Computational Linguistics

Algorithms for Scope Underspecification

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Today

- Refining the solver from the last lecture
- Hypernormally connected dominance graphs
- Tree Automata (in a nutshell)
- Redundancy elimination (very short)

(Bodirsky &al., 2004)

Solving dominance graphs

solve(G = ⟨V, E $\uplus$ D⟩) = (*)

choose a free fragment F of G else fail

let $G_1, \ldots, G_k$ be the WCC's of G\F

let $\langle V_i, E_i \uplus D_i \rangle$ = solve(G_i)

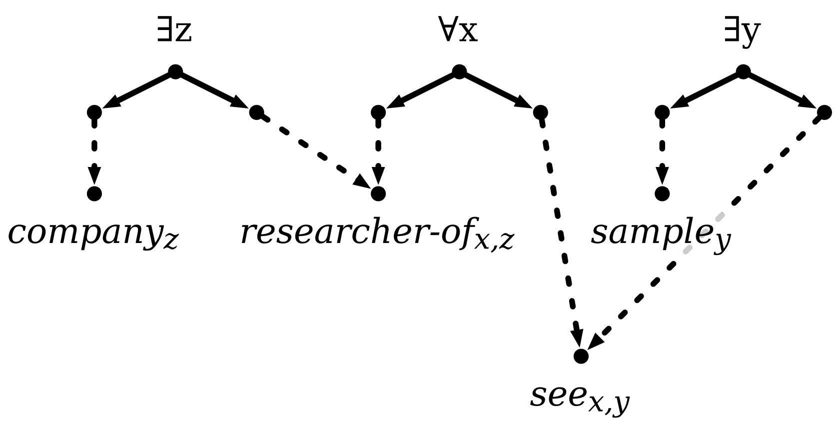
return $\langle V, E \uplus D_1 \uplus \cdots \uplus D_k \uplus D' \rangle$

where D' are dominance edges that connect the holes of F with the solved form of one of the corresponding G_i

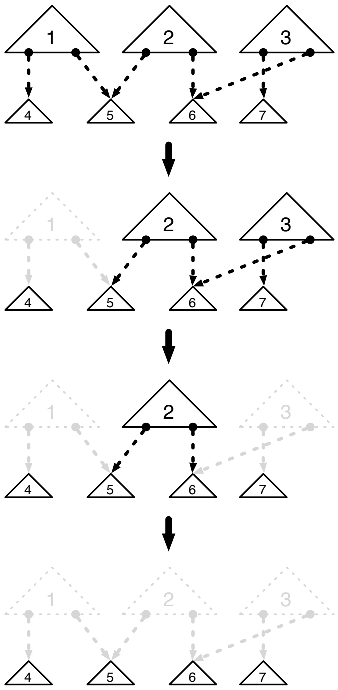
(*) slightly simplified version, works only for connected normal dominance graphs

An Example

- *Every researcher of a company saw a sample.*

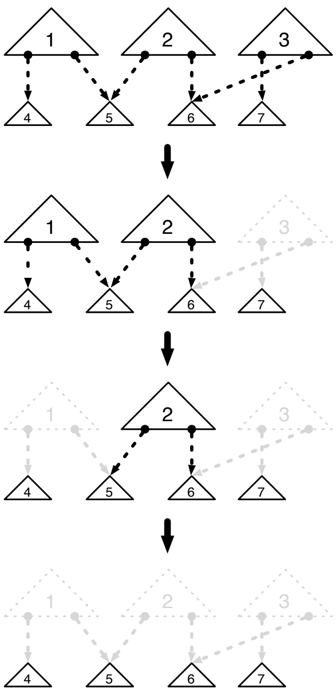


An Example: Run #1



subgraph	free	wccs
{1,...,7}	1	{4}, {2,3,5,6,7}
{2,3,5,6,7}	3	{7}, {2,5,6}
{2,5,6}	2	{5}, {6}

An Example: Run #2

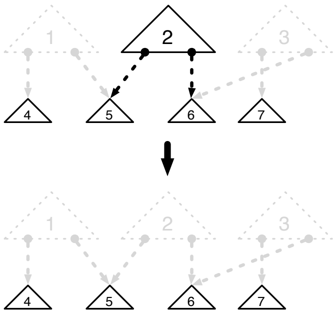


subgraph	free	wccs
{1,...,7}	3	{7}, {1,2,4,5,6}
{1,2,4,5,6}	1	{4}, {2,5,6}
{2,5,6}	2	{5}, {6}

An Example

Problem: The algorithm is applied twice to the subgraph {2,5,6}

subgraph	free	wccs
{1,...,7}	1	{4}, {2,3,5,6,7}
{2,3,5,6,7}	3	{7}, {2,5,6}
{2,5,6}	**2**	**{5}, {6}**



subgraph	free	wccs
{1,...,7}	3	{7}, {1,2,4,5,6}
{1,2,4,5,6}	1	{4}, {2,5,6}
{2,5,6}	**2**	**{5}, {6}**

Chart-based solver

- **Basic idea:** use dynamic programmic techniques and store intermediate results in a chart-like datastructure
- The chart records how graphs are decomposed into smaller subgraphs if free fragments are removed
  - The chart assigns each subgraph a “split.”
  - Splits consist of (references to) a free fragment F and the weakly connected components of G\F.
    - Notation: $F(G_1, \ldots, G_n)$

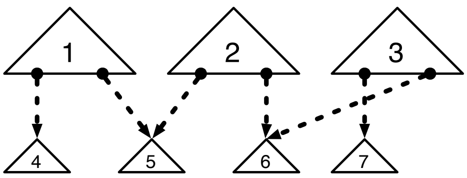
(Koller & Thater, 2005)

The algorithm

```
GRAPH-SOLVER-CHART(G')
if there is an entry for G' in the chart
   then return true
free ← FREE-FRAGMENTS(G')
if free = ∅
   then return false
if G' contains only one fragment
   then return true

for each F ∈ free
do split ← SPLIT(G', F)
      for each S ∈ WCCS(G' − F)
      do if GRAPH-SOLVER-CHART(S) = false
            then return false
      add (G', split) to the chart
return true
```


An Example

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      for each S ∈ WCCS(G' − F)
      do if GRAPH-SOLVER-CHART(S) = false
            then return false
      add (G', split) to the chart
return true
```

subgraph	splits
{1,...,7}	1({4}, {2,3,5,6,7})
	2({1,4,5}, {3,6,7})
	3({1,2,4,5,6}, {7})
{2,3,5,6,7}	2({5}, {3,6,7})
	3({2,5,6}, {7})
{1,2,4,5,6}	1({4}, {2,5,6})
	2({1,4,5}, {6})
{2,5,6}	2({5}, {6})
{3,6,7}	3({7}, {6})
{1,4,5}	1({4}, {5})

Complexity

- Let G be a dominance graph with n nodes and m edges
- The computation of free fragments takes time O(n + m)
- The time to compute the chart is O(n(n + m) wcsg(G))
  - wcsg(G) = the number of weakly connected subgraphs of G
- Worst case complexity in the number of nodes is $O(n^2 2^n)$
  - ⇒ big improvement compared to $O(n^2 n!)$ of the basic solver
  - (n! = upper bound for the number of solved forms)

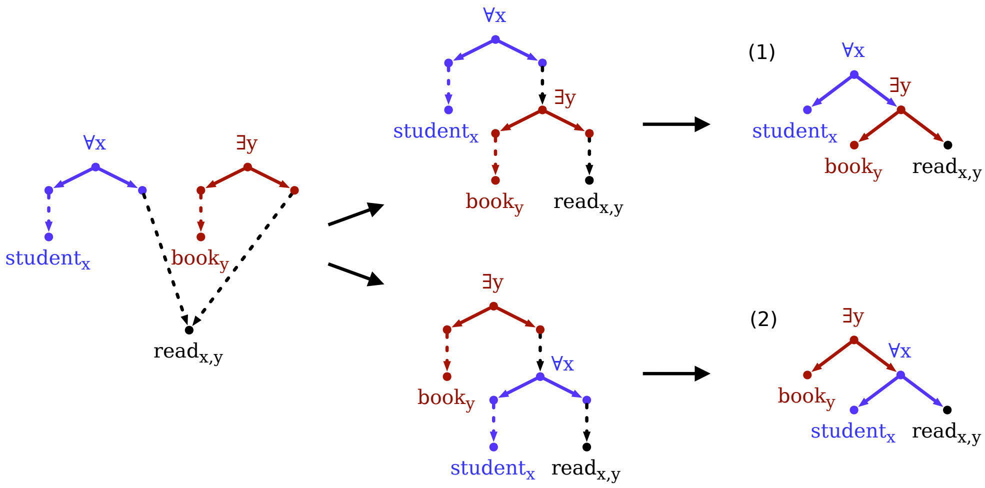
Hypernormally Connected Dominance Graphs

The big picture

- *Every student reads a book*

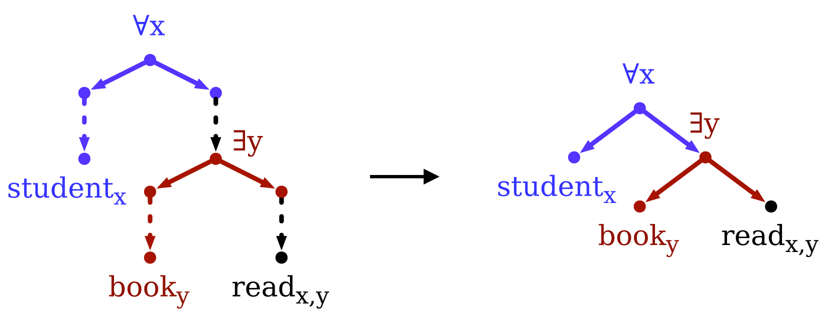
  (1) ∀x(student(x), ∃y(book(y), read(x, y))

  (2) ∃y(book(y), ∀x(student(x), read(x, y))



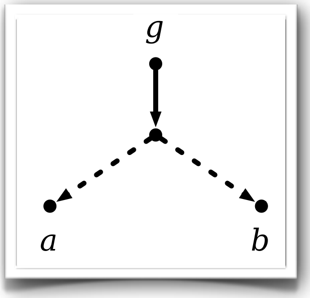
Simple solved forms

- We are usually interested only in solved forms in which every hole is related to exactly one root
  - let's call such solved forms "simple"
- Simple solved forms can be mapped to "proper" trees simply by "plugging" the hole with the root connected to it by a dominance edge.



Not all solved forms are simple

- **Problem:**
  Not all solved forms are simple
- **Solution:**
  Identify a class of dominance graphs that only have simple solved forms



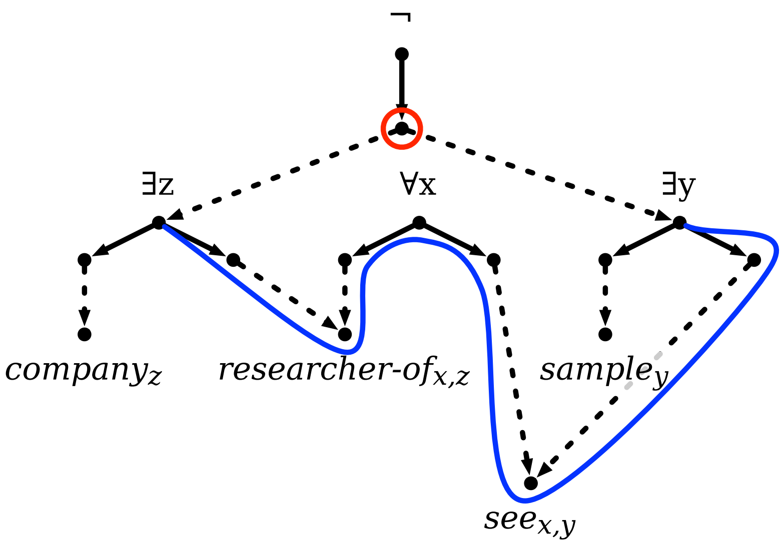
- ⇒ Hypernormally connected dominance graphs

Hypernormally connected dominance graphs

- A **hypernormal path** in a (normal) dominance graph G is a path in the undirected version of G that does not use two dominance edges incident to the same hole.
- A (normal) dominance graph G is **hypernormally connected** if each pair of nodes is connected by some hypernormal path.

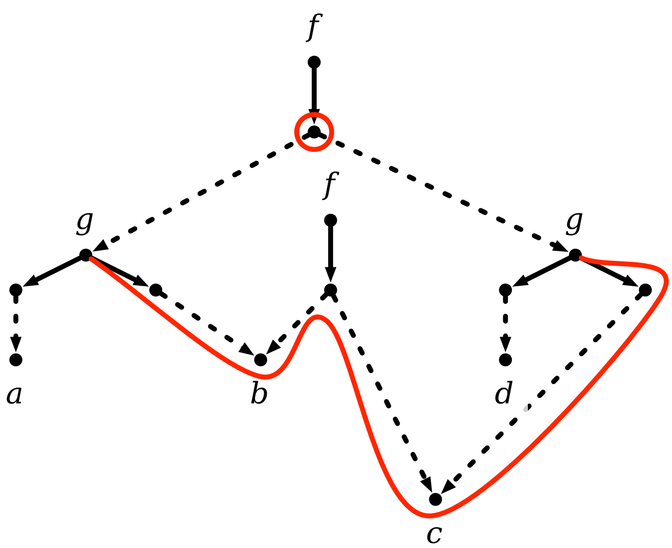
Hypernormally connected dominance graphs

- Hypernormally connected:



Hypernormally connected dominance graphs

- Not hypernormally connected:



Hypernormally connected dominance graphs

- **Lemma:** if G is a hypenormally connected normal dominance graph with free fragment F, then all WCCs of G\F are hypernormally connected.
- **Proposition:** if a normal dominance graph is hypernormally connected, then all its (minimal) solved forms are simple.

Tree Automata

(see Comon &al., 2007)

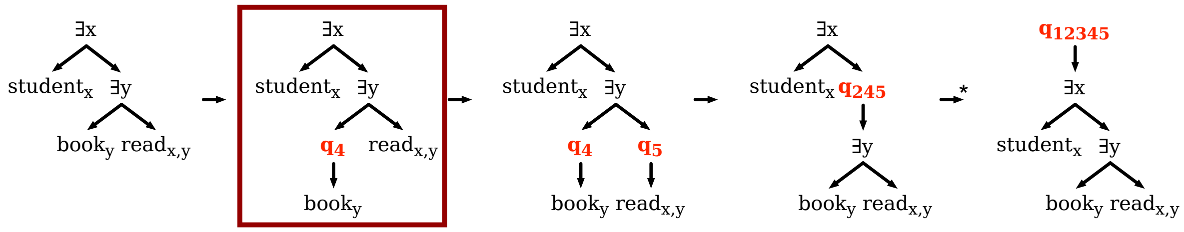
Bottom-up Tree Automaton

- **A tree automaton** is a tuple $A = \langle Q, \Sigma, Q_f, \Delta \rangle$
  - Σ a finite ranked signature
  - Q a finite set of states
  - $Q_f \subseteq Q$ a finite set of final (accepting) states
  - Δ a finite set of transition rules
- **Transition rules:**
  - $f(q_1(x_1), \ldots, q_n(x_n)) \to q(f(x_1, \ldots, x_n))$
    - $f \in \Sigma$
    - $q, q_1, \ldots, q_n \in Q$
    - $x_1, \ldots, x_n$ different variables

An Example Computation

- $Q = \{q_3, q_4, \ldots\}$, $\Sigma = \{\exists x_{|2}, \text{book}_{y|0}, \ldots\}$, $Q_f = \{q_{12345}\}$

$$\exists x(q_3(x_1), q_{245}(x_2)) \to q_{12345}(\exists x(x_1, x_2))$$
$$\exists y(q_4(x_1), q_5(x_2)) \to q_{245}(\exists y(x_1, x_2))$$
$$\text{student}_x \to q_3(\text{student}_x)$$
$$\text{book}_y \to q_4(\text{book}_y)$$
$$\text{read}_{x,y} \to q_5(\text{read}_{x,y})$$



Bottom-up Tree Automaton

- A tree t is **accepted** by an atomaton A = $\langle Q, \Sigma, Q_f, \Delta \rangle$ if
  - $t \rightarrow^{*} q(t)$
  - $q \in Q_f$
- **The language L(A)** of trees recognized by A is the set of trees accpeted by A.

Another Example

- **Automaton** (rules)

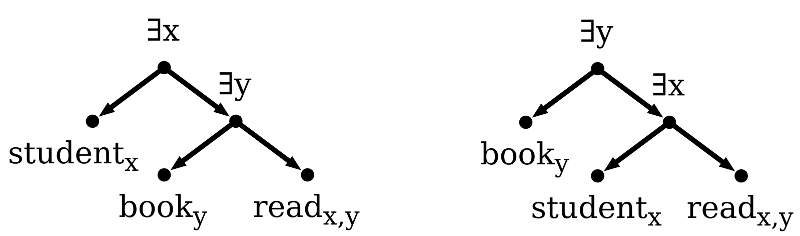
$\exists x(q_3(x_1), q_{245}(x_2)) \rightarrow q_{12345}(\exists x(x_1, x_2))$ $\quad$ $student_x \rightarrow q_3(student_x)$

$\exists y(q_4(x_1), q_{135}(x_2)) \rightarrow q_{12345}(\exists y(x_1, x_2))$ $\quad$ $book_y \rightarrow q_4(book_y)$

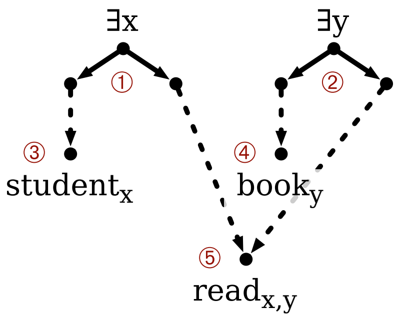
$\exists y(q_4(x_1), q_5(x_2)) \rightarrow q_{245}(\exists y(x_1, x_2))$ $\quad$ $read_{x,y} \rightarrow q_5(read_{x,y})$

$\exists x(q_3(x_1), q_5(x_2)) \rightarrow q_{135}(\exists x(x_1, x_2))$

- **Accepted Trees:**



Back to dominance charts



dominance chart	
{1,2,3,4,5}	1({3}, {2,4,5})
	2({4}, {1,3,5})
{2,4,5}	2({4}, {5})
{1,3,5}	1({3}, {5})

Compare

tree automaton	
$\exists x(q_3(x_1), q_{245}(x_2))$	$\to q_{12345}(\exists x(x_1, x_2))$
$\exists y(q_4(x_1), q_{135}(x_2))$	$\to q_{12345}(\exists y(x_1, x_2))$
$\exists y(q_4(x_1), q_5(x_2))$	$\to q_{245}(\exists y(x_1, x_2))$
$\exists x(q_3(x_1), q_5(x_2))$	$\to q_{135}(\exists x(x_1, x_2))$
$student_x$	$\to q_3(student_x)$
$book_y$	$\to q_4(book_y)$
$read_{x,y}$	$\to q_5(read_{x,y})$

dominance chart	
{1,2,3,4,5}	1({3}, {2,4,5})
	2({4}, {1,3,5})
{2,4,5}	2({4}, {5})
{1,3,5}	1({3}, {5})

Charts = Tree automata

- Dominance charts can be seen as (or translated into) tree automata
  - … provided that the original dominance graph is hypernormally connected (why?)

Charts = Tree automata

- The class of recognizable languages is closed under
  - Union
  - Complement
  - Intersection
- ⇒ We can model certain inferences on the level of dominance charts by intersecting regular tree languages
  - Redundancy elimination
  - Weakest readings

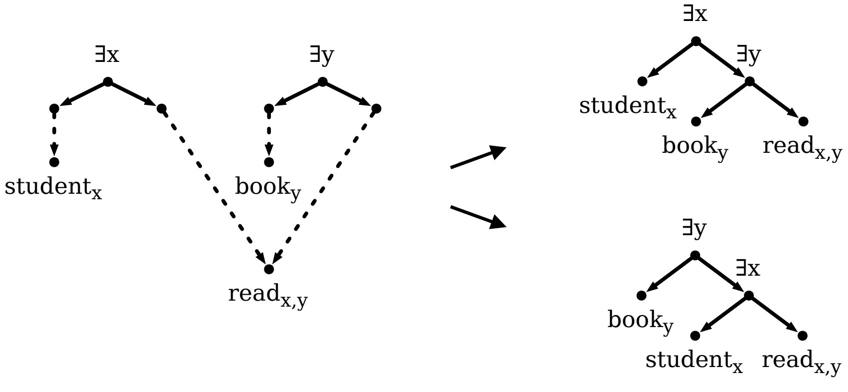
Redundancy Elimination (very short)

Ambiguity, revisited

- *A student reads a book*

  (1) ∃x(student(x), ∃y(book(y), read(x, y))

  (2) ∃y(book(y), ∃x(student(x), read(x, y))

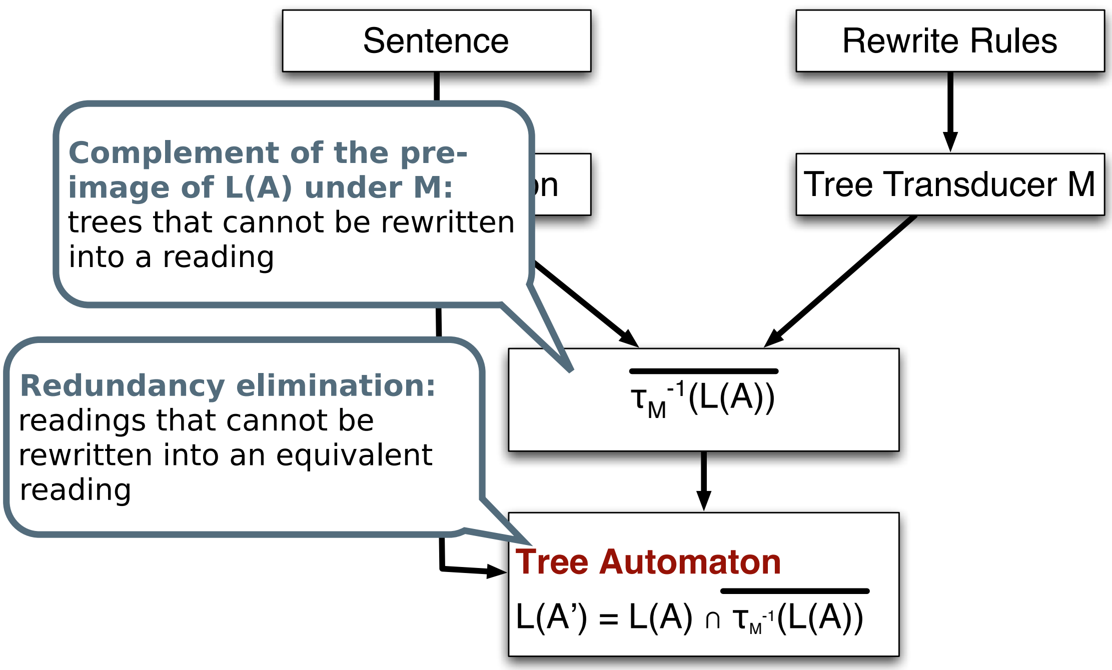


Redundancy Elimination

- *A student reads a book*
  (1) ∃x(student(x), ∃y(book(y), read(x, y))
  (2) ∃y(book(y), ∃x(student(x), read(x, y))
- Readings (1) and (2) are logically equivalent!
- Basic idea:
  - Model the relation between (1) and (2) by rewrite rules
  - Translate the rewrite rules into a tree automaton
  - Redundancy elimination = Intersection of regular tree languages

(Koller & Thater, 2010)

Redundancy Elimination



References

- Alexander Koller and Stefan Thater (2005). The evolution of dominance constraint solvers [PDF]. In Proceedings of the ACL Workshop on Software.
- Alexander Koller and Stefan Thater (2010). Computing Weakest Readings [PDF]. In Proceedings of the 48th Annual Meeting of the Association for Computational Linguistics.
- Hubert Comon, Max Dauchet, Remi Gilleron, Florent Jacquemard, Denis Lugiez, Christof Löding, Sophie Tison, Marc Tommasi. Tree Automata, Techniques and Applications. [PDF]