP$		$VP \rightarrow \cdot v$ $VP \rightarrow \cdot v NP$			
3	$S \rightarrow NP VP \cdot$		$VP \rightarrow v \cdot$ $VP \rightarrow v \cdot NP$	$NP \rightarrow \cdot det n$		
4				$NP \rightarrow det \cdot n$		
5	$S \rightarrow NP VP \cdot$		$VP \rightarrow v NP \cdot$	$NP \rightarrow det n \cdot$		

- 1 $S \rightarrow NP VP$
- 2 $VP \rightarrow v NP$
- 3 $VP \rightarrow v$
- 4 $NP \rightarrow det n$

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Probabilistic Context-Free Grammar

An PCFG is a quintuple: $\langle V_T, V_N, \mathcal{P}, S, \mathbf{Pr} \rangle$

- $\mathbf{Pr} : \mathcal{P} \rightarrow [0, 1]$ s.t.

$$\forall A \in V_N, \sum_{A \rightarrow \alpha \in \mathcal{P}} \mathbf{Pr}(A \rightarrow \alpha) = 1$$

- $\mathbf{Pr}(A \rightarrow \alpha)$ can be understood as the conditional probability of observing $A \rightarrow \alpha$ in the derivation given A : $P(A \rightarrow \alpha | A)$

Various Probabilities

Joint Probability: $P(x, y)$

- Input sequence: x
- A Parse: y with corresponding derivation sequence:

$$S \xRightarrow[r_1]{} \phi_1 \xRightarrow[r_2]{} \phi_2 \xRightarrow[r_3]{} \dots \xRightarrow[r_k]{} x$$

where r_i is the production rule used in the i^{th} derivation step

- $P(x, y) = \prod_{i=1}^k \mathbf{Pr}(r_i)$
- $\sum_{y \in \mathcal{T}(G), x = \text{yield}(y)} P(x, y) = 1$
- More generally, $P(x, y|A) = \prod_{i=1}^k \mathbf{Pr}(r_i)$ is the probability of a sub-parse y rooted by A and generate input x by derivation sequence

$$A \xRightarrow[r_1]{} \phi_1 \xRightarrow[r_2]{} \phi_2 \xRightarrow[r_3]{} \dots \xRightarrow[r_k]{} x$$

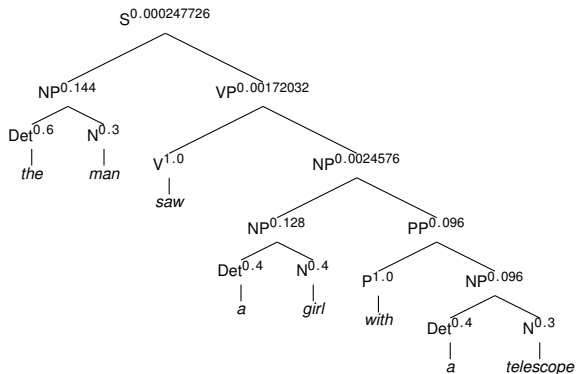
Various Probabilities

Structural Language Model: $P(x)$

- $P(x) = \sum_{y \in \mathcal{T}(x)} P(x, y)$
- $\mathcal{T}(x)$ is the set of parse trees for input sequence x

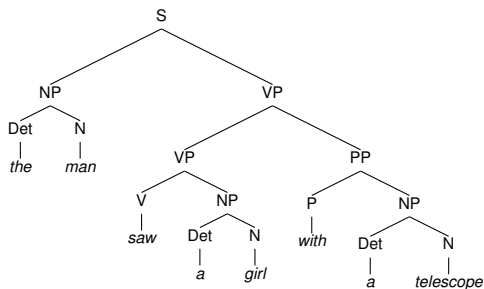
PCFG Example

- 1 $S \rightarrow NP VP$ 1.0
- 2 $NP \rightarrow Det N$ 0.8
- 3 $NP \rightarrow NP PP$ 0.2
- 4 $VP \rightarrow V NP$ 0.7
- 5 $VP \rightarrow VP PP$ 0.3
- 6 $PP \rightarrow P NP$ 1.0
- 7 $V \rightarrow \textit{saw}$ 1.0
- 8 $N \rightarrow \textit{man}$ 0.3
- 9 $N \rightarrow \textit{girl}$ 0.4
- 10 $N \rightarrow \textit{telescope}$ 0.3
- 11 $Det \rightarrow \textit{a}$ 0.4
- 12 $Det \rightarrow \textit{the}$ 0.6
- 13 $P \rightarrow \textit{with}$ 1.0



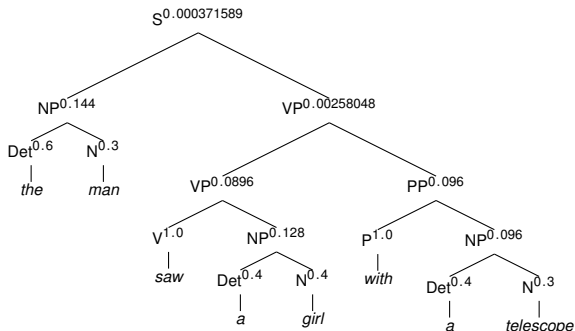
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- 10 $N \rightarrow \textit{telescope}$ 0.3
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- 13 $P \rightarrow \textit{with}$ 1.0



Parsing with PCFG

- Earley and CYK algorithms can be adapted to carry probabilities
- Best parse tree y^* for a sentence x

$$y^* = \operatorname{argmax}_{y \in \mathcal{T}(x)} P(x, y)$$

- N -best parse can be recovered with Viterbi-like algorithm

Learning PCFG Probabilities

- Given a treebank, with Maximum-Likelihood Estimation (MLE):

$$P(A \rightarrow \beta) = \frac{\#(A \rightarrow \beta)}{\#(A)}$$

- When the grammar is large (e.g. by lexicalization), smoothing is necessary to overcome data sparseness

Inside-Outside Algorithm

- When there is no labeled data (treebank), probabilities of a PCFG can be updated to maximize the likelihood over a set of unlabeled sentences

$$\mathbf{Pr}^* = \operatorname{argmax}_{\mathbf{Pr}} \prod_x P(x) = \operatorname{argmax}_{\mathbf{Pr}} \prod_x \sum_{y \in \mathcal{T}(x)} P(x, y)$$

- An Expectation-Maximization procedure can be used to iteratively find \mathbf{Pr}^*

Inside Probability

Definition

Inside probability $\beta_j(p, q)$ is the probability of sequence $w_{p+1} \dots w_q$ being generated with a tree rooted by node N^j

$$\beta_j(p, q) = P(w_{p+1} \dots w_q | N_{pq}^j)$$

- $\beta_1(0, n) = P(w_1 w_2 \dots w_n) \quad N^1 = S$
- Calculation can be carried out bottom-up

$$\beta_j(k-1, k) = Pr(N^j \rightarrow w_k) \quad N^j \in V_N \quad (1)$$

$$\beta_j(p, q) = \sum_{r,s} \sum_{d=p+1}^{q-1} Pr(N^j \rightarrow N^r N^s) \cdot \beta_r(p, d) \cdot \beta_s(d, q) \quad (2)$$

Outside Probability

Definition

Outside probability $\alpha_j(p, q)$ is the total probability of beginning with the start symbol and generating N_{pq}^j and all the words outside

$$\alpha_j(p, q) = P(w_1 \dots w_p, N_{pq}^j, w_{q+1} \dots w_n)$$

- N_{pq}^j means

$$N^j \xRightarrow{*} w_{p+1} \dots w_q$$

- $P(w_1 w_2 \dots w_n, N_{pq}^j) = \alpha_j(p, q) \cdot \beta_j(p, q)$
- $P(w_1 w_2 \dots w_n) = \sum_j \alpha_j(k-1, k) Pr(N^j \rightarrow w_k)$ for any k

Outside Probability (cont.)

- Calculation is top-down

$$\alpha_j(0, n) = \begin{cases} 1 & N^j = S \\ 0 & \text{otherwise} \end{cases} \quad (3)$$

$$\begin{aligned} \alpha_j(p, q) = & \sum_{f,g} \sum_{q < e < n} \alpha_f(p, e) \cdot Pr(N^f \rightarrow N^j N^g) \cdot \beta(q, e) \\ & + \sum_{f,g} \sum_{0 < e < p} \alpha_f(e, q) \cdot Pr(N^f \rightarrow N^g N^j) \cdot \beta(e, p) \end{aligned} \quad (4)$$

Calculating Expected Counts

The expected times N^j is used in the derivation for sentence $w_1 \dots w_n$

$$\begin{aligned}
 E[N^j | w_1 \dots w_n] &= \sum_{p=0}^{n-1} \sum_{q=p+1}^n P(N_{pq}^j | w_1 \dots w_n) & (5) \\
 &= \sum_{p=0}^{n-1} \sum_{q=p+1}^n \frac{P(N_{pq}^j, w_1 \dots w_n)}{P(w_1 \dots w_n)} = \sum_{p=0}^{n-1} \sum_{q=p+1}^n \frac{\alpha_j(p, q) \cdot \beta_j(p, q)}{P(w_1 \dots w_n)}
 \end{aligned}$$

Calculating Expected Counts (cont.)

The expected times $N^j \rightarrow N^r N^s$ and N^j is used in the derivation for sentence $w_1 \dots w_n$

$$\begin{aligned}
 E[N^j \rightarrow N^r N^s | w_1 \dots w_n] &= \sum_{p=0}^{n-1} \sum_{q=p+1}^n P(N_{pq}^j, N^j \rightarrow N^r N^s | w_1 \dots w_n) \quad (6) \\
 &= \frac{\sum_{p=0}^{n-1} \sum_{q=p+1}^n \sum_{d=p+1}^{q-1} \alpha_j(p, q) \cdot Pr(N^j \rightarrow N^r N^s) \cdot \beta_r(p, d) \cdot \beta_s(d, q)}{P(w_1 \dots w_n)}
 \end{aligned}$$

Update Formula

For a single sentence, rule probabilities can be reestimated

$$\begin{aligned} \hat{P}_r(N^j \rightarrow N^r N^s) &= \frac{E[N^j \rightarrow N^r N^s, N^j | w_1 \dots w_n]}{E[N^j | w_1 \dots w_n]} \\ &= \frac{\sum_{p=0}^{n-1} \sum_{q=p+1}^n \sum_{d=p+1}^{q-1} \alpha_j(p, q) \cdot Pr(N^j \rightarrow N^r N^s) \cdot \beta_r(p, d) \cdot \beta_s(d, q)}{\sum_{p=0}^{n-1} \sum_{q=p+1}^n \alpha_j(p, q) \cdot \beta_j(p, q)} \end{aligned} \quad (7)$$

Similarly, for unary rules,

$$\hat{P}_r(N^j \rightarrow w^k) = \frac{\sum_{h=1}^n \alpha_j(h-1, h) \cdot P(w_h = w^k) \cdot \beta_j(h-1, h)}{\sum_{p=0}^{n-1} \sum_{q=p+1}^n \alpha_j(p, q) \cdot \beta_j(p, q)} \quad (8)$$

Multiple Training Sentences

For each sentence \vec{w}^j in the training corpus

$$f_i(p, q, j, r, s) = \frac{\sum_{d=p+1}^{q-1} \alpha_j(p, q) \cdot Pr(N^j \rightarrow N^r N^s) \cdot \beta_r(p, d) \cdot \beta_s(d, q)}{P(w_1 \dots w_n)} \quad (9)$$

$$g_i(h, j, k) = \frac{\alpha_j(h-1, h) \cdot P(w_h = w^k) \cdot \beta_j(h-1, h)}{P(w_1 \dots w_n)} \quad (10)$$

$$h_i(p, q, j) = \frac{\alpha_j(p, q) \cdot \beta_j(p, q)}{P(w_1 \dots w_n)} \quad (11)$$

then

$$\hat{P}_r(N^j \rightarrow N^r N^s) = \frac{\sum_{i=1}^m \sum_{p=0}^{n_i-1} \sum_{q=p+1}^{n_i} f_i(p, q, j, r, s)}{\sum_{i=1}^m \sum_{p=0}^{n_i-1} \sum_{q=p+1}^{n_i} h_i(p, q, j)} \quad (12)$$

$$\hat{P}_r(N^j \rightarrow w^k) = \frac{\sum_{i=1}^m \sum_{h=1}^{n_i} g_i(h, j, k)}{\sum_{i=1}^m \sum_{p=0}^{n_i-1} \sum_{q=p+1}^{n_i} h_i(p, q, j)} \quad (13)$$

Inside-Outside Algorithm

Initialize an arbitrary set of rule probabilities \mathbf{Pr}^0

repeat

$F = G = H \leftarrow 0$

for $\vec{w}^k = w_1^k \dots w_n^k$ in the corpus **do**

Calculate inside probabilities $\beta_j(p, q)$

Calculate outside probabilities $\alpha_j(p, q)$

Accumulate counts F G and H

end for

Update rule probabilities $Pr^{i+1}(N^j \rightarrow N^r N^s)$ and $Pr^{i+1}(N^j \rightarrow w^h)$

until $|P_{Pr^{i+1}}(W) - P_{Pr^i}(W)| \leq \epsilon$

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Statistical Constituent Parsing

- Collins parser [Collins, 1997]
- Charniak Parser [Charniak, 2000]
- Reranking model [Collins and Koo, 2005]
- Self-training [McClosky et al., 2006]

Statistical Dependency Parsing

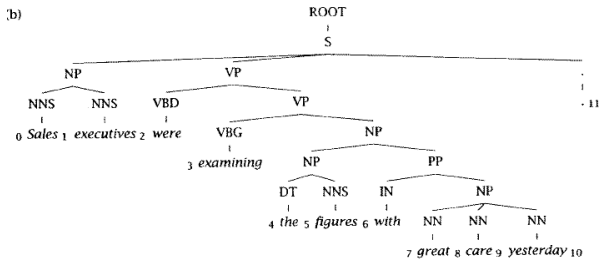
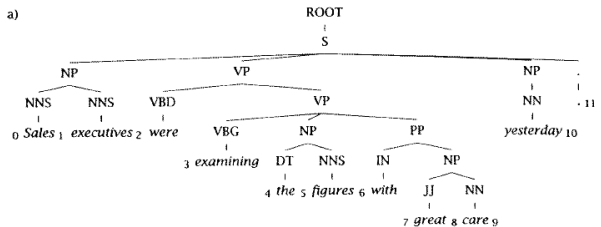
- Graph-based approach [Eisner, 1996, McDonald et al., 2005]
 - Edge-factorized scoring model
 - Efficient algorithms to find maximal spanning tree
 - Allows non-projective dependency structures
- Transition-based approach [Nivre et al., 2007, Sagae and Tsujii, 2008]
 - (Near) deterministic parsing
 - Projective/pseudo-projective

Parsing with Richer Formalisms

- TAG
- CCG
- LFG
- HPSG

Parser Evaluation

- Evaluation against “gold-standard”
 - E.g. PARSEVAL
- Application-based evaluation



(c) Brackets in gold standard tree (a.):

S-(0:11), NP-(0:2), VP-(2:9), VP-(3:9), NP-(4:6), PP-(6:9), NP-(7,9), *NP-(9:10)

(d) Brackets in candidate parse (b.):

S-(0:11), NP-(0:2), VP-(2:10), VP-(3:10), NP-(4:10), NP-(4:6), PP-(6:10), NP-(7,10)

(e) Precision:	3/8 = 37.5%	Crossing Brackets:	0
Recall:	3/8 = 37.5%	Crossing Accuracy:	100%
Labeled Precision:	3/8 = 37.5%	Tagging Accuracy:	10/11 = 90.9%
Labeled Recall:	3/8 = 37.5%		

Domain Adaptability and Multilinguality

- Statistical parsing models usually performs well in in-domain tests and suffer significant accuracy drop when tested on out-of-domain data
- Differences between languages require different parsing models (morphology, word order, etc.)

Open Questions

- How relevant is linguistic study to the development of parsers?
- How do we evaluate a parser?
- How to make trade-offs between adequacy, accuracy and efficiency?

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